Topics in Matroid Union

by

Ali M Hameed

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Abstract

The operation of matroid union was introduced by Nash-Williams in 1966. A matroid is indecomposable if it cannot be written in the form $M = M_1 \vee M_2$, where $r(M_1), r(M_2) > 0$. In 1971 Welsh posed the problem of characterizing indecomposable matroids, this problem has turned out to be extremely difficult. As a partial solution towards its progress, Cunningham characterized binary indecomposable matroids in 1977. In this thesis we present numerous results in topics of matroid union. Those include a link between matroid connectivity and matroid union, such as the implication of having a 2-separation in the matroid union, and under what conditions is the union 3connected. We also identify which elements in binary and ternary matroids are non-fixed. Then we create a link between having non-fixed elements in binary and ternary matroids and the decomposability of such matroids, and the effect of removing non-fixed elements from binary and ternary matroids. Moreover, we show results concerning decomposable 3-connected ternary matroids, such as what essential property every decomposable 3-connected ternary matroid must have, how to compose a ternary matroid, and what a 3-connected ternary matroid decomposes into. We also give an alternative statement and an alternative proof of Cunningham's theorem from the perspective of fixed and non-fixed elements.

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Introduction

The operation of matroid union was introduced by Nash-Williams in 1966. A matroid is indecomposable if it cannot be the union of matroids. In 1971 Welsh posed the problem of characterizing indecomposable matroids, this problem has turned out to be extremely difficult. As a partial solution towards its progress, Cunningham characterized binary indecomposable matroids in 1977. In this thesis I present numerous results in topics of matroid union. We start by looking at general properties of matroid union, then proceed to show the link between matroid connectivity and matroid union. Then we look at properties of non-fixed elements in binary and ternary matroids, and how the presence of those elements can inform us about the decomposability of those matroids. Finally we show some results concerning decomposable 3-connected ternary matroids. The concept of fixed and non-fixed elements plays an essential role in proving many results in this thesis.

This thesis has been written with the assumption that the reader is familiar with the basic definitions, ideas and concepts of Matroid Theory. Although Chapter 1 contains all the preliminaries needed, it is only for the purposes of this thesis, and it is far from being an exhaustive reference for matroid theory as a whole. For an exhaustive reference on the basics of matroid theory, the reader is encouraged to read the early sections of [5].

The first section of Chapter 2 introduces the operation of matroid union in a formal way. The second section discusses the important issue of the ground sets of the matroids whose union we are considering. The third section explores some properties of matroid union, and finally we show which minors of a decomposable matroid are also decomposable.

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Chapter 3 discusses the relationship between matroid union and connectivity. First we note why we are not interested in disconnected unions. Then we look at the rank of the connected union, the link with series connections, and what it means to have a 2-separation in a matroid union. The final section answers the question: 'When is the union 3-connected?'.

Chapter 4 investigates properties of fixed and non-fixed elements in the union of two matroids in general, and in binary and ternary matroids in particular. Cunningham's result is restated and reproved in terms of fixed and non-fixed elements, so that an alternative picture of decomposable binary matroids is presented.

Chapter 5 contains results relating to decomposable 3-connected ternary matroids, including how to compose a ternary matroid. Finally, Chapter 6 presents some conjectures that may be catalysts for further research in the problem of matroid union.

A final note on which results in this thesis are original and which are known. Chapter 1 contains the preliminary information needed for this thesis, so everything in this chapter is known. In Chapter 2, everything prior to Section 2.4 is known, the results of Section 2.4 are not found in the literature, however, they are straightforward. The material of Sections 3.1 and 3.2 of Chapter 3 is known, and the remaining sections of this chapter contain new material. The definitions and results of Section 4.1 of Chapter 4 are known. Section 4.2 contains new material. Theorem 4.3.3 and Corollary 4.3.5 of Section 4.3 are new, and although Corollary 4.3.2 is straightforward, it is not currently present in the literature. The rest of the material in this section is known. The material in Section 4.4 is known but is restated and re-proven using the new material in Section 4.2, and using Lemma 4.3.1 which is known. The results of Sections 5.1, 5.2 and 5.3 in Chapter 5 are new, apart from Proposition 5.3.6 and Theorem 5.3.8. Propositions 5.4.2 - 5.4.4 and Theorem 5.4.5 of Section 5.4 are known, and the rest of the material in the section is new.

Chapter 1

Matroids and matroid operations

This chapter catalogues some basic facts of matroid theory that will be needed later on in this thesis.

1.1 Independent sets and circuits

A matroid M is an ordered pair (E,\mathcal{I}) consisting of a finite set E and a collection \mathcal{I} of subsets of E satisfying the following three conditions:

- (I1) $\emptyset \in \mathcal{I}$.
- (I2) If $I \in \mathcal{I}$ and $I' \subseteq I$, then $I' \in \mathcal{I}$.
- (I3) If I_1 and I_2 are in \mathcal{I} and $|I_1| < |I_2|$, then there is an element e of $I_2 I_1$ such that $I_1 \cup e \in \mathcal{I}$.

Condition (I3) is called the independence augmentation axiom.

If M is the matroid (E, \mathcal{I}) , then M is a matroid on E. The members of \mathcal{I} are the *independent sets* of M, and E is the *ground set* of M. A subset of E that is not in \mathcal{I} is a *dependent* set. A *basis* of M is a maximal independent set.

A *circuit* of an arbitrary matroid M is a minimal dependent set all of

whose proper subsets are independent, and we shall denote the set of circuits of M by C or C(M). Circuits have the following properties:

- (C1) $\emptyset \notin \mathcal{C}$.
- (C2) If C_1 and C_2 are members of \mathcal{C} and $C_1 \subseteq C_2$, then $C_1 = C_2$.
- (C3) If C_1 and C_2 are distinct members of \mathcal{C} and e in $C_1 \cap C_2$, then there is a member C_3 of \mathcal{C} such that $C_3 \subseteq (C_1 \cup C_2) e$.

Condition (C3) is the *circuit elimination axiom*. Next we state some facts about matroids, but first we recall what it means for a matroid to be representable over a field F.

The two fundamental classes of matroids arise from matrices and from graphs. A vector matroid is the matroid that arises from a matrix, and a graphic matroid is the matroid that arises from a graph. If M is isomorphic to the vector matroid of a matrix D over a field F, then M is said to be representable over F or F-representable.

1.2 Rank

In this section we define what a rank of a matroid is, and state a fact about ranks. First we define a fundamental matroid construction. Let M be the matroid (E,\mathcal{I}) and suppose that $X\subseteq E$. Let $\mathcal{I}|X$ be $\{I\subseteq X:I\in\mathcal{I}\}$. Then the pair (X,\mathcal{I}) is a matroid. This matroid is the restriction of M to X or the deletion of E-X from M. It is denoted by M|X or $M\setminus (E-X)$. The rank r(X) of X is the size of a basis B of M|X. Evidently r maps 2^E into the set of non-negative integers. This function, the rank function of M, will be written as r_M or r.

Lemma 1.2.1. [5, Lemma 1.3.1] The rank function r of a matroid M on a set E satisfies the following condition:

(R3) If X and Y are subsets of E, then

$$r(X \cup Y) + r(X \cap Y) \le r(X) + r(Y). \tag{1.2.1}$$

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1.3 Duality

In this section we define the dual of a matroid and state a useful property. Note that $\mathcal{B}(M)$ is the set of bases of a matroid M.

Theorem 1.3.1. [5, Theorem 2.1.1] Let M be a matroid and $B^*(M)$ be $\{E(M) - B : B \in \mathcal{B}(M)\}$. Then $B^*(M)$ is the set of bases of a matroid on E(M).

The matroid obtained in the theorem above, whose ground set is E(M), and whose set of bases is $B^*(M)$ is the dual of M, and is denoted by M^* . The circuits, hyperplanes, independent sets, and spanning sets of M are the cocircuits, cohyperplanes, coindependent sets, and cospanning sets of M^* respectively.

Proposition 1.3.2. [5, Proposition 2.1.6] Let M be a matroid on a set E and suppose $X \subseteq E$. Then

- (i) X is independent if and only if E X is cospanning.
- (ii) X is spanning if and only if E X is coindependent.
- (iii) X is a hyperplane if and only if E X is a cocircuit.
- (iv) X is a circuit if and only if E X is a cohyperplane.

1.4 Minors

In this section we define what minors are. If M is a matroid on E and T is a subset of E, then the *contraction* of T from M is denoted M/T and is given by $M/T = (M^* \setminus T)^*$. A sequence of deletions and contractions from M can be written in the form $M \setminus X/Y$ for some pair of disjoint sets X and Y, where $X,Y \subseteq E(M)$, either of which may be empty. Matroids of this form are *minors* of M.

Proposition 1.4.1. [5, Corollary 3.1.8] Suppose that B_T is a basis of M|T. Then

```
\mathcal{I}(M/T) = \{ I \subseteq E - T : I \cup B_T \in \mathcal{I}(M) \}
= \{ I \subseteq E - T : M \mid T \text{ has a basis } B \text{ such that } B \cup I \in \mathcal{I}(M) \}.
```

Proposition 1.4.2. [5, Proposition 3.1.26] Let T_1 and T_2 be disjoint subsets of E(M). Then

- (i) $(M \setminus T_1) \setminus T_2 = M \setminus (T_1 \cup T_2) = (M \setminus T_2) \setminus T_1$;
- (ii) $(M/T_1)/T_2 = M/(T_1 \cup T_2) = (M/T_2)/T_1$; and
- (iii) $(M \setminus T_1)/T_2 = (M/T_2) \setminus T_1$.

1.5 Connectivity

In this section we define connectivity for matroids and state some important facts about connectivity.

Proposition 1.5.1. [5, Proposition 4.1.4] The matroid M is connected if and only if, for every pair of distinct elements of E(M), there is a circuit containing both.

The following observation follows from the above.

Proposition 1.5.2. A connected matroid with at least two elements does not have any coloops.

Proof. Suppose M is a connected matroid that has a coloop e. By Proposition 1.5.1, there must be a circuit containing e and any other element f of E(M). But since a coloop is not contained in any circuit, then e is not contained in any circuit, giving a contradiction. Therefore a connected matroid does not have any coloops.

Similarly, the following is true for connected matroids.

Proposition 1.5.3. A connected matroid with at least two elements does not have any loops.

Proof. Suppose M is a connected matroid that has a loop e. By Proposition 1.5.1, there must be a circuit containing e and any other element f of E(M). But a loop is a 1-element circuit, and therefore is not contained in any circuit other than itself. Hence e is not contained in any circuit other than itself, giving a contradiction. Therefore a connected matroid does not have any loops.

The following proposition links the connectivity of a matroid with its dual.

Proposition 1.5.4. [5, Corollary 4.2.8] M is connected if and only if M^* is connected.

The following proposition describes the structure of a disconnected matroid.

Proposition 1.5.5. [5, Proposition 4.2.11] The matroid M is disconnected if and only if, for some proper non-empty subset T of E(M), $\mathcal{I}(M) = \{I_1 \cup I_2 : I_1 \in \mathcal{I}(M|T), I_2 \in \mathcal{I}(M|E-T)\}$.

1.6 Direct sums

In this section we define and state facts about direct sums. Note the following.

Proposition 1.6.1. [5, Proposition 4.2.12] Let M_1 and M_2 be matroids on disjoint sets E_1 and E_2 . Let $E = E_1 \cup E_2$ and $\mathcal{I} = \{I_1 \in \mathcal{I}(M_1), I_2 \in \mathcal{I}(M_2)\}$. Then (E, \mathcal{I}) is a matroid.

The matroid (E,\mathcal{I}) in the proposition above is the *direct sum* or 1-sum of M_1 and M_2 and is denoted $M_1 \oplus M_2$.

What follows are equivalent ways of describing direct sums.

Proposition 1.6.2. [5, Proposition 4.2.16] C is a circuit of $M_1 \oplus M_2$ if and only if C is either a circuit of M_1 or a circuit of M_2 .

Proof. Let C be a circuit of $M_1 \oplus M_2$, then we will show that C is either in $C(M_1)$ or in $C(M_2)$. If for all $x \in C$, x was in $E(M_1)$ then C is in $C(M_1)$, and if for all $x \in C$, x was in $E(M_2)$ then C is in $C(M_2)$ and we are done. So suppose the elements of C are from $E(M_1)$ and from $E(M_2)$. Without loss of generality, choose $x \in C$, such that x is in $E(M_1)$. Then $C - \{x\} = I$ is independent in $M_1 \oplus M_2$, and hence of the form $I_1 \cup I_2$ where I_1 is independent in I_2 is independent in I_3 . It is also true that for any other element $I_3 \oplus I_4$ is independent. But $I_4 \oplus I_4$ contains I_5 and contains I_6 . Since I_6 is a proper subset of I_6 , then

it is an independent set of M_1 . Meaning $I' = I_1 \cup \{x\}$ is independent in M_1 , because I' is a proper subset of C. So that $I' \cup I_2$ must be independent in $M_1 \oplus M_2$, but $I' \cup I_2 = C$, contradiction. Therefore if C is a circuit of $M_1 \oplus M_2$, then it is either in $C(M_1)$ or in $C(M_2)$.

Conversely, if C is either a circuit of M_1 or M_2 , then its elements are either completely from $E(M_1)$ or completely from $E(M_2)$ respectively. Suppose without loss of generality that all its elements are from $E(M_1)$. Consider C in $M_1 \oplus M_2$. C is either dependent or independent in $M_1 \oplus M_2$. It cannot, however, be independent because then it has to be independent in M_1 when of course it is not. So it must be dependent in $M_1 \oplus M_2$. Moreover, every proper subset of C is independent in M_1 , because C is a circuit of M_1 . It follows that every proper subset of C is independent in $M_1 \oplus M_2$ as well. Therefore C is a minimal dependent set in $M_1 \oplus M_2$, that is a circuit of $M_1 \oplus M_2$.

The rank of a set in the direct sum is calculated as follows.

Proposition 1.6.3. [5, Proposition 4.2.17] If $X \subseteq E(M_1 \oplus M_2)$ then $r_{M_1 \oplus M_2}(X) = r_{M_1}(X \cap E(M_1)) + r_{M_2}(X \cap E(M_2))$.

Proof. Let Y be the maximal independent subset of X. Then Y can be partitioned into the two independent sets $Y_1 \subseteq E(M_1)$ and $Y_2 \subseteq E(M_2)$. But in order for Y to be a maximal independent set, both Y_1 and Y_2 have to be maximal independent sets in M_1 and M_2 respectively. It follows that $|Y_1| = r_{M_1}(X \cap E(M_1))$, and $|Y_2| = r_{M_2}(X \cap E(M_2))$. Therefore $r_{M_1 \oplus M_2}(X) = |Y| = |Y_1| + |Y_2| = r_{M_1}(X \cap E(M_1)) + r_{M_2}(X \cap E(M_2))$.

Next we look at the bases in the direct sum.

Proposition 1.6.4. [5, Proposition 4.2.18] The set of bases of $M_1 \oplus M_2$ are $\{B_1 \cup B_2 : B_1 \in \mathcal{B}(M_1), B_2 \in \mathcal{B}(M_2)\}.$

Proof. Let B be a basis of $M_1 \oplus M_2$. B is a maximal independent set. So it follows from Proposition 1.6.1, that B is the union of two independent sets I_1 and I_2 , where I_1 is independent in M_1 and I_2 independent in M_2 . But in order for B to be a maximal independent set, both I_1 and I_2 have to be maximal independent sets in M_1 and M_2 respectively. So that I_1 and I_2 are bases of M_1 and M_2 respectively.

The flats of the direct sum are as follows.

Proposition 1.6.5. [5, Proposition 4.2.20] F is a flat of $M_1 \oplus M_2$ if and only if $F \cap E(M_1)$ and $F \cap E(M_2)$ are flats of M_1 and M_2 respectively.

Proof. Let $F_{M_1 \oplus M_2}$ denote a flat of $M_1 \oplus M_2$. If F is a flat of $M_1 \oplus M_2$, then cl(F) = F. Let $F = F_1 \cup F_2$, where F_1 and F_1 are disjoint sets such that $F_1 = F \cap E(M_1)$, and $F_2 = F \cap E(M_2)$. Then $cl(F) = F = F_1 \cup F_2$. By Proposition 1.6.3, any increase in the rank of F_1 , will increase the rank of $F_{M_1 \oplus M_2}$, and any increase in the rank of F_2 will increase the rank of $F_{M_1 \oplus M_2}$. So it follows that $cl(F) = cl_{M_1}(F_1) \cup cl_{M_2}(F_2)$. So now we have $cl(F) = F = F_1 \cup F_2 = cl_{M_1}(F_1) \cup cl_{M_2}(F_2)$. Since $cl_{M_1}(F_1)$ only contains elements from $E(M_1)$, then $cl_{M_1}(F_1)$ must be equal to F_1 . Similarly $cl_{M_2}(F_2)$ must be equal to F_2 . Therefore $F \cap E(M_1)$ and $F \cap E(M_2)$ are flats of M_1 and M_2 respectively.

Conversely, suppose $F \cap E(M_1)$ is a flat of M_1 and $F \cap E(M_2)$ is a flat of M_2 . Let $F_1 = F \cap E(M_1)$ and $F_2 = F \cap E(M_2)$, and let $F = F_1 \cup F_2$. Since F_1 and F_2 are flats of M_1 and M_2 respectively, then $cl_{M_1}(F_1) = F_1$ and $cl_{M_2}(F_2) = F_2$. So $F_1 \cup F_2 = F = cl_{M_1}(F_1) \cup cl_{M_2}(F_2)$. But the elements that you can add to F_1 without increasing its rank in M_1 , together with the elements that you can add to F_2 without increasing its rank in M_2 , are the elements that you can add to F without increasing its rank in $M_1 \oplus M_2$. In other words $cl_{M_1}(F_1) \cup cl_{M_2}(F_2) = F = cl(F)$. Therefore F is a flat of $M_1 \oplus M_2$.

The closure operator of the direct sum is as follows.

Proposition 1.6.6. [11, Proposition 7.6.1]. If X is a set in $E(M_1 \oplus M_2)$, then the closure of X in the direct sum is $cl_{M_1}(X_1) \cup cl_{M_2}(X_2)$ for some partition $\{X_1, X_2\}$ of X, where $X = X_1 \cup X_2$, $X_1 = X \cap E(M_1)$, and $X_2 = X \cap E(M_2)$.

Proof. By Proposition 1.6.3 $r_{M_1 \oplus M_2}(X) = r_{M_1}(X_1) + r_{M_2}(X_2)$. So $cl(X) = \{x : r(X \cup \{x\}) = r(X) \text{ for all } x \in E(M_1 \oplus M_2)\}$. Element x can be either in $E(M_1)$, or in $E(M_2)$. So $cl(X) = \{x_1 : r(X_1 \cup \{x_1\}) = r(X_1)\}$ for $x_1 \in E(M_1)\} \cup \{x_2 : r(X_2 \cup \{x_2\}) = r(X_2) \text{ for } x_2 \in E(M_2)\}$. The aforementioned is equal to $cl_{M_1}(X_1) \cup cl_{M_2}(X_2)$.

1.7 Representable matroids

In this section we state some facts about representable matroids. Recall from Section 1.1 that an F-representable matroid, is a matroid isomorphic to the vector matroid of a matrix D over the field F. A matroid is binary if it is representable over GF(2), a matroid is ternary if it is representable over GF(3), a matroid is regular if it is both binary and ternary.

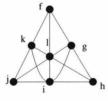
Proposition 1.7.1. [5, Corollary 2.2.9] If M is representable over the field F, then M^* is also representable over F.

Proposition 1.7.2. [5, Proposition 3.2.4] Every minor of an F-representable matroid is F-representable.

The next result lists some classes of matroids that are closed under direct sums. Please note that transversal matroids will be defined later.

Proposition 1.7.3. [5, Proposition 4.2.15] The classes of F-representable, graphic, cographic, transversal, and regular matroids are all closed under the operation of direct sum.

One of the important matroids in Matroid Theory, that will be referred to a lot in this thesis is known as the Fano matroid, denoted F_7 and shown in Figure 1.1 below. The next proposition shows when F_7 is representable.



The Fano matroid F7

Figure 1.1: The Fano matroid

Theorem 1.7.4. [5, Theorem 6.4.8] Let F be a field. Then F_7 is F-representable if and only if the characteristic of F is two.

The remaining results of this section show what conditions are needed for a matroid to be binary, ternary or regular. **Theorem 1.7.5.** [5, Theorem 6.5.4] A matroid is binary if and only if it has no $U_{2,4}$ -minor.

Theorem 1.7.6. [5, Theorem 6.5.7] A matroid is ternary if and only if it has no minor isomorphic to any of the matroids $U_{2.5}, U_{3.5}, F_7$ or F_7^* .

Theorem 1.7.7. [5, Theorem 6.6.3] The following statements are equivalent for a matroid M:

- (i) M is regular.
- (ii) M is representable over every field.
- (iii) M is binary and, for some field F of characteristic other than two, M is F-representable.

Theorem 1.7.8. [5, Theorem 6.6.4] A matroid is regular if and only if it has no minor isomorphic to any of the matroids $U_{2,4}$, F_7 or F_7^* .

1.8 Series and parallel connections

In this section we state some facts about series and parallel connections. We also define and state facts about the 2-sum operation. We assume the reader is familiar with series connections and parallel connections of matroids, for background information on this subject see [5, Section 7.1].

Let M_1 and M_2 be two matroids with $E(M_1) \cap E(M_2) = \{p\}$. Let $E = E(M_1) \cup E(M_2)$. If the collection of circuits of the series connection of M_1 and M_2 is C_S , and the collection of circuits of the parallel connection of M_1 and M_2 is C_P , then the following is true:

Proposition 1.8.1. [5, Theorem 7.1.4] Each of C_S and C_P is the collection of circuits of a matroid on E.

The matroids on E that have C_S and C_P as their sets of circuits will be denoted by $S(M_1, M_2)$ and $P(M_1, M_2)$ respectively. In this context, p is the base-point of the connections. What follows are facts about the base-point, series and parallel connections, and the link between series connections and direct sums.

Proposition 1.8.2. [5, Theorem 7.1.8] If the base-point p is a coloop of M_1 , we define $S(M_1, M_2) = M_1 \oplus (M_2 \setminus p)$.

Proposition 1.8.3. [5, Theorem 7.1.11] The base-point p is a coloop of $S(M_1, M_2)$ if and only if it is a coloop of M_1 or M_2 .

Proposition 1.8.4. [5, Theorem 7.1.12] The base-point p is a loop of $P(M_1, M_2)$ if and only if it is a loop of M_1 or M_2 .

Proposition 1.8.5. [5, Proposition 7.1.13] Let M_1 and M_2 be matroids with $E(M_1) \cap E(M_2) = \{p\}$. Let $E = E(M_1) \cup E(M_2)$ and B be a subset of E.

- (i) Assume that, in at least one of M_1 and M_2 , p is not a coloop. Then B is a basis of $S(M_1, M_2)$ if and only if B can be written as a disjoint union of bases of M_1 and M_2 .
- (ii) Assume that, in at least one of M_1 and M_2 , p is not a loop. Then B is a basis of $P(M_1, M_2)$ containing p if and only if $B \cap E(M_i)$ is a basis of M_i containing p for i = 1,2. Moreover, B is a basis of $P(M_1, M_2)$ not containing p if and only if $p \notin B$ and, for some distinct i and j in $\{1,2\}$, the set $[B \cap E(M_i)] \cup p$ and $B \cap E(M_j)$ are bases of M_i and M_j , respectively.

Theorem 1.8.6. [5, Theorem 7.1.16] Let p be an element of a connected matroid M.

(i) If $M \setminus p = M_1 \oplus M_2$ where both M_1 and M_2 are non-empty, then $M = S(M/E(M_1), M/E(M_2))$.

Dually,

(ii) If $M/p = M_1 \oplus M_2$ where both M_1 and M_2 are non-empty, then $M = P(M \setminus E(M_1), M \setminus E(M_2))$.

A matroid is said to have a k-separation if there is a partition (X, Y) of the ground set E such that:

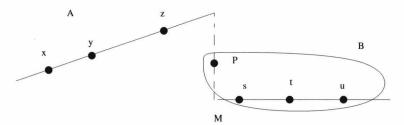
- (i) |X| and |Y| are each at least k, and
- (ii) r(X) + r(Y) r(E) is less than or equal to k 1. The smallest k such that M has a k-separation is called the *connectivity* of M. If k is at

least 2, we say that M is n-connected for all integers n no larger than k. If k = 1, then we say that M is disconnected. So in general:

- (i) If a matroid has k-separation then it is not (k+1)-connected.
- (ii) If a matroid is k-connected then it is (k-1)-connected.

The next example illustrates parts (i) and (ii) of Theorem 1.8.6.

Example 1.8.7. Figures 1.2 and 1.3 illustrate parts (i) and (ii) of Theorem 1.8.6 respectively. In Figure 1.2 M is a series connection on base-point p. Let $A = \{x, y, z\}, B = \{s, t, u, p\}$. So that $r_M(A) + r_M(B) = 2 + 3 = 5 = r(M) + 1 = 4 + 1$, thus indicating that (A, B) is a 2-separation of M. The series connection can be disconnected by deleting the base-point element p. So in $M \setminus p$ we get $r_{M \setminus p}(A) + r_{M \setminus p}(B - \{p\}) = 2 + 2 = 4 = r(M)$, thus indicating that $(A, B - \{p\})$ is a 1-separation of $M \setminus p$.



Series connections get disconnected by deleting the basepoint element

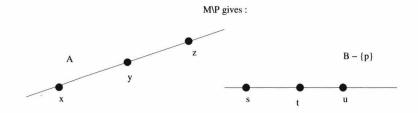


Figure 1.2: Series connection and deletion.

Figure 1.3 shows how a parallel connection can be disconnected by contracting the base-point p. Let $A = \{x, y, z\}, B = \{s, t, u, p\}$. So that $r_M(A) + r_M(B) = 2 + 2 = 4 = r(M) + 1 = 3 + 1$, thus indicating that (A, B) is a 2-separation of M. The parallel connection can be disconnected by contracting the base-point element p. So in M/p we get $r_{M/p}(A) + r_{M/p}(B - \{p\}) = 1 + 1 = 2 = r(M/p)$, thus indicating that $(A, B - \{p\})$ is

a 1-separation of M/p.



M is a parallel connection on basepoint p

Figure 1.3: Parallel connection and contraction.

The following proposition shows that the connectivity of the series and parallel connection of two matroids depends on the connectivity of the two matroids.

Proposition 1.8.8. [5, Proposition 7.1.17] If the matroids M_1 and M_2 each have at least two elements and $E(M_1) \cap E(M_2) = \{p\}$, then the following statements are equivalent:

- (i) Both M_1 and M_2 are connected.
- (ii) $S(M_1, M_2)$ is connected.
- (iii) $P(M_1, M_2)$ is connected.

Let M and N be matroids each with at least three elements, and let $E(M) \cap E(N) = \{p\}$. Suppose that neither M nor N has $\{p\}$ as a separator. Then the 2-sum of M and N denoted $M \oplus_2 N$ is $S(M_1, M_2)/p$ or, equivalently, $P(M_1, M_2) \setminus p$. Clearly $M \oplus_2 N = N \oplus_2 M$. M and N are the parts of the 2-sum. What follows are results about 2-sums.

Proposition 1.8.9. [5, Proposition 7.1.19] Both M and N are isomorphic to proper minors of $M \oplus_2 N$.

Proposition 1.8.10. [5, Proposition 7.1.20] (i) $(M \oplus_2 N)^* = M^* \oplus_2 N^*$. (ii) If $|E(N_i)| \geq 2$ for i = 1, 2, then $P(N_1, N_2) \setminus p$ is connected if and only if both N_1 and N_2 are connected. In particular, $M \oplus_2 N$ is connected if and only if both M and N are connected.

Some classes of matroids are closed under 2-sums, as shown in the next result.

Proposition 1.8.11. [5, Corollary 7.1.23] The classes of graphic, cographic, F-representable, and regular matroids are all closed under 2-sums.

The following proposition describes the set of circuits in the 2-sum.

Proposition 1.8.12. [6] The 2-sum of M_1 and M_2 is the matroid whose ground set is $(E_1 \cup E_2) - \{p\}$ and whose set of circuits consists of all circuits of $M_1 \setminus p$ together with all circuits of $M_2 \setminus p$ and all sets of the form $(C_1 \cup C_2) - \{p\}$ where each C_i is a circuit of M_i containing p.

The next figure gives a general illustration of the three different ways 2-separation occurs.

Figure 1.4(a) illustrates a 1-separation in M. It has a partition (A, B) such that r(A) + r(B) = r(M). Figure 1.4(b), (c) and (d) illustrate the three ways 2-separation can occur. Part (a) shows a parallel connection, (b) shows a series connection and (c) shows a 2-separation where there is no base-point. Figure 1.4 only gives the general view. The figure in the next example shows concrete instances of 2-separation.

Example 1.8.13. Figure 2.1(a) shows a 1-separation of matroid M. The two disjoint sets are $A = \{x\}, B = \{a, b\}$, so that $M = (M|A) \oplus (M|B)$. Figure 2.1(b) shows a 2-separation as a parallel connection. In (b) $A = \{p, x, y\}$ and $B = \{p, b, c\}$. So that M = P(M|A, M|B). Figure 2.1(c) shows a 2-separation as a series connection. In (c) $A = \{p, x, y, a\}$ and $B = \{p, b, c, d, e\}$. So that M = S(M|A, M|B). Finally Figure 2.1(d) shows a 2-separation where there is no base-point p. Here $A = \{p, a, x, y\}$, and $B = \{p, d, b, c\}$. So that $M = (M|A) \oplus_2 (M|B)$.

F-representable matroids are also closed under the operation of series and parallel connections, as shown in the next result.

Proposition 1.8.14. [5, Proposition 7.1.21] Let F be a field. If M_1 and M_2 are F-representable matroids such that $E(M_1) \cap E(M_2) = \{p\}$, then both $P(M_1, M_2)$ and $S(M_1, M_2)$ are F-representable.

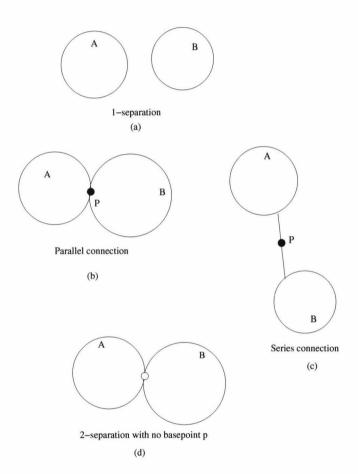


Figure 1.4: General view of 1-separation and 2-separations.

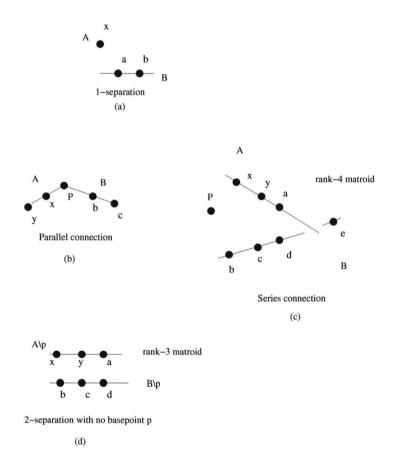


Figure 1.5: Examples of 1- and 2-separations.

The following proposition ties together the operations of 2-sum, parallel and series connections.

Lemma 1.8.15. Let M be the 2-sum of the two matroids M_A and M_B where $\{p\} = E(M_A) \cap E(M_B)$ is the base-point. Then the following are equivalent:

- (i) $M = M_A \oplus_2 M_B$.
- (ii) $M = P(M_A, M_B) \setminus p$.
- (iii) $M = S(M_A, M_B)/p$.

The operation of 2-sum and 3-connected matroids are linked as will soon be shown in Theorem 1.8.19, but first we state an important fact about connectivity. The next proposition says that given a matroid M, if you delete or contract an element from M, then the connectivity of M drops by at most one.

Proposition 1.8.16. [5, Proposition 8.1.13] If e is an element of an n-connected matroid M, then, provided $|E(M)| \geq 2(n-1)$, both $M \setminus e$ and M/e are (n-1)-connected.

The connectivity of a matroid M, is determined by its connectivity function r(X) + r(E(M) - X) - r(M) defined on all subsets X of E(M).

Lemma 1.8.17. [5, Lemma 8.1.14] A matroid and its dual have the same connectivity function.

The following lemma shows that if the condition $|E(M)| \ge 2(n-1)$ does not hold in Proposition 1.8.16, then the proposition fails.

Lemma 1.8.18. [2, Section 5] Let e be an element of a 3-connected matroid M, where $M \cong U_{2,3}$. Then $M \setminus e$ is disconnected.

The following theorem shows the link between 3-connected matroids and 2-sums.

Theorem 1.8.19. [5, Theorem 8.3.1] A 2-connected matroid M is not 3-connected if and only if $M = M_1 \oplus_2 M_2$ for some matroids M_1 and M_2 , each of which is isomorphic to a proper minor of M.

The two following results show how matroids that are not 3-connected can be constructed.

Theorem 1.8.20. [5, Theorem 8.3.4] Every matroid that is not 3-connected can be constructed from 3-connected proper minors of it by a sequence of the operations of direct sum and 2-sum.

Proposition 1.8.21. [5, Proposition 8.3.5] Let M be a connected matroid. Then, for some positive integer k, there is a collection $M_1, M_2, ..., M_k$ of 3-connected matroids and a k-vertex tree T with edges labeled $e_1, e_2, ..., e_{k-1}$ and vertices labeled $M_1, M_2, ..., M_k$ such that

- (i) $E(M_1) \cup E(M_2) \cup ... \cup E(M_k) = E(M) \cup \{e_1, e_2, ..., e_k\};$
- (ii) if the edge e_i joins the vertices M_{j1} and M_{j2} , then $E(M_{j1}) \cap E(M_{j2})$ is $\{e_i\}$; and
- (iii) if no edge joins the vertices M_{j1} and M_{j2} , then $E(M_{j1}) \cap E(M_{j2})$ is empty.

Moreover, M is the matroid that labels the single vertex of the tree $T/e_1, e_2, ..., e_{k-1}$ at the conclusion of the following process: contract the edges $e_1, e_2, ..., e_{k-1}$ of T one by one in order; when e_i is contracted, its ends are identified and the vertex formed by this identification is labeled by the 2-sum of the matroids that previously labeled the ends of e_i .

Chapter 2

Fundamentals of matroid union

Given two matroids M_1 and M_2 , we can consider a way of joining the two matroids, this operation is called 'matroid union'. Here we define it for two matroids, but that can be easily extended for any finite number of matroids.

2.1 Matroid union

Let M_1 and M_2 be two matroids, and let $\mathcal{I} = \{I_1 \cup I_2 : I_1 \in \mathcal{I}(M_1), I_2 \in \mathcal{I}(M_2)\}$. Then the following is true.

Theorem 2.1.1. \mathcal{I} is the collection of independent sets of a matroid on $E(M_1) \cup E(M_2)$.

Proof. Consider the set \mathcal{I} . Since \emptyset is in both $\mathcal{I}(M_1)$ and $\mathcal{I}(M_2)$ by definition of matroids, and since $\emptyset \cup \emptyset = \emptyset$, then we have $\emptyset \in \mathcal{I}$, and so (I1) is satisfied. Now if $I \in \mathcal{I}$ and $I' \subseteq I$, then we have $I = I_1 \cup I_2$ where $I_1 \in \mathcal{I}(M_1)$ and $I_2 \in \mathcal{I}(M_2)$. If either $I' \subseteq I_1$, or $I' \subseteq I_2$, then $I' \in \mathcal{I}(M_1)$, or $I' \in \mathcal{I}(M_2)$ respectively. It follows that $I' \in \mathcal{I}$. Or else $I' \subseteq (I_1 \cup I_2)$, without being a subset of either I_1 or I_2 . But in this case I' itself is the union of two independent sets, one being the set $I' \cap I_1 = I'_1$, and the other being $I' \cap I_2 = I'_2$. That gives $I'_1 \subseteq I_1$, implying $I'_1 \in \mathcal{I}(M_1)$, and similarly $I'_2 \subseteq I_2$, implying $I'_2 \in \mathcal{I}(M_2)$. Hence $I' \in \mathcal{I}$, and (I2) is satisfied. If X and Y are in

 \mathcal{I} with |X| < |Y|, then $X = X_1 \cup X_2$, where X_1 in $\mathcal{I}(M_1)$, and X_2 in $\mathcal{I}(M_2)$, and similarly $Y = Y_1 \cup Y_2$, where Y_1 in $\mathcal{I}(M_1)$, and Y_2 in $\mathcal{I}(M_2)$. So we have $|X_1 \cup X_2| < |Y_1 \cup Y_2|$. Without loss of generality suppose $|X_1| < |Y_1|$. Then there exists an element e in $Y_1 - X_1$ such that $(X_1 \cup \{e\}) \in \mathcal{I}(M_1)$. Since X_2 is already in $\mathcal{I}(M_2)$, then $(X_1 \cup X_2 \cup \{e\}) \in \mathcal{I}$. Now that we have (I1), (I2) and (I3) satisfied, we conclude that \mathcal{I} is the collection of independent sets of a matroid on $E(M_1) \cup E(M_2)$.

The matroid constructed via Theorem 2.1.1 is the union of M_1 and M_2 , and is denoted by $M_1 \vee M_2$. It follows that we can also talk about the operation of decomposition of matroids, where a matroid is written as the union of two matroids. Note that any matroid can be written as the union of itself with a rank-zero matroid. However, since this is a trivial operation, we will not be concerned with it. What we are concerned with is a proper decomposition. A matroid M is decomposable if there exist matroids M_1 and M_2 such that $M = M_1 \vee M_2$ where $r(M_1) \geq 1$ and $r(M_2) \geq 1$. Otherwise a matroid is indecomposable. Some authors use the words reducible for decomposable, and irreducible for indecomposable. In this thesis, however, we will always use the words decomposable and indecomposable. The following figure gives an example of a matroid union.

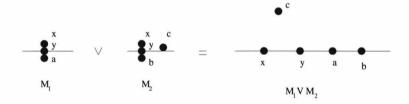


Figure 2.1: An example of a matroid union.

2.2 Ground sets of the union

If $M = M_1 \vee M_2$, then it follows from the definition of matroid union that M_1 and M_2 may have different ground sets. However, at times it is convenient to have $E(M_1) = E(M_2)$. The next proposition shows that we can always do that.

Proposition 2.2.1. If a matroid M is decomposable such that $M = M_1 \vee M_2$, then there are matroids M_1^+, M_2^+ on a common ground set E such that $M = M_1^+ \vee M_2^+$, where $r(M_1^+), r(M_2^+) \geq 1$.

Proof. If M_1 and M_2 above have the same ground set E then we are done. So let us assume that M_1 and M_2 do not have a common ground set and that M_1 is defined over E_1 and M_2 defined over E_2 , and let $E = E_1 \cup E_2$. Define M_1^+ to be the following extension of M_1 . Take E_1 and adjoin to it a set of loops where these loops are labeled by the elements of $E - E_1$. The same is done for M_2^+ . Thus $M_i^+|E_i=M_i$ and $M_i^+|(E-E_i)$ has rank zero, for $i \in \{1,2\}$. Note that now M_1^+ and M_2^+ have a common ground set E. Since the loops in M_1^+ cannot be members of any of the independent sets of M_1^+ , then the set of independent sets of M_1^+ is the same as the set of independent sets of M_1 . Similarly for M_2^+ . So we have $\mathcal{I}(M_1^+ \vee M_2^+) = \mathcal{I}(M_1 \vee M_2)$. Therefore $M_1^+ \vee M_2^+ = M_1 \vee M_2$, with the only difference being that now M_1^+ and M_2^+ have common ground sets. So $E(M_1^+) = E(M_2^+) = E(M_1^+ \vee M_2^+) = E(M)$.

The disadvantage of assuming that we have the same ground sets is that M_1 and M_2 may have loops that play no structural role. At times it is convenient to assume that M_1 and M_2 are loopless. The next proposition shows that we can always do that.

Proposition 2.2.2. If M is a loopless decomposable matroid, then there exist loopless matroids M_1 and M_2 , each having rank at least one, such that $M = M_1 \vee M_2$.

Proof. Since M is decomposable, then suppose $M = N_1 \vee N_2$. Suppose that N_1 and N_2 are not loopless. Consider $e \in E(M)$. If $e \in E(N_1) \cap E(N_2)$ and e is a loop in N_1 and a loop in N_2 , then e is a loop in M contradicting the fact that M is loopless. So let A be a set of loops of N_1 without loss of generality. A is a set of loops of N_1 only by the above. Then $\mathcal{I}(N_1) = \mathcal{I}(N_1 \setminus A)$, because A is a set of loops, and therefore none of its elements are contained in any independent set of N_1 . Let $M_1 = N_1 \setminus A$. Similarly, let B be the set of loops of N_2 . Let $M_2 = N_2 \setminus B$. Then $\mathcal{I}(N_2) = \mathcal{I}(N_2 \setminus B)$. It follows that $M = M_1 \vee M_2$, where both M_1 and M_2 are loopless. \square

In this thesis we will at times assume that M_1 and M_2 have different ground sets, and at other times we will assume they have the same ground sets according to convenience. We will always be clear as to which case we are in.

2.3 Properties of matroid union

In this section we state some properties of matroid union. We also look at equivalent ways of describing matroid unions when the ground sets are the same. The two following propositions are elementary, they show that matroid union is commutative and associative. In these two propositions we make no assumption about the ground sets.

Proposition 2.3.1. Let M_1 and M_2 be two matroids, then $M_1 \vee M_2 = M_2 \vee M_1$.

Proposition 2.3.2. Let M_1, M_2 and M_3 be three matroids, then $(M_1 \lor M_2) \lor M_3 = M_1 \lor (M_2 \lor M_3)$.

Proof. We observe that a set I in $(M_1 \vee M_2) \vee M_3$ is independent if and only if it is the union of two independent sets I_{12} and I_3 , where I_{12} is independent in $M_1 \vee M_2$, and I_3 is independent in M_3 . But I_{12} itself must be the union of two independent sets I_1 and I_2 , where I_1 is independent in M_1 and I_2 independent in M_2 . Since the order in which we take the union of I_1 , I_2 and I_3 does not matter, then it follows that $(M_1 \vee M_2) \vee M_3$ and $M_1 \vee (M_2 \vee M_3)$ have exactly the same independent sets. Therefore they are equal.

The following theorem shows what the rank function of the union is when the ground sets are the same.

Theorem 2.3.3. [5, Theorem 12.3.1] Let M_1 and M_2 be two matroids on E. If $X \subseteq E$, its rank in $M_1 \vee M_2$ is

$$min\{r_1(Y) + r_2(Y) + |X - Y| : Y \subseteq X\}$$

where r_1 and r_2 are the rank functions of M_1 and M_2 respectively.

To prove Theorem 2.3.3, first we state a couple of results. Note that the notion of matching mentioned in the next theorem will be defined later in case the reader is not already familiar with it.

Theorem 2.3.4. [5, Theorem 12.2.12] Let \triangle be a bipartite graph with vertex classes S and J and let M be a matroid on J. Suppose that $\triangle(\mathcal{I})$ is the set of subsets X of S that are matched in \triangle onto a member of $\mathcal{I}(M)$. Then $\triangle(\mathcal{I})$ is the set of independent sets of a matroid $\triangle(M)$ on S.

Note that the symbol \sqcup used in the proof for the next proposition indicates a disjoint union.

Proposition 2.3.5. [5, Corollary 12.2.14] If $X \subseteq S$, its rank in $\triangle(M)$ is equal to $min\{r_M(N(Y)) + |X - Y| : Y \subseteq X\}$.

Proof of Theorem 2.3.3. Construct the bipartite graph \triangle as follows. Let ϕ_1 and ϕ_2 be bijections from E onto disjoint sets E_1 and E_2 . Then for i=1,2, the function ϕ_i induces an isomorphic copy of M_i on E_i . Hence on $E_1 \sqcup E_2$, we have a matroid isomorphic to $M_1 \oplus M_2$. Let \triangle have E as one vertex class and $E_1 \cup E_2$ as the other. Join each element e of E to $\phi_1(e)$ and $\phi_2(e)$ and to no other vertices in $E_1 \cup E_2$ as in Figure 2.2. Now consider the matroid N on E induced from $M_1 \oplus M_2$ by \triangle . Since $M_1 \oplus M_2$ has rank function $r_1 + r_2$, Proposition 2.3.5 implies that the rank function of N is as specified in the theorem. Moreover, one easily checks that $\mathcal{I}(N) = \{I_1 \cup I_2 : I_1 \in \mathcal{I}(M_1), I_2 \in \mathcal{I}(M_2)\}$.

When the ground sets are the same, then the collection of circuits in the union is as follows.

Proposition 2.3.6. [11, Proposition 7.6.12] Let M_1 and M_2 be two matroids on E. The collection of circuits C in $M_1 \vee M_2$ is as follows: C consists of minimal sets of the form $\{C: A \cap C \text{ contains a circuit of } M_1, \text{ or } C - A \text{ contains a circuit of } M_2 \text{ for all } A \subseteq C\}.$

The closure operator for the matroid union M, where the ground sets are the same is as follows.

Proposition 2.3.7. [11, Proposition 7.6.12] Let M_1 and M_2 be two matroids on E. For a set X in $M = M_1 \vee M_2$, $cl_{M_1 \vee M_2}(X) = X \cup \{x : x \in M_1 \cap M_2\}$

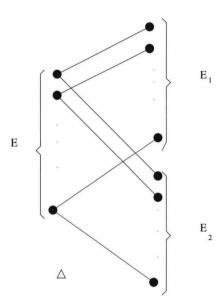


Figure 2.2: $M_1 \vee M_2$ is induced from $M_1 \oplus M_2$ by \triangle .

there exists $X_1, X_2 \subseteq X$, with $X_1 \cap X_2 = \emptyset$, such that $x \in cl_{M_1}(X_1)$ and $x \in cl_{M_2}(X_2), X_1 \subseteq cl_{M_2}(X_2)$, and $X_2 \subseteq cl_{M_1}(X_1)$.

The bases in the union, when the ground sets are the same are as follows:

Proposition 2.3.8. [11, Proposition 7.6.12] Let M_1 and M_2 be two matroids on E. A basis B in the union is a maximal set of the form $\{B_1 \cup B_2 : where B_1 \text{ is a basis of } M_1 \text{ and } B_2 \text{ is a basis of } M_2\}.$

Proof. Let B be a basis of $M_1 \vee M_2$. B is a maximal independent set. So it follows from the definition of matroid union, that B is the union of two independent sets I_1 and I_2 , where I_1 is independent in M_1 and I_2 independent in M_2 . But in order for B to be a maximal independent set, at least one of I_1 and I_2 has to be a maximal independent set in M_1 or M_2 respectively, say I_1 . The independent set I_2 , however, can always be chosen to be a maximal independent set, even if that requires having common elements between I_1 and I_2 . So that they are bases of M_1 and M_2 respectively.

Having introduced the operation of matroid union, and some properties of the operation, we now state a well known unsolved problem in Matroid Theory. **Problem 2.3.9.** [5, Problem 12.3.9] Characterize all indecomposable matroids.

The problem mentioned above has been solved for some classes of matroids. Some examples of such classes are the classes of graphic, transversal, and binary matroids. The following is a characterization of graphic matroids M which are indecomposable, meaning that the only decomposition possible is the trivial decomposition in which a matroid is the union of itself with the empty matroid.

Theorem 2.3.10. [8] The cycle matroid M(G) of a graph G is indecomposable if and only if the removal of any edge leaves the graph 2-connected.

It is known that if M(G) is the union of two matroids then it is the union of two graphic matroids as well, see [8].

Next we look at a characterization of the decomposition of transversal matroids. If \mathcal{A} is the family (A_1, A_2, \ldots, A_m) of subsets of a set S and $J = \{1, 2, \ldots, m\}$, then the bipartite graph $\Delta[\mathcal{A}]$ associated with \mathcal{A} has vertex set $S \cup J$; its edge set is $\{xj : x \in S, j \in J, \text{ and } x \in A_j\}$. A matching in a graph is a set of edges in the graph no two of which have a common endpoint. A subset X of S is a partial transversal of \mathcal{A} if and only if there is a matching in $\Delta[\mathcal{A}]$ in which every edge has one endpoint in X. Note the following.

Theorem 2.3.11. [5, Theorem 1.6.2] Let A be a family $(A_1, A_2, ..., A_m)$ of subsets of a set S. Let I be the set of partial transversals of A. Then I is the collection of independent sets of a matroid on S.

The matroid obtained above from the set of partial transversals of \mathcal{A} is denoted by $M[\mathcal{A}]$. If M is an arbitrary matroid and $M \cong M[\mathcal{A}]$ for some family \mathcal{A} of sets, then M is a transversal matroid.

In other words, if M is a transversal matroid, then there is a bipartite graph G with vertex sets V and E(M), such that $I \subseteq E(M)$ is independent in M if and only if I can be matched in G. In this case I is called a partial transversal of V. For I to be matched means that every element x in I is connected to exactly one unique point in V.

The following construction will be used in the next proposition. Say $V = \{1, 2, ..., n\}$. For i in V, define M_i on E(M) as follows: if e is in E(M) and there is an edge from e to i then $\{e\}$ is independent in M_i , otherwise e is a loop in M_i . If $\{e, f\} \subseteq E(M)$ and e and f are both connected to i in V, then $\{e, f\}$ is dependent in M_i . M_i in this case is a rank-1 matroid. The following is the characterization of the decomposition of transversal matroids.

Proposition 2.3.12. [5, Proposition 12.3.7] A matroid M with rank at least one is transversal if and only if it is a union of rank-1 matroids.

Proof. Since we are working with a transversal matroid M, then we can construct a bipartite graph G as follows. It has vertex sets V and E(M), where V is the set of rank-1 matroids M_1, M_2, \ldots, M_n , and E(M) is the set of elements of M. If I is independent in the transversal matroid M, then I can be matched in G. Meaning that every element in I has exactly one endpoint in V. Since each of those single element sets have exactly one endpoint in V, then I meets the requirement of being independent in M_i . Implying that I can be partitioned into single element sets, each of which is a match in G. Therefore each element of I is independent in some rank-1 matroid M_i . But I is the union of those single element independent sets. So I is independent in $M_1 \vee M_2 \vee \cdots \vee M_n$. Therefore if I is independent in M then I is independent in $M_1 \vee M_2 \vee \cdots \vee M_n$.

Conversely, if I is independent in $M_1 \vee M_2 \vee \cdots \vee M_n$, then I is the union of n number of independent sets, each having exactly one element. Meaning each one-element set is a match in G. But because each one-element set of I is a match in G, then I itself is a match in G. So it follows that I is a transversal of $M_1 \vee M_2 \vee \cdots \vee M_n = M$. But if I is a transversal of M, then I is independent in M by definition. Therefore if I is independent in $M_1 \vee M_2 \vee \cdots \vee M_n$, then I is independent in M. The two directions show that a matroid M is transversal if and only if it is the union of rank-1 matroid.

The class of uniform matroids is also another class of decomposable matroids. But first we note the following.

Proposition 2.3.13. All uniform matroids are transversal.

Proof. Let \mathcal{A} be a family of independent sets of M, where M is a uniform matroid. Consider the general form of uniform matroids $U_{m,n}$. All the independent sets in M have size less than or equal to m. Since $m \leq n$, then we can construct a bipartite graph $\Delta[\mathcal{A}]$ as follows. It has vertex sets N of n elements and M of m elements. Each edge represents an element in a set, the element is the endpoint of the edge in M, and so a certain collection of edges represents a certain set. It follows each and every independent set of M is a matching in $\Delta[\mathcal{A}]$, because $m \leq n$. Meaning that the set \mathcal{I} of independent sets of M are partial transversals of \mathcal{A} . Therefore by Theorem 2.3.11, M is a transversal matroid.

By combining Propositions 2.3.12 and 2.3.13 we get:

Corollary 2.3.14. All uniform matroids of rank at least 2 are decomposable.

Another class for which there exists a characterization of decomposability is the class of binary matroids. It has been solved by Cunningham. In Chapter 4 we will give an alternative proof of his result. Here we state Cunningham's original result.

Theorem 2.3.15. [1] A binary matroid M on E is indecomposable if and only if M is connected and $M \setminus e$ is connected for each $e \in E$.

As a corollary of Theorem 2.3.15 we get the following.

Corollary 2.3.16. With the exception of $U_{2,3}$, every 3-connected binary matroid is indecomposable.

Proof. A binary matroid is decomposable by Theorem 2.3.15 if and only if it contains an element whose deletion disconnects the matroid. But by Proposition 1.8.16 a 3-connected matroid cannot be disconnected by deleting one element. The only exception to this case, however, is $U_{2,3}$ as stated in [2, Section 5]. Therefore with the exception of $U_{2,3}$, 3-connected binary matroids are indecomposable.

2.4 Minors of a decomposable matroid

In this section we show which minors of a decomposable matroid are also decomposable. First we look at minors obtained by deletions. In the special

case when the ground sets of M_1 and M_2 are the same, the following is true.

Proposition 2.4.1. Let M be a decomposable matroid so that $M = M_1 \vee M_2$, where $E(M_1) = E(M_2)$. If $X \subseteq E(M)$, then

- (i) $M \setminus X = (M_1 \setminus X) \vee (M_2 \setminus X)$.
- (ii) If neither $M_1 \setminus X$ nor $M_2 \setminus X$ consists entirely of loops, then $M \setminus X$ is decomposable.

Proof. (i) If I is independent in $M \setminus X$, then I is independent in M, so that $I = I_1 \cup I_2$, where I_1 is independent in M_1 , and I_2 is independent in M_2 . As $I \cap X = \emptyset$, then I_1 is independent in $M_1 \setminus X$, and I_2 is independent in $M_2 \setminus X$. Therefore I is independent in $(M_1 \setminus X) \vee (M_2 \setminus X)$. Conversely, if I is independent in $(M_1 \setminus X) \vee (M_2 \setminus X)$, then $I = I_1 \cup I_2$, where I_1 is independent in $M_1 \setminus X$, and I_2 is independent in $M_2 \setminus X$. It follows that I_1 is independent in M_1 , and I_2 is independent in M_2 . Meaning I is independent in M. But $I \cap X = \emptyset$. Therefore I is independent in $M \setminus X$. (ii) By (i) $M \setminus X = (M_1 \setminus X) \vee (M_2 \setminus X)$. If neither $M_1 \setminus X$ nor $M_2 \setminus X$ consists entirely of loops, then $r(M_1 \setminus X) > 0$ and $r(M_2 \setminus X) > 0$. So by definition of decomposable matroids in Section 2, $M \setminus X$ is decomposable.

More generally, the following is true.

Proposition 2.4.2. Let $M = M_1 \vee M_2$, and suppose $x \in E(M)$. Then (i) $M \setminus x = (M_1 \setminus x) \vee (M_2 \setminus x)$ if $x \in E(M_1) \cap E(M_2)$. (ii) $M \setminus x = (M_1 \setminus x) \vee M_2$ if $x \in E(M_1) - E(M_2)$.

Proof. Proof of (i) is the same as that of Proposition 2.4.1(i). Consider (ii). If I is independent in $M \setminus x$, then I is independent in M, so that $I = I_1 \cup I_2$, where I_1 is independent in M_1 , and I_2 is independent in M_2 . As $I \cap x = \emptyset$, then I_1 is independent in $M_1 \setminus x$, and I_2 is independent in M_2 . Therefore I is independent in $(M_1 \setminus x) \vee M_2$. Conversely, if I is independent in $(M_1 \setminus x) \vee M_2$, then $I = I_1 \cup I_2$, where I_1 is independent in I_2 is independent in I_3 . Meaning I_4 is independent in I_4 and I_5 is independent in I_4 . So that I_4 is independent in I_4 . But $I_4 \cap I_4 \cap I_4$ is independent in I_4 . Therefore I_4 is independent in I_4 .

The following are two immediate corollaries.

Corollary 2.4.3. If $M = M_1 \vee M_2$, where $D \subseteq E(M)$, then $M|D = M_1|(D \cap E(M_1)) \vee M_2|(D \cap E(M_2))$.

Corollary 2.4.4. If $M = M_1 \vee M_2$, where $D \subseteq E(M_1) - E(M_2)$, then $M_1|D = (M_1 \vee M_2)|D$. In particular, if $I \subseteq D$, then I is independent in $M_1 \vee M_2$ if and only if I is independent in M_1 .

Second, we look at minors obtained by contractions.

Proposition 2.4.5. Let M_1 and M_2 be two matroids such that $x \in E(M_1) - E(M_2)$. Then

- (i) $(M_1 \vee M_2)/x = (M_1/x) \vee M_2$.
- (ii) If M_1/x does not consist entirely of loops, then M/x is decomposable.

Proof. (i) Consider $(M_1 \vee M_2)/x$. If x is a loop, then $(M_1 \vee M_2)/x =$ $(M_1 \vee M_2) \setminus x$ and it follows from Proposition 2.4.2(ii) that $(M_1 \vee M_2) \setminus x =$ $(M_1 \setminus x) \vee M_2 = (M_1/x) \vee M_2$. So assume x is not a loop. If I is independent in $(M_1 \vee M_2)/x$, then $I \cup \{x\}$ is independent in $M_1 \vee M_2$ by Proposition 1.4.1. Hence $I \cup \{x\} = I_1 \cup I_2$, where I_1 is independent in M_1 and I_2 is independent in M_2 . As $x \notin E(M_2)$, $x \in I_1$, so that $I_1 = I'_1 \cup \{x\}$. Now I'_1 is independent in M_1/x . Therefore $I = I'_1 \cup I_2$ where I'_1 is independent in M_1/x and I_2 independent in M_2 . Therefore I is independent in $(M_1/x) \vee M_2$. Conversely, if I is independent in $(M_1/x) \vee M_2$, then $I = I_1 \cup I_2$, where I_1 is independent in M_1/x and I_2 independent in M_2 . So that $I'_1 = I_1 \cup \{x\}$ is independent in M_1 . It follows that $I \cup \{x\} = I'_1 \cup I_2$ is independent in $M_1 \vee M_2$. Therefore I is independent in $(M_1 \vee M_2)/x$ by Proposition 1.4.1. (ii) By (i) $(M_1 \vee M_2)/x = (M_1/x) \vee M_2$. If M_1/x does not consist entirely of loops, then $r(M_1/x) > 0$. Therefore by the definition of decomposable matroids in Section 2, M/x is decomposable.

We shall give two examples, the first one shows how $(M_1 \vee M_2)/x$ is not equal to $(M_1/x) \vee (M_2/x)$ when x is in both M_1 and M_2 . The second example shows that a decomposable matroid may become indecomposable by contracting an element from $E(M_1) \cap E(M_2)$.

Example 2.4.6. Consider the matroid union shown in Figure 2.3. Here $E(M_1) = \{x, y, a\}$, $E(M_2) = \{x, y, b\}$ and $E(M) = \{x, y, a, b\}$. $M_1, M_2 \cong$

 $U_{1,3}$ and $M \cong U_{2,4}$ as shown in (a). We see that $M_1/x \cong U_{0,2}$, and $M_2/x \cong U_{0,2}$. Therefore $(M_1/x) \vee (M_2/x) \cong U_{0,3}$. Whereas $M/x \cong U_{1,3}$. Therefore $M/x \neq (M_1/x) \vee (M_2/x)$ as shown in (b).

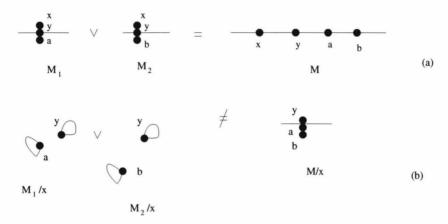


Figure 2.3: Contracting an intersection element.

Example 2.4.7. This example shows that if $x \in E(M_1) \cap E(M_2)$, then $(M_1 \vee M_2)/x$ is indecomposable. Consider the matroid $M = M_1 \oplus_2 M_2$ shown in Figure 2.4. Here $E(M_1) = \{a, b, c, d, x, r, p\}$, $E(M_2) = \{f, g, h, i, q, e, p\}$, where $M_1, M_2 \cong F_7$. By Lemma 1.8.15(iii) we have $M = S(M_1, M_2)/p$. Proposition 3.2.1 in the next section gives $M_1 \vee M_2 = S(M_1, M_2)$. So $M = (M_1 \vee M_2)/p$. But by combining Theorem 3.3.1 with Lemma 4.2.3 that we shall encounter later, we get that M is indecomposable because both M_1 and M_2 are indecomposable.

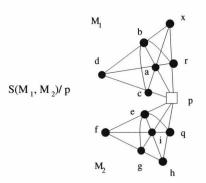


Figure 2.4: Contracting an intersection element may give an indecomposable matroid.

Chapter 3

Matroid union and connectivity

A disconnected matroid M is simply the direct sum of two disjoint matroids by Proposition 1.6.1. That of course is the same as saying M is the union of two disjoint matroids. So the case of decomposing disconnected matroids is trivial. Therefore we shall focus entirely on connected matroids from now on.

Sometimes the connectivity of a matroid can tell us whether the matroid is decomposable or not. Being decomposable does not necessarily mean having a low connectivity. In general it is possible to be decomposable and at the same time be highly connected. An example of this is $U_{100,200}$ which is highly connected but by corollary 2.3.14 is decomposable. However, if a matroid is the union of two matroids whose ground sets have few elements in common, then that can tell us something about the connectivity of the matroid M. In this chapter we explore the connection between matroid union and connectivity. First we look at the rank of the connected union. Second we look at the relationship between matroid union and series connection, then we look at the union that has a 2-separation, and finally we look at when the union is 3-connected.

3.1 The rank of a connected matroid union

The following is a Theorem of Cunningham [1, Theorem 2]. First we need a definition. Let M be a matroid and i be a positive integer not exceeding r(M). Then M' is a truncation of M denoted $T^i(M)$, if $\mathcal{I}(T^i(M)) = \{X \subseteq \mathcal{I}(M) : |X| \leq r(M) - i\}$. We state Cunningham's result below for two matroids only.

Theorem 3.1.1. If $M = M_1 \vee M_2$, then there exists a truncation M'_1 for M_1 , and M'_2 for M_2 , such that $r(M) = r(M'_1) + r(M'_2)$, and we have $M = M'_1 \vee M'_2$. Moreover, if M has no coloop, then the only choices for M'_1 and M'_2 are M_1 and M_2 respectively.

The two following results are corollaries of Theorem 3.1.1, the first one is immediate. The second is for the case when the matroid union is connected.

Corollary 3.1.2. If $M = M_1 \vee M_2$ is a matroid with no coloops, then $r(M) = r(M_1) + r(M_2)$.

Corollary 3.1.3. If $M = M_1 \vee M_2$ is a connected matroid with rank at least two, then $r(M) = r(M_1) + r(M_2)$.

Proof. By Theorem 3.1.1 $r(M) = r(M'_1) + r(M'_2)$, where M'_1 and M'_2 are truncations of M_1 and M_2 respectively. Since M is connected, then by Proposition 1.5.2 M has no coloops. So it follows by Theorem 3.1.1 that the only choices for M'_1 and M'_2 are M_1 and M_2 respectively. Therefore $r(M) = r(M_1) + r(M_2)$.

The following is a corollary about the bases in a connected matroid union.

Corollary 3.1.4. If $M = M_1 \vee M_2$ has no coloops, then B is a basis of M if and only if $B = B_1 \sqcup B_2$ where B_1 and B_2 are bases of M_1 and M_2 respectively.

Proof. If B is a basis of M, then by definition of matroid union we have $B = I_1 \sqcup I_2$, where I_1 is independent in M_1 and I_2 independent in M_2 . Since M has no coloops, then by Corollary 3.1.2 $r(M) = r(M_1) + r(M_2)$.

Let B_1 be a basis of M_1 , and B_2 be a basis of M_2 . Since r(M) = |B|, $r(M_1) = |B_1|$ and $r(M_2) = |B_2|$, then

$$|B| = |B_1| + |B_2|. (3.1.1)$$

But $|B| = |I_1 \cup I_2| = |I_1| + |I_2|$, and hence I_1 must be a basis of M_1 and I_2 a basis of M_2 . Conversely, if $B = B_1 \sqcup B_2$, where B_1 and B_2 are bases of M_1 and M_2 respectively, then $|B_1| = r(M_1)$, and $|B_2| = r(M_2)$. So

$$|B| = |B_1 \cup B_2| = |B_1| + |B_2| = r(M_1) + r(M_2). \tag{3.1.2}$$

But $r(M_1) + r(M_2) = r(M)$ by Corollary 3.1.2. So that r(M) = |B|, hence B is a basis of M.

3.2 Matroid union and series connection

In this section we look at the relationship between matroid union and series connection. We prove the following proposition, which is left as an exercise in [5].

Proposition 3.2.1. [5, Proposition 12.3.6] Let M_1 and M_2 be two matroids. If $E(M_1) \cap E(M_2) = \{p\}$, then $M_1 \vee M_2 = S(M_1, M_2)$.

First we observe the following.

Lemma 3.2.2. Let M_1 and M_2 be matroids and assume that z is a coloop of M_1 or M_2 . Then z is a coloop of $M_1 \vee M_2$.

Proof. Assume without loss of generality that z is a coloop of M_1 , and assume that B is a basis of $M_1 \vee M_2$. Then $B = B_1 \cup B_2$, where B_1 is a basis of M_1 , and B_2 is a basis of M_2 . If $z \notin B$, then $z \notin B_1$. But as z is a coloop of M_1 , it follows that z should be in every basis of M_1 , so we have a contradiction. Therefore z is in every basis of M meaning z is a coloop of M.

Proof of Proposition 3.2.1. There are two cases to consider.

Case 1. p is not a coloop in at least one of M_1 and M_2 . By Proposition 1.8.5, B is a basis of $S(M_1, M_2)$ if and only if $B = B_1 \sqcup B_2$, where B_1 is a

basis of M_1 and B_2 is a basis of M_2 . It follows from the definition of matroid union that B is a basis of $M_1 \vee M_2$. Therefore $S(M_1, M_2) = M_1 \vee M_2$.

Case 2. p is a coloop of both M_1 and M_2 . By Lemma 3.2.2, p is a coloop of $M_1 \vee M_2$. So that, without loss of generality, $M_1 \vee M_2 = (M_1 \setminus p) \oplus (M_2 \setminus p) \oplus (M$

$$M_1 \vee M_2 = (M_1 \setminus p) \oplus M_2. \tag{3.2.1}$$

By Proposition 1.8.2, 3.2.1 is equal to $S(M_1, M_2)$.

3.3 2-separation in the union

In this section we look at the connected matroid union that has a 2-separation. The following theorem is the main result of this section.

Theorem 3.3.1. Let (A, B) be a 2-separation of M such that $M = M_A \oplus_2 M_B$. If M is decomposable, then at least one of M_A and M_B is decomposable.

The next few results will be used in constructing a proof for the theorem above.

Proposition 3.3.2. Let M be a connected matroid, then the following is true:

- (i) If $M \setminus x$ is not connected, where x is in E(M), then M is decomposable.
- (ii) If M has a series pair then M is decomposable.

Proof. (i) If $M \setminus x$ is disconnected, then it is a series connection by Theorem 1.8.6(i), and then it follows by Proposition 3.2.1 that it is a union of two matroids. (ii) If M has a series pair $\{a, b\}$, then $M \setminus a$ and $M \setminus b$ have coloops, which implies they are both disconnected by Proposition 1.5.2. From part(i) now of the proposition it follows that M is decomposable.

Lemma 3.3.3. Let (A,B) be a 2-separation of a connected matroid M where $M = M_A \oplus_2 M_B$ and p is the base-point. If c is a coloop of M|A, then $\{c,p\}$ is a series pair of M_A , similarly, if c is a coloop of M|B, then $\{c,p\}$ is a series pair of M_B .

Proof. By a property of 2-sum, $M_A \setminus p = M|A$. It follows that c is a coloop of $M_A \setminus p$. So that A - c is a hyperplane of $M_A \setminus p$ by Proposition 1.3.2(iii). If $p \in cl_{M_A}(A-c)$, then c is a coloop of M_A , implying that M_A is disconnected by Proposition 1.5.2. But because M is connected, then M_A must also be connected by Proposition 1.8.10(ii), giving a contradiction. So $p \notin cl_{M_A}(A-c)$, meaning A-c is a hyperplane of M_A . Hence $\{p,c\}$ is a cocircuit of M_A , that is a series pair of M_A .

Lemma 3.3.4. If M|A has a coloop, then M_A is decomposable.

Proof. If M|A has a coloop, then by Lemma 3.3.3 above M_A has a series pair. But if M_A has a series pair, then it follows by Proposition 3.3.2(ii) that M_A is decomposable.

Finally before constructing a proof for Theorem 3.3.1, we need a definition. If a matroid M is obtained from a matroid N by deleting a non-empty subset T of E(N), then N is an extension of M. In particular, if |T| = 1, then N is a single-element extension of M. Two obvious ways to extend a matroid M are to adjoin a loop or to adjoin a coloop. In these cases the resulting matroids are isomorphic to $M \oplus U_{0,1}$ and $M \oplus U_{1,1}$, respectively. Another type of well-defined extension is parallel extension of M. Here we add a new element in parallel to some existing element of M. For more details see [5, Section 7.2].

Now we prove Theorem 3.3.1.

Proof of Theorem 3.3.1. Assume that M is decomposable. If M is not connected, then Proposition 1.8.10(ii) tells us that either M_A or M_B is disconnected and hence decomposable, and the theorem holds. So let us assume that M is connected. Let N_1 and N_2 be two matroids such that $M = N_1 \vee N_2$. We also assume that $E(N_1) = E(N_2) = E(M)$.

If M|A has a coloop, then by Lemma 3.3.4 above, M_A is decomposable and we are done. So assume that neither M|A nor M|B have coloops. Since M is connected, then $r(M) = r(N_1) + r(N_2)$ by Corollary 3.1.3. It also

follows by Corollary 3.1.2 that

$$r_M(A) = r_{N_1}(A) + r_{N_2}(A).$$
 (3.3.1)

Similarly,

$$r_M(B) = r_{N_1}(B) + r_{N_2}(B).$$
 (3.3.2)

Since (A, B) is a 2-separation of M, then

$$r(M) = r_M(A) + r_M(B) - 1$$

= $r_{N_1}(A) + r_{N_2}(A) + r_{N_1}(B) + r_{N_2}(B) - 1.$ (3.3.3)

The aforementioned tells us on one hand that either $r(N_1)$ equals $r_{N_1}(A) + r_{N_1}(B) - 1$, or it equals $r_{N_1}(A) + r_{N_1}(B)$. In the former (A, B) is a 2-separation of N_1 , in the later (A, B) is a 1-separation of N_1 . On the other hand, either $r(N_2)$ equals $r_{N_2}(A) + r_{N_2}(B) - 1$, or it equals $r_{N_2}(A) + r_{N_2}(B)$. In the former (A, B) is a 2-separation of N_2 , in the later (A, B) is a 1-separation of N_2 .

Without loss of generality we assume from now on that (A, B) is a 2-separation of N_1 and a 1-separation of N_2 .

Define a single element extension M' of M as follows. By Lemma 1.8.15 $M = P(M_A, M_B) \setminus p$. Let M' be $P(M_A, M_B)$. Then M' is a single element extension of M with ground set $E(M') = E(M) \cup \{p\}$, as shown in Figure 3.1.



Figure 3.1: $M' = P(M_A, M_B)$.

We now define two single element extensions N_1' and N_1' of N_1 and N_2 respectively. We have assumed that (A, B) is a 2-separation of N_1 , meaning N_1 is the 2-sum of two matroids. Let N_1 be the 2-sum of N_{1A} and N_{1B} . So by Lemma 1.8.15 $N_1 = P(N_{1A}, N_{1B}) \setminus p$. Let N_1' be $P(N_{1A}, N_{1B})$, as shown in Figure 3.2.



Figure 3.2: $N'_1 = P(N_{1A}, N_{1B})$.

Since (A, B) is a 1-separation of N_2 , then N_2 is the 1-sum of two matroids. So $N_2 = (N_2|A) \oplus (N_2|B)$. Here the single element extension of N_2 giving N_2' is as follows. Add $\{p\}$ as a loop to N_2 and get the new matroid $N_2' = (N_2|A) \oplus (N_2|B) \oplus U_{0,1}$, on the ground set $E(N_2') = E(N_2) \cup \{p\}$ as shown in Figure 3.3.

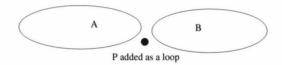


Figure 3.3: $N_2' = (N_2|A) \oplus (N_2|B) \oplus U_{0,1}$.

Now that we have extended M, N_1 , and N_2 to M', N'_1 and N'_2 respectively, we show that $M' = N'_1 \vee N'_2$. We will analyze the bases for both sides. We have two cases to consider, one where bases contain p and the other where they do not.

Case 1. Consider the bases of $M' = P(M_A, M_B)$ containing p. They are by Proposition 1.8.5(ii) of the form

$$B = B_A \cup B_B, \tag{3.3.4}$$

where $B_A = B \cap E(M_A)$ is a basis of M_A containing p, and $B_B = B \cap E(M_B)$ is a basis of M_B containing p. A basis of $N'_1 \vee N'_2$ containing p is of the form

$$B' \cup B'', \tag{3.3.5}$$

where B' is a basis of N'_1 , and B'' is a basis of N'_2 . Element p cannot be in B'' as p is a loop of N'_2 . Therefore $p \in B'$.

A basis of $N'_1 = P(N_{1A}, N_{1B})$ containing p is of the form

$$B' = B'_A \cup B'_B, (3.3.6)$$

where $B'_A = B' \cap E(N_{1A})$ is a basis of N_{1A} containing p, and $B'_B = B' \cap E(N_{1B})$ is a basis of N_{1B} containing p. A basis of N'_2 is of form

$$B'' = B_A'' \cup B_B'', \tag{3.3.7}$$

where $B_A'' = B'' \cap E(N_2|A)$ is a basis of $N_2|A$, and $B_B'' = B'' \cap E(N_2|B)$ is a basis of $N_2|B$.

A basis of $N_1' \vee N_2'$ is by 3.3.6 and 3.3.7 $B' \cup B'' = B_A' \cup B_B' \cup B_A'' \cup B_B'' = (B_A' \cup B_A'') \cup (B_B' \cup B_B'')$. Since $(N_{1A}) \cup (N_2|A) = M_A$, and $(N_{1B}) \cup (N_2|B) = M_B$, then $(B_A' \cup B_A'')$ is a basis of M_A , and $(B_B' \cup B_B'')$ is a basis of M_B . Therefore $B = B_A \cup B_B = B' \cup B''$, and so the bases of M' and $N_1' \vee N_2'$ are equal.

Case 2. Now consider the bases that do not contain p. Those bases of $M' = P(M_A, M_B)$ are by Proposition 1.8.5(ii) of the form

$$B = (B_A \cup B_B) - \{p\},\tag{3.3.8}$$

where $B_A = (B \cap E(M_A)) \cup \{p\}$ is a basis of M_A containing p, and $B_B = B \cap E(M_B)$ is a basis of M_B not containing p. Note that we lose no generality if p was in B_B and not in B_A .

Bases of $N'_{1} = P(N_{1A}, N_{1B})$ are

$$B' = (B'_A \cup B'_B) - \{p\},\tag{3.3.9}$$

where $B'_A = (B' \cap E(N_{1A})) \cup \{p\}$ is a basis of N_{1A} containing p, and $B'_B = (B' \cap E(N_{1B}))$ is a basis of N_{1B} not containing p. Again we lose no generality if p was in B'_B not B'_A .

As for the bases of N'_2 , they are of the form

$$B'' = B_A'' \cup B_B'' \tag{3.3.10}$$

not containing p, where $B''_A = (B'' \cap A)$ is a basis of $(N_2|A)$, and $B''_B = (B'' \cap B)$ is a basis of $(N_2|B)$. We see that a basis of $N'_1 \vee N'_2$ not containing p by 3.3.9 and 3.3.10 is $B' \cup B'' = (B'_A \cup B'_B \cup B''_A \cup B''_B) - \{p\} = (B'_A \cup B''_A) \cup (B'_B \cup B''_B) - \{p\}$. Which is equal to $(B_A \cup B_B) - \{p\}$ as discussed in Case 1. From Cases 1 and 2 we conclude that $M' = N'_1 \vee N'_2$.

If every element of E(M) is either a loop in N_1 or a loop in N_2 , then $N_1 \vee N_2$ will be equivalent to a direct sum and hence disconnected, giving a contradiction. So there must exist an element $z \in E(M)$, such that z is not a loop in N_1 and not a loop in N_2 . So without loss of generality we assume that $z \in A$. Since $M' = N'_1 \vee N'_2$, then $M' \setminus B = (N'_1 \setminus B) \vee (N'_2 \setminus B)$ by Proposition 2.4.1(i), where

$$M' \setminus B = M_A. \tag{3.3.11}$$

Element $z \in A$ and z is not a loop. Also $A \subseteq E(N'_1 \setminus B)$, and $A \subseteq E(N'_2 \setminus B)$ as shown in Figures 3.2 and 3.3. So it follows that neither $N'_1 \setminus B$ nor $N'_2 \setminus B$ consists entirely of loops. Therefore it follows from 3.3.11 and Proposition 2.4.1(ii) that M_A is decomposable. Therefore we conclude that if $M = M_A \oplus_2 M_B$ is decomposable, then at least one of M_A and M_B is decomposable.

The converse of Theorem 3.3.1 is not necessarily correct as we shall see shortly. The next theorem, however, is a weaker result.

Theorem 3.3.5. Let $M = M_A \oplus_2 M_B$, such that $\{p\} = E(M_A) \cap E(M_B)$. If one of M_A and M_B is decomposable, say M_A , such that $M_A = N_1 \vee N_2$, and $p \notin [E(N_1) \cap E(N_2)]$, then M is decomposable.

Proof. If M is disconnected, then it is decomposable. So assume that M is connected. By Proposition 1.8.10, both M_A and M_B are connected, therefore neither of them has a coloop. Suppose without loss of generality that $p \in E(N_2)$, and $p \notin E(N_1)$. By Lemma 1.8.15, $M = S(M_A, M_B)/p$. Since we have $M = (N_1 \vee N_2) \oplus_2 M_B$, then Proposition 3.2.1 gives

$$M = S((N_1 \vee N_2), M_B)/p = [(N_1 \vee N_2) \vee M_B]/p.$$
 (3.3.12)

That in turn by Proposition 2.3.2 is equal to

$$[N_1 \lor (N_2 \lor M_B)]/p.$$
 (3.3.13)

Since p is not in $E(N_1)$, then by Proposition 2.4.5, 3.3.13 is equal to

$$N_1 \vee [(N_2 \vee M_B)/p].$$
 (3.3.14)

We already know that N_1 and N_2 do not consist entirely of loops. Since N_2 and M_B have element p in common, then p cannot be in parallel with any element in $N_2 \vee M_B$, meaning contracting p from $N_2 \vee M_B$ can never make $(N_2 \vee M_B)/p$ consist entirely of loops. So $(N_2 \vee M_B)/p$ does not consist entirely of loops. Then it follows from Proposition 2.4.1(ii) that M is decomposable.

To see why the converse of Theorem 3.3.1 may not work, consider the following example.

Example 3.3.6. Let M_1, M_2 and M_3 be three matroids, such that $M_1 \cong M_2 \cong M_3 \cong F_7$, where element p is the only element they have in common. Consider $S(M_1, M_2) \oplus_2 M_3$ as shown in Figure 3.4. $S(M_1, M_2)$ is a matroid union by Proposition 3.2.1. Moreover, since the series connection of two binary matroids is also binary by Proposition 1.8.14, then $S(M_1, M_2)$ is binary. Also, since F-representable matroids are closed under the operation of 2-sum by Proposition 1.8.11, then $S(M_1, M_2) \oplus_2 M_3$ is binary too. Every element now in this 2-connected binary matroid $S(M_1, M_2) \oplus_2 M_3$ is fixed, a notion that will be explained in the next chapter. In Section 4.4.1, Theorem 4.4.3 shows why this structure gives an indecomposable matroid.

3.4 When is the union 3-connected?

In this section we characterize when the union of two matroids is 3-connected under the assumption that there are two elements in common in their ground sets. First we need a definition. Let x and y be elements of the connected matroid M. Then x and y are split if there is no 2-separation (A, B) in M such that $x, y \in A$ or $x, y \in B$. In other words, whenever you see a 2-separation in M, then x is never in the same part of the 2-separation as y. Note that, if M is 3-connected, then every two elements are split. The following theorem characterizes when the union is 3-connected, in the case

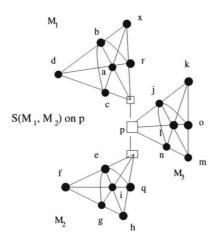


Figure 3.4: $S(M_1, M_2) \oplus_2 M_3$.

that $E(M_1)$ and $E(M_2)$ intersect in two elements, it is the main theorem of this section.

Theorem 3.4.1. Let M_1 and M_2 be two loopless matroids with $E(M_1) \cap E(M_2) = \{x, y\}$. $M_1 \vee M_2$ is 3-connected if and only if the following hold:

- (i) M_1 and M_2 are connected with no series pairs.
- (ii) $|E(M_1)|, |E(M_2)| \geq 3$ and
- (iii) x and y are split in M_1 and split in M_2 .

The example below illustrates Theorem 3.4.1, that the union of two matroids is 3-connected if and only if the two matroids satisfy conditions (i)-(iii) in Theorem 3.4.1.

Example 3.4.2. Let M_1 be the matroid on $E(M_1) = \{a, b, c, x, y, h, g\}$, and M_2 be a matroid on $E(M_2) = \{f, e, d, x, y, i, j\}$ as shown in Figure 3.5(a). It is seen that neither M_1 nor M_2 is disconnected, and neither of them has a series pair. Moreover, x and y are split in both M_1 and M_2 . As shown in Figure 3.5(b), $M_1 \vee M_2$ is the rank-6 whirl \mathcal{W}^6 , which is 3-connected.

To prove Theorem 3.4.1, we first establish some lemmas. The first is elementary.

Lemma 3.4.3. Let x be an element of the matroid M. If C^* contains a cocircuit of $M \setminus x$, then $C^* \cup x$ contains a cocircuit of M.

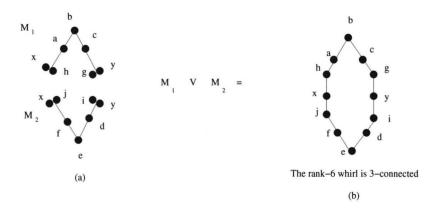


Figure 3.5: A 3-connected matroid union.

In the next lemma we make no assumptions about the cardinality of the intersections of the ground sets of M_1 and M_2 .

Lemma 3.4.4. Let M_1 and M_2 be matroids. If C^* is a cocircuit of M_1 or M_2 , then C^* contains a cocircuit of $M_1 \vee M_2$.

Proof. Assume C^* is a cocircuit of M_1 . If $|C^*| = 1$, then the lemma follows from Lemma 3.2.2. Thus we may assume that $|C^*| = k$ and, for induction, that the lemma holds for any cocircuit of M_1 with k-1 elements. Say $z \in C^*$, then $C^* - \{z\}$ is a cocircuit of $M_1 \setminus z$. If $z \in E(M_2)$, then $M \setminus z = (M_1 \setminus z) \vee (M_2 \setminus z)$ by Proposition 2.4.1, and if $z \notin E(M_2)$, then $M \setminus z = (M_1 \setminus z) \vee (M_2)$ by Proposition 2.4.2. In either case it follows from the induction assumption that $C^* - \{z\}$ contains a cocircuit of $M \setminus z$. The lemma now follows from Lemma 3.4.3.

Now we consider the different obstacles preventing $M = M_1 \vee M_2$ from being 3-connected.

Lemma 3.4.5. Let $M = M_1 \vee M_2$, where M_1 and M_2 are two matroids with $\{x,y\} = E(M_1) \cap E(M_2)$. If M_1 or M_2 is disconnected, then M is not 3-connected.

Proof. Without loss of generality, suppose that M_1 is disconnected. Suppose that M_1 has two components N_1 and N_2 , so that $M_1 = N_1 \oplus N_2 = N_1 \vee N_2$. There are two cases to consider.

Case 1. The elements x, y are in one component of M_1 , say N_1 without loss of generality. Note that N_1 and M_2 have x and y in common as shown in Figure 3.6, so let $N_1M_2 = N_1 \vee M_2$. It follows by Propositions 2.3.2 and 2.3.1 that

$$M = M_1 \lor M_2 = (N_2 \lor N_1) \lor M_2 = N_2 \lor (N_1 \lor M_2). \tag{3.4.1}$$

Sine N_2 and N_1M_2 are disjoint, then 3.4.1 is equal to $N_2 \oplus N_1M_2$. Therefore M is disconnected and hence not 3-connected.

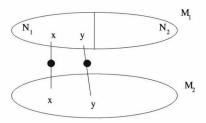


Figure 3.6: Intersection elements in one component.

Case 2. The element x is in one component of M_1 , say N_1 , and y is in the other component N_2 as shown in Figure 3.7. Let $M_{12} = N_1 \vee M_2$. Since $E(N_1) \cap E(M_2) = \{x\}$, then by Proposition 3.2.1 $M_{12} = S(N_1, M_2)$. By Propositions 2.3.2 and 2.3.1, we get

$$M = M_2 \lor (N_1 \lor N_2) = (M_2 \lor N_1) \lor N_2$$

= $S(M_2, N_1) \lor N_2 = M_{12} \lor N_2$. (3.4.2)

Since $E(M_{12}) \cap E(N_2) = \{y\}$, then $M_{12} \vee N_2 = S(M_{12}, N_2)$. But $M_{12} \vee N_2 = M$, so M is a series connection, showing M is not 3-connected.

Lemma 3.4.6. If $M = M_1 \vee M_2$ is such that at least one of M_1 and M_2 is isomorphic to $U_{2,3}$, and $\{x,y\} = E(M_1) \cap E(M_2)$, then M is not 3-connected.

Proof. Suppose without loss of generality that $M_1 \cong U_{2,3}$. Since $x \in E(M_1) \cap E(M_2)$, then by Proposition 2.4.1 $M \setminus x = (M_1 \setminus x) \vee (M_2 \setminus x)$. Now $M_1 \setminus x$ and $M_2 \setminus x$ have only the element y in common. Therefore by

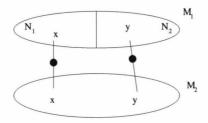


Figure 3.7: One intersection element in each component.

Proposition 3.2.1 $(M_1 \setminus x) \vee (M_2 \setminus x) = S(M_1 \setminus x, M_2 \setminus x)$. But $M_1 \setminus x$ is disconnected, and hence by Proposition 1.8.8, $M \setminus x = S(M_1 \setminus x, M_2 \setminus x)$ is also disconnected. It follows from Proposition 1.8.16 that M cannot be 3-connected.

The reason M was not 3-connected in Lemma 3.4.6 above, even though M_1 was 3-connected, is that $U_{2,3}$ has series pairs. So using Lemma 3.4.4 presented earlier, we can generalize Lemma 3.4.6 to the following:

Lemma 3.4.7. If $M = M_1 \vee M_2$ is such that $|E(M)| \geq 4$, M_1 or M_2 has a series pair, and $E(M_1) \cap E(M_2) = \{x, y\}$, then M is not 3-connected.

Proof. Suppose without loss of generality that M_1 has a series pair $C_1^* = \{s, t\}$. Since a series pair is a 2-element cocircuit, then by Lemma 3.4.4, C_1^* contains a cocircuit C^* in the matroid union M. But C_1^* only has two elements. So either C^* itself is a series pair in M, implying that M can be disconnected by removing one of the elements s and t from C^* . Hence by Proposition 1.8.16 M is not 3-connected, or C^* contains a 1-element cocircuit, that is a coloop. Showing that M is not 3-connected.

The results above are used in proving Theorem 3.4.1 part (i). The next Lemma is for proving part (ii).

Lemma 3.4.8. Let $M = M_1 \vee M_2$ is such that $|E(M)| \geq 4$, $E(M_1) \cap E(M_2) = \{x,y\}$, where M_1 and M_2 are two loopless matroids. If one of $E(M_1)$ and $E(M_2)$ has cardinality 2 or less, then M is not 3-connected.

Proof. Suppose without loss of generality that $|E(M_1)| \le 2$. $|E(M_1)|$ cannot be one or zero, because M_1 and M_2 must have exactly two elements in

common. So assume $|E(M_1)| = 2$, so that $E(M_1) = \{x,y\}$. Consider $(M_1 \vee M_2) \setminus x$. By Proposition 2.4.1 $(M_1 \vee M_2) \setminus x = (M_1 \setminus x) \vee (M_2 \setminus x)$. But now $M_1 \setminus x$ has only one element left in it which is y. Since M_1 is loopless, then y is not a loop in M_1 , so y is independent in M_1 , and hence it is in every basis of $M \setminus x$. Meaning y is a coloop of $M \setminus x$. Therefore by Proposition 1.5.2 $M \setminus x$ is disconnected. This shows by Proposition 1.8.16 that M cannot be 3-connected.

The next two Lemmas will be used to prove Theorem 3.4.1 Part (iii).

Lemma 3.4.9. Let $M = M_1 \vee M_2$ be a connected matroid, with $E(M_1) \cap E(M_2) = \{x, y\}$, where both M_1 and M_2 are connected, neither of them has a series pair, and $|E(M_1)|, |E(M_2)| \geq 3$. If at least one of M_1 and M_2 does not split x and y, then M is not 3-connected.

Proof. Suppose without loss of generality that M_1 does not split x and y, so that it has a partition (A, B) such that $x, y \in A$ as shown in Figure 3.8. Then

$$r(M_1) = r_{M_1}(A) + r_{M_1}(B) - 1. (3.4.3)$$

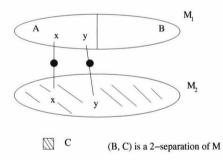


Figure 3.8: When is the union not 3-connected?

Say $E(M_2) = C$. So $C \cap A = \{x, y\}$. Now, as M is connected, then by Corollary 3.1.3 $r(M) = r(M_1) + r(M_2)$. But since $M_2 | C = M_2$, then by Corollary 2.4.3 $M | (A \cup C) = (M_1 | A) \vee (M_2 | C) = M_2 \vee (M_1 | A)$. By Lemma 1.2.1

$$r_M(A \cup C) \le r(M_2) + r_{M_1}(A).$$
 (3.4.4)

Also by Lemma 2.4.4

$$M|B = M_1|B. (3.4.5)$$

Hence by 3.4.3 we get

$$r(M) = r(M_1) + r(M_2) = r_{M_1}(A) + r_{M_1}(B) - 1 + r(M_2), (3.4.6)$$

and that by 3.4.4 gives

$$r_{M_1}(A) + r_{M_1}(B) - 1 + r(M_2) \ge r_M(A \cup C) + r_{M_1}(B) - 1$$

= $r_M(A \cup C) + r_M(B) - 1$. (3.4.7)

But by Lemma 1.2.1 $r(M) \leq r_M(A \cup C) + r_M(B)$. Therefore $(A \cup C, B)$ is either a 2-separation or 1-separation of M. As M is connected, then it must be that $(A \cup C, B)$ is a 2-separation of M. Hence M is not 3-connected.

For the converse of Theorem 3.4.1, we have the following:

Lemma 3.4.10. Let $M = M_1 \vee M_2$, where both M_1 and M_2 are connected, $\{x,y\} = E(M_1) \cap E(M_2)$, neither M_1 nor M_2 has a series pair, and $|E(M_1)|, |E(M_2)| \geq 3$. If neither M_1 nor M_2 splits x and y, then M is 3-connected.

The proof will use the following lemmas.

Lemma 3.4.11. Let N be a connected matroid with no series pairs. If $\{x,y\} \subseteq E(N)$, such that x and y are split in N, then $N \setminus x$ is connected.

Proof. Assume that $N \setminus x$ is disconnected. Let (A, B) be a separation of $N \setminus x$, where say $y \in A$. Then $(A \cup \{x\}, B)$ is a 2-separation of N unless |B| = 1. But if |B| = 1, then $B \cup \{x\}$ is a series pair in N, contradicting the fact that N has no series pairs. But now both x and y are in $A \cup \{x\}$, meaning x and y are not split in N, contradiction. Therefore $N \setminus x$ is connected.

Lemma 3.4.12. Let N be a connected matroid with no series pairs. Let (A, B) be a 2-separation of N, such that $z \in A$. Then z is a coloop of N|A

if and only if $N \setminus z$ is not connected, where $(A - \{z\}, B)$ is a separation of $N \setminus z$. Moreover, if the above holds, then $(A - \{z\}, B \cup \{z\})$ is a 2-separation of N.

Proof. Since (A, B) is a 2-separation of N, then r(A) + r(B) - 1 = r(N). If z is a coloop of N|A, then $r(A - \{z\}) = r(A) - 1$. So that $r(A - \{z\}) = r(N) - r(B) + 1 - 1 = r(N) - r(B)$. Therefore $N \setminus z$ is disconnected, where $(A - \{z\}, B)$ is a separation of $N \setminus z$. Conversely, if $N \setminus z$ is disconnected, then $(A - \{z\}, B)$ is a separation of $N \setminus z$. So that $r(A - \{z\}) + r(B) = r(N)$. But r(A) + r(B) - 1 = r(N). Meaning $r(A - \{z\}) = r(A) - 1$. Implying that z is a coloop of N|A. Moreover, if |A| = 2, then $(A - \{z\}, B) = r(A) - 1$. In a separation of $N \setminus z$, then z is in a series pair in N, contradicting the fact that N has no series pairs. □

Proof of Lemma 3.4.10. By Lemma 3.4.11, both $M_1 \setminus x$ and $M_2 \setminus x$ are connected. By Proposition 3.2.1 $(M_1 \setminus x) \vee (M_2 \setminus x) = S(M_1 \setminus x, M_2 \setminus x)$. So that by Proposition 1.8.8, $(M_1 \setminus x) \vee (M_2 \setminus x)$ is also connected. By Proposition 2.4.1 $M \setminus x = (M_1 \setminus x) \vee (M_2 \setminus x)$, so that $M \setminus x$ is connected. Similarly, $M \setminus y$ is connected. Since M_1 and M_2 are both connected, then neither M_1 nor M_2 have loops or coloops, therefore M does not have loops or coloops. Hence M is connected. But because $M \setminus x$ is also connected, then x is not in a series pair with any other element.

Now suppose that M has a 2-separation (A, B). Let $A_1 = A \cap E(M_1)$, $A_2 = A \cap E(M_2)$, $B_1 = B \cap E(M_1)$ and $B_2 = B \cap E(M_2)$. We have two cases to consider.

Case 1. $\{x,y\}\subseteq A$ or $\{x,y\}\subseteq B$. Without loss of generality, assume that $\{x,y\}\subseteq A$. By Lemmas 3.4.11 and 3.4.12, neither x nor y is a coloop of M|A. Hence we may assume that M|A has no coloops. So that by Corollary 3.1.2 we have

$$r(M|A) = r(M_1|A_1) + r(M_2|A_2). (3.4.8)$$

But as $\{x,y\} \cap B = \emptyset$, then

$$r(M|B) = r(M_1|B_1) + r(M_2|B_2). (3.4.9)$$

By 3.4.8 and 3.4.9 we have

$$r(M) = r(M|A) + r(M|B) - 1$$

= $r(M_1|A_1) + r(M_2|A_2) + r(M_1|B_1) + r(M_2|B_2) - 1$. (3.4.10)

But since M has no coloops, then by Corollary 3.1.2 $r(M) = r(M_1) + r(M_2)$. So that

$$r(M_1) + r(M_2) = r(M_1|A_1) + r(M_2|A_2) + r(M_1|B_1) + r(M_2|B_2) - 1.$$
 (3.4.11)

It may be that some of the subsets A_1, A_2, B_1 and B_2 are empty. Without loss of generality we consider the two following cases.

Case 1.1. $B_1 = \emptyset$. If $B_1 = \emptyset$, and A_1 is also empty, then $E(M_1) = A_1 \cup B_1$ would be empty, giving a contradiction. Therefore A_1 is not empty. Also, if B_2 is empty, then $B = B_1 \cup B_2$ is empty. But (A, B) is a 2-separation of M, meaning |B| must be at least 2, giving a contradiction. Therefore B_2 is not empty. Finally if A_2 is empty, then $r(M_1) + r(M_2) = r(M_1|A_1) + r(M_2|B_2) - 1$. But since B_1 is empty, then $M_1|A_1$ must be equal to M_1 . Similarly, $M_2|B_2$ must be equal to M_2 , giving a contradiction. Therefore if B_1 is empty, then it must be that $r(M_1) + r(M_2) = r(M_1|A_1) + r(M_2|A_2) + r(M_2|B_2) - 1$. Showing that (A_2, B_2) is a 2-separation of M_2 that does not split x and y, contradicting the assumption that x and y are split in M_2 .

Case 1.2. $B_1 \neq \emptyset$. If any of the other subsets are empty, then we get the same contradiction shown in Case 1.1. So suppose that none of the subsets are empty. Then either $r(M_1) = r(M_1|A_1) + r(M_1|B_1)$, showing that (A_1, B_1) is a separation of M_1 . Or $r(M_2) = r(M_2|A_2) + r(M_2|B_2)$, showing that (A_2, B_2) is a separation of M_2 . Both possibilities contradict the fact that both M_1 and M_2 are connected.

Case 2. Assume without loss of generality that $x \in A, y \in B$. Suppose that z is a coloop of M|A. Since neither x nor y are coloops of M|A, then $z \neq x$. As M has no series pairs, then $M \setminus z$ has no coloops. Assume that

 $z \in E(M_1)$. Then by Corollary 3.1.2 we have

$$r_{M\setminus z}(A - \{z\}) = r_{M_1}(A_1 - \{z\}) + r_{M_2}(A_2),$$
 (3.4.12)

and

$$r_{M\setminus z}(B) = r_{M_1}(B_1) + r_{M_2}(B_2).$$
 (3.4.13)

But

$$r(M \setminus z) = r_{M \setminus z}(A - \{z\}) + r_{M \setminus z}(B)$$

$$= r_{M_1}(A_1 - \{z\}) + r_{M_2}(A_2) + r_{M_1}(B_1) + r_{M_2}(B_2)$$

$$= r(M_1 \setminus z) + r(M_2). \quad (3.4.14)$$

This shows that (A_2, B_2) is a separation of M_2 , contradicting the fact that M_2 is connected. Thus M|A, and similarly M|B have no coloops. Therefore

$$r_M(A) = r_{M_1}(A_1) + r_{M_2}(A_2).$$
 (3.4.15)

$$r_M(B) = r_{M_1}(B_1) + r_{M_2}(B_2).$$
 (3.4.16)

So that

$$r(M) = r_M(A) + r_M(B) - 1 = r(M_1) + r(M_2)$$

= $r_{M_1}(A_1) + r_{M_2}(A_2) + r_{M_1}(B_1) + r_{M_2}(B_2) - 1$. (3.4.17)

But since $x \in A_1 \cap A_2$ and $x \in B_1 \cap B_2$, then none of the subsets A_1 , A_2 , B_1 or B_2 is empty. It follows that one of M_1 and M_2 has a separation, contradicting the fact that they are both connected. So we conclude that M has no 2-separation.

Now by combining all the aforementioned results we have a proof for Theorem 3.4.1.

Proof of Theorem 3.4.1. Let M_1 and M_2 be two loopless matroids with $E(M_1) \cap E(M_2) = \{x, y\}$. If $M_1 \vee M_2$ is 3-connected, then: (i) If at least one of M_1 and M_2 is disconnected, then by Lemma 3.4.5 M is not 3-connected, giving a contradiction. Therefore both M_1 and M_2 must be connected. If

at least one of M_1 and M_2 has a series pair, then by Lemma 3.4.7 M is not 3-connected, also giving a contradiction. Therefore neither M_1 nor M_2 can have a series pair. (ii) If at least one of $E(M_1)$ and $E(M_2)$ has size less than 3, then by Lemma 3.4.8 M is not 3-connected, giving a contradiction. So it must be that $|E(M_1)|, |E(M_2)| \geq 3$. (iii) If at least one of M_1 and M_2 does not split x and y, then Lemma 3.4.9 implies that M is not 3-connected, giving a contradiction. So x and y must be split in M_1 and split in M_2 .

Conversely, suppose both M_1 and M_2 are loopless connected matroids, neither of them has a series pair, x and y are split in M_1 and M_2 , and $|E(M_1)|, |E(M_2)| \geq 3$. If $M_1 \vee M_2$ was not 3-connected, then by Lemma 3.4.10 at least one of M_1 and M_2 does not split x and y, giving a contradiction. Hence M is 3-connected.

Chapter 4

Freedom of elements

4.1 Fixed elements and clones

Certain elements of a matroid have 'freedom', while others are 'fixed'. These properties of elements will be used a lot in the material to come, in this chapter and the next. So in this section we present the basic facts about fixed elements and clones. For more details see [3, Sections 2 and 4].

Elements x and x' of a matroid M are clones if the map that interchanges x and x' and acts as the identity on $E(M) - \{x, x'\}$ is an automorphism of M. In other words, x and x' are clones if they are indistinguishable up to labeling. An element z of M is fixed in M if there is no single-element extension of M by z' in which z and z' are independent clones. Dually, the element z is cofixed in M if there is no single-element coextension of M by z' in which z and z' are coindependent clones.

Let x be an element of a matroid M. The matroid M' is obtained by cloning x with x' if M' is a single-element extension of M by x', and x and x' are clones in M'.

If x is not fixed, then there is a matroid M' obtained by cloning x with x' such that $\{x, x'\}$ is independent in M'. We say that M' is obtained by independently cloning x with x'.

Proposition 4.1.1. [3, Section 4] If $\{x, x'\}$ is a pair of loops, a pair of coloops, a parallel pair, or a series pair, then x and x' are clones.

Next we state equivalent conditions for x and x' to be clones in a matroid

M, and then state further properties of fixed and non-fixed elements.

Proposition 4.1.2. [3, Proposition 4.2]. Let x and x' be elements of a matroid M. Then the following are equivalent.

- (i) x and x' are clones in M.
- (ii) The function that replaces x by x' and acts as the identity on other elements is an isomorphism from $M \setminus x$ to $M \setminus x'$.
 - (iii) $M/x \setminus x' = M/x' \setminus x$ and $r(\lbrace x \rbrace) = r(\lbrace x' \rbrace)$.

Proposition 4.1.3. [3, Corollary 4.4]. Let x be an element of a matroid M, and X be a subset of E(M) - x.

(i) If x is not fixed in M, then x is not fixed in $M \setminus X$.

Proposition 4.1.4. [3, Proposition 4.6]. Let x be an element of a matroid M.

(i) If M has an element a such that x is fixed in $M \setminus a$, then x is fixed in M.

Proposition 4.1.5. [3, Proposition 4.8]. If x and x' are independent clones in M, then x is fixed in neither M nor $M \setminus x'$.

4.2 Freedom of elements in the union

The aim of this section is to prove the following theorem. The theorem shows that when taking the union of two matroids, then the elements those two matroids have in common will not be fixed in the union.

Theorem 4.2.1. Let M_1 and M_2 be two matroids such that $M = M_1 \vee M_2$, and let $x \in E(M_1) \cap E(M_2)$. If x is not a loop in M_1 and not a loop in M_2 , then x is not fixed in M.

First we establish a lemma.

Lemma 4.2.2. Let M_1 and M_2 be two matroids with $\{x, x'\} \subseteq [E(M_1) \cap E(M_2)]$. If x and x' are clones in M_1 and clones in M_2 , then they are clones in $M_1 \vee M_2$.

Proof. By Proposition 4.1.2(ii) we need to show that replacing x by x' and fixing all other elements is an isomorphism between $(M_1 \vee M_2) \setminus x$ and $(M_1 \vee M_2) \backslash x'$. To do this it suffices to show that if $I \subseteq E(M_1 \vee M_2) - \{x, x'\}$, then $I \cup \{x'\}$ is independent in $(M_1 \vee M_2) \setminus x$ if and only if $I \cup \{x\}$ is independent in $(M_1 \vee M_2) \setminus x'$. By Proposition 2.4.1 $(M_1 \vee M_2) \setminus x =$ $(M_1 \setminus x) \vee (M_2 \setminus x)$ and $(M_1 \vee M_2) \setminus x' = (M_1 \setminus x') \vee (M_2 \setminus x')$. Let us suppose that $I \cup \{x'\}$ is independent in $(M_1 \vee M_2) \setminus x$, then $I \cup \{x'\}$ is independent in $(M_1 \setminus x) \vee (M_2 \setminus x)$. It follows by definition of matroid union that there is a partition (J, K) of $I \cup \{x'\}$, where J is independent in $M_1 \setminus x$ and K is independent in $M_2 \setminus x$. Without loss of generality, assume that $x' \in J$. So that $J = J' \cup \{x'\}$. Now since x and x' are clones in M_1 and clones in M_2 , then it follows by Proposition 4.1.2 that we can replace x by x' and have $J = J' \cup \{x\}$ is independent in $M_1 \setminus x'$ and K is independent in $M_2 \setminus x'$. Hence $(J' \cup \{x\}) \cup K = I \cup \{x\}$ is independent in $(M_1 \setminus x') \vee (M_2 \setminus x')$. Implying that $I \cup \{x\}$ is independent in $(M_1 \vee M_1) \setminus x'$. Similarly, if $I \cup \{x\}$ is independent in $(M_1 \vee M_1) \setminus x'$, then $I \cup \{x'\}$ is independent in $(M_1 \vee M_1) \setminus x$. It follows that x and x' are indeed clones in $M_1 \vee M_2$. \square

Proof of Theorem 4.2.1. Assume that x is not a loop in M_1 and not a loop in M_2 . Let M'_1 be the parallel extension of M_1 obtained by adding x' parallel to x, and similarly M'_2 be the parallel extension of M_2 obtained by adding x' parallel to x. Now since x and x' are parallel to each other in M'_1 , and parallel to each other in M'_2 , then by Proposition 4.1.1, x and x' are clones in M'_1 and clones in M'_2 . Therefore by Lemma 4.2.2, x and x' are clones in $M'_1 \vee M'_2$. Now $\{x, x'\}$ is independent in $M'_1 \vee M'_2$, because it is the union of $\{x\}$ and $\{x'\}$, where $\{x\}$ is independent in M'_1 and $\{x'\}$ is independent in M'_2 . Meaning $\{x\}$ and $\{x'\}$ are independent clones in $M'_1 \vee M'_2$. By Proposition 4.1.5 x is not fixed in $M'_1 \vee M'_2$. By Proposition 4.1.3(i), x is not fixed in $(M'_1 \vee M'_2) \setminus x'$. But by Proposition 2.4.1 $(M'_1 \vee M'_2) \setminus x' = (M'_1 \setminus x') \vee (M'_2 \setminus x') = M_1 \vee M_2$. Therefore x is not fixed in $M_1 \vee M_2$ as required.

As a consequence of Theorem 4.2.1, we get the following result.

Lemma 4.2.3. If every element in a connected matroid M is fixed, then M is indecomposable.

Proof. Suppose M is decomposable. By Proposition 2.2.2 we can assume that M is the union of two matroids M_1 and M_2 , such that M_1 and M_2 are loopless. If $x \in E(M_1) \cap E(M_2)$, then it follows from Theorem 4.2.1 that x must be non-fixed in M. But every element in M is fixed. Therefore there are no elements in the intersection of $E(M_1)$ and $E(M_2)$. So that M is disconnected, contradicting the fact that M is connected. Hence M is indecomposable.

4.3 Freedom of elements in binary and ternary matroids

In this section we look at which elements in binary and ternary matroids are fixed and not fixed. First we look at binary matroids.

Lemma 4.3.1. [2, Lemma 5.6] Let M be a connected binary matroid having at least two elements and let x be an element of M. If $M \setminus x$ is connected, then x is fixed in M.

As a result of combining Lemma 4.3.1, Proposition 1.8.16 and Lemma 1.8.18 we get the following.

Corollary 4.3.2. With the exception of $U_{2,3}$, every element of a 3-connected binary matroid M is fixed.

Proof. Suppose that M has a non-fixed element a. Then $M \setminus a$ is disconnected by Lemma 4.3.1. But that by Proposition 1.8.16 contradicts the fact that M is 3-connected. By Lemma 1.8.18, $U_{2,3}$ is an exception to this. Therefore every element of a 3-connected binary matroid is fixed, except $U_{2,3}$.

The following theorem provides a link between binary and 3-connected ternary matroids using non-fixed elements in ternary matroids.

Theorem 4.3.3. Let M be a 3-connected ternary matroid. If $x \in E(M)$ and x is not fixed, then $M \setminus x$ is a regular matroid.

Before we prove Theorem 4.3.3 we need definitions. A *well-closed* class of matroids is a class of matroids that is closed under isomorphism, minors,

and duality. A 3-connected matroid N in a well-closed class \mathcal{N} is a universal stabilizer for \mathcal{N} if the following holds for all 3-connected matroids M in \mathcal{N} and all x in E(M): if $M \setminus x$ is 3-connected and has an N-minor, then x is fixed in M, and if M/x is 3-connected and has an N-minor, then x is cofixed in M. For more details on universal stabilizers see [2] and [3]. The next theorem will be used in the proof.

Theorem 4.3.4. [2, Theorem 5.3] Let N be a 3-connected matroid in the well-closed class \mathcal{N} . Then N is a universal stabilizer for \mathcal{N} if and only if the following condition holds for every 3-connected matroid M in \mathcal{N} and every x in E(M):

If $M \setminus x$ is connected with an N-minor, then x is fixed in M, and if M/x is connected with an N-minor, then x is cofixed in M.

Proof of Theorem 4.3.3. Suppose $M \setminus x$ is not binary, then by Theorem 1.7.5 it must have a $U_{2,4}$ -minor. It is shown in [3, Section 2] that $U_{2,4}$ is a universal stabilizer for the class of all ternary matroids. Now if $M \setminus x$ has a $U_{2,4}$ -minor, then it follows that M has a $U_{2,4}$ -minor. Since M is 3-connected, and given that x is not fixed in M, then it follows by Theorem 4.3.4 that $M \setminus x$ is disconnected. By Lemma 1.8.18, the only case where this would happen is when $M \cong U_{2,3}$. But in that case M would not have a $U_{2,4}$ -minor. So we have shown that if M has a $U_{2,4}$ -minor, then it does not have a $U_{2,4}$ -minor, contradiction. Then it must be true that $M \setminus x$ is binary. Moreover, by Proposition 1.7.2 $M \setminus x$ is also ternary. Meaning $M \setminus x$ is binary and ternary. Therefore by Theorem 1.7.8, $M \setminus x$ is regular.

The following corollary shows that by deleting two non-fixed elements from a ternary matroid, it gets disconnected.

Corollary 4.3.5. Let M be a 3-connected ternary matroid not isomorphic to $U_{2,3}$. Then for all x and y in E(M) where x and y are not fixed in M, $M \setminus x, y$ is a disconnected matroid.

Proof. Since M is 3-connected, then by Proposition 1.8.16 $M \setminus x$ is connected. Since x is not fixed, then by Theorem 4.3.3 $M \setminus x$ is a binary matroid. By Proposition 4.1.3 y is not fixed in $M \setminus x$. By Corollary 4.3.2 every element in a 3-connected binary matroid is fixed. So that $M \setminus x$ cannot

be 3-connected. Therefore $M \setminus x$ is a binary connected but not 3-connected matroid containing the non-fixed element y. It follows by Lemma 4.3.1 that $M \setminus x, y$ is disconnected.

4.4 Decomposable binary matroids

4.4.1 Alternative proof of Cunningham's theorem

In this section we give an alternative proof for Cunningham's Theorem, which we also restate for convenience. We also restate and give an alternative proof for Corollary 2.3.16, and give an example showing that matroid union is not closed under duality.

Theorem 4.4.1. [1] A binary matroid M on E is indecomposable if and only if M is connected and $M \setminus e$ is connected for each $e \in E$.

Proof. Let M be a connected binary matroid. If M is decomposable, then by Lemma 4.2.3 it must have at least one non-fixed element x. Consider $M \setminus x$. By Lemma 4.3.1 $M \setminus x$ is disconnected. So if M is decomposable, then it contains at least one element x such that $M \setminus x$ is disconnected. Conversely, if M contains an element x such that $M \setminus x$ is disconnected, then it has a separation (A, B). By Theorem 1.8.6(i) M = S(M/A, M/B). It follows then by Proposition 3.2.1 that $M = (M/A) \vee (M/B)$. But since M = S(M/A, M/B), then both M/A and M/B must have rank strictly greater than zero. Hence M is decomposable.

Next we restate and re-prove Corollary 2.3.16.

Corollary 4.4.2. With the exception of $U_{2,3}$, every 3-connected binary matroid is indecomposable.

Proof. Suppose M is 3-connected and is binary. Then by Corollary 4.3.2, every element in M is fixed. So that by Lemma 4.2.3, M is indecomposable. Since every uniform matroid is decomposable by corollary 2.3.14, then the only exception to this case is $U_{2,3}$.

We can now give a restatement of Cunningham's theorem using the notion of non-fixed elements.

Theorem 4.4.3. Let M be a connected binary matroid. M is decomposable if and only if it contains at least one non-fixed element x. Moreover, if x is not fixed, then $M \setminus x$ is not connected, so that there is a separation (A, B) in $M \setminus x$. In this case $M = (M/A) \vee (M/B)$.

Next we illustrate an application of Lemma 4.2.3. We will show that if M is decomposable, then M^* is not necessarily also decomposable. Consider the following example.

Example 4.4.4. Consider the two matroids M_1 and M_2 where $M_1 \cong M_2 \cong F_7^*$. $E(M_1) = \{a, b, c, d, e, f, h\}$, $E(M_2) = \{h, i, j, k, l, m, g\}$. $S(M_1, M_2)$ on the base-point h is by Proposition 3.2.1 decomposable. The dual of that, however, is $P(M_1^*, M_2^*)$, where $M_1^* \cong M_2^* \cong F_7$. Every element in this parallel connection is fixed as shown in Figure 4.1. Which by Lemma 4.2.3 means it is indecomposable.

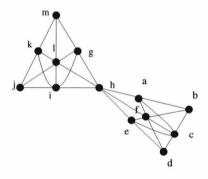


Figure 4.1: A parallel connection of two Fano matroids.

Chapter 5

Decomposable ternary matroids

In this chapter we show what essential property every decomposable 3-connected ternary matroid must have. We also present a theorem showing how a 3-connected ternary matroid can be composed from two other matroids, and also present a partial converse of that. From now on we will be assuming that our matroid union is loopless. First we start by identifying the ternary matroids that for certain are indecomposable, as shown in the next section.

5.1 4-connected ternary matroids are indecomposable

The main result of this section is the following:

Lemma 5.1.1. With the exception of $U_{2,3}$, 4-connected ternary matroids are indecomposable.

First we need a couple of results.

Proposition 5.1.2. If M is a 4-connected ternary matroid, then with the exception of $U_{2,3}$, M has at most one non-fixed element.

Proof. Suppose that M has two non-fixed elements x and y. Then by Corollary 4.3.5 $M \setminus x, y$ is disconnected. But M is 4-connected, and hence by

Proposition 1.8.16 cannot be disconnected by deleting two elements, so we get a contradiction. $U_{2,3}$ is the exception here. Therefore with the exception of $U_{2,3}$, M has at most one non-fixed element.

Proposition 5.1.3. Let M_1 and M_2 be two loopless matroids such that $M = M_1 \vee M_2$. If S is the set of non-fixed elements of M, then $E(M_1) \cap E(M_2) \subseteq S$.

Proof. If $x \in E(M_1) \cap E(M_2)$, then by Theorem 4.2.1 x is not fixed in M. Therefore $x \in S$, so that $E(M_1) \cap E(M_2) \subseteq S$.

Now we prove Lemma 5.1.1.

Proof of Lemma 5.1.1. Suppose M is decomposable. Then by Proposition 2.2.2 there exist loopless matroids M_1 and M_2 such that $M = M_1 \vee M_2$. If M has every element fixed, then by Theorem 4.2.3 M is indecomposable. So assume it has some non-fixed elements. But if M has some non-fixed elements, then by Proposition 5.1.2 it has at most one non-fixed element, call this element a. It follows by Proposition 5.1.3 that $E(M_1) \cap E(M_2) = \{a\}$. By Proposition 3.2.1 $M = S(M_1, M_2)$. So that M has a 2-separation, meaning it is not 4-connected.

5.2 Property of decomposability

The aim of this section is to prove the following theorem, which shows what essential property every decomposable 3-connected ternary matroid has.

Theorem 5.2.1. Let M be a loopless 3-connected ternary matroid not isomorphic to $U_{2,3}$. If M is decomposable, then it contains at least two non-fixed elements p and q.

Proof. If M is 4-connected, then by Lemma 5.1.1 it is indecomposable. So assume that M is not 4-connected. If M is decomposable, then since it is loopless, we can assume by Proposition 2.2.2 that $M = N_1 \vee N_2$, where both N_1 and N_2 are loopless. Since M is connected, then N_1 and N_2 cannot be disjoint. If N_1 and N_2 have only one element in common, then by Proposition 3.2.1 M is a series connection, and hence not 3-connected, giving a

contradiction. So N_1 and N_2 must have at least two elements in common, call them p and q. Since p and q are not loops, then by Theorem 4.2.1, p and q are not fixed in M. So if M is decomposable, then it contains at least two non-fixed elements.

5.3 Composing ternary matroids

The aim of this section is to prove the following theorem which shows how a ternary matroid can be constructed from two other matroids. Recall that if $M = M_1 \vee M_2$ is disconnected, then $M = N_1 \oplus N_2$, for some disjoint matroids N_1 and N_2 . If $M = M_1 \vee M_2$ is 2-connected, then $|E(M_1) \cap E(M_2)| \geq 1$, if $M = M_1 \vee M_2$ is 3-connected, then $|E(M_1) \cap E(M_2)| \geq 2$, and if $M = M_1 \vee M_2$ is 4-connected, then $|E(M_1) \cap E(M_2)| \geq 3$. Due to the fact that 4-connected ternary matroids are indecomposable as shown in Section 5.1, we are not interested in ternary matroids that are 4-connected. For this reason, we shall focus on the union of two matroids that have exactly two elements in common, thus making sure the union is not 4-connected. Note also that having two elements in common, x and y, means that x and y are non-fixed in the union by Theorem 4.2.1. Thus the 3-connected decomposable ternary matroid meets the decomposition property mentioned in Theorem 5.2.1.

Theorem 5.3.1. Let M_1 and M_2 be two loopless regular matroids with $|E(M_1) \cap E(M_2)| \leq 2$, such that $M_1 \vee M_2$ is 3-connected. Then $M = M_1 \vee M_2$ is a ternary matroid.

We first establish some lemmas.

Lemma 5.3.2. Let M be a matroid such that $M = M_1 \vee M_2$ where M_1 and M_2 are both loopless binary matroids with $|E(M_1) \cap E(M_2)| = 1$, then M is a binary matroid.

Proof. Since $|E(M_1) \cap E(M_2)| = 1$, then by Proposition 3.2.1 $M_1 \vee M_2 = S(M_1, M_2)$. But since M_1 and M_2 are both binary, then it follows by Proposition 1.8.14 that $S(M_1, M_2) = M$ is also binary.

Lemma 5.3.3. Let M be a matroid such that $M = M_1 \vee M_2$ where M_1 and M_2 are both loopless binary matroids with $E(M_1) \cap E(M_2) = \{x, y\}$. Then $M \setminus x$ and $M \setminus y$ are both binary matroids.

Proof. Without loss of generality, consider $M \setminus x$. By Proposition 2.4.1 $M \setminus x = (M_1 \vee M_2) \setminus x = (M_1 \setminus x) \vee (M_2 \setminus x)$. $(M_1 \setminus x)$ and $(M_2 \setminus x)$ have only one element in common which is y. Moreover, $(M_1 \setminus x)$ and $(M_2 \setminus x)$ are both binary by Proposition 1.7.2. So by Lemma 5.3.2 $M \setminus x$ is binary. Similarly, $M \setminus y$ is also binary.

Lemma 5.3.4. Let M be a matroid and let $x \in E(M)$. If $M/x \cong U_{2,5}$, then for all $z \in E(M) - \{x\}, M \setminus z$ is non-binary.

Proof. Since $M/x \cong U_{2,5}$, then $M/x \setminus z \cong U_{2,4}$. By Proposition 1.4.2(iii) $M/x \setminus z = M \setminus z/x$. Hence $M \setminus z$ has a $U_{2,4}$ -minor, so that $M \setminus z$ is non-binary.

Lemma 5.3.5. Let M be a matroid and let $x \in E(M)$. If $M \setminus x \cong U_{2,5}$, then for all $z \in E(M) - \{x\}, M \setminus z$ is non-binary.

Proof. Since $M \setminus x \cong U_{2,5}$, then $M \setminus x, z \cong U_{2,4}$. By Proposition 1.4.2(i) $M \setminus x, z = M \setminus z, x$. Meaning $M \setminus z$ has a $U_{2,4}$ -minor, and hence non-binary.

The following proposition will also be used in proving Theorem 5.3.1.

Proposition 5.3.6. [5, Corollary 11.2.20] A 3-connected matroid whose rank and corank exceed two has a minor isomorphic to one of $M(W_3)$, W^3 , Q_6 , P_6 and $U_{3,6}$.

First we prove a lemma that shows how to construct a matroid that has no $U_{2,5}$ or $U_{3,5}$ -minors.

Lemma 5.3.7. Suppose that M_1 and M_2 are both loopless binary matroids with $|E(M_1) \cap E(M_2)| \leq 2$, such that $M_1 \vee M_2$ is 3-connected. Then $M = M_1 \vee M_2$ has no $U_{2,5}$ or $U_{3,5}$ -minors.

Proof. The proof is by induction. First we show the theorem holds for matroids with $|E(M)| \leq 7$, and that becomes our base case. We then show

that it is not possible for M to have a $U_{2,5}$ or $U_{3,5}$ -minors when $|E(M)| \ge 8$. But first we show that $E(M_1)$ and $E(M_2)$ cannot have less that two elements in common.

If $|E(M_1) \cap E(M_2)| = 0$, then M is disconnected, giving a contradiction. If $|E(M_1) \cap E(M_2)| = 1$, then M is a series connection and hence not 3-connected, giving a contradiction. Therefore we may assume from now on that $|E(M_1) \cap E(M_2)| = 2$.

Base case step. Case 1. Suppose that |E(M)| = 5. We break this into two further cases, one for each minor.

Case 1.1. Suppose M has a $U_{2,5}$ -minor, then since |E(M)| = 5, then $M \cong U_{2,5}$. Now we examine the two different possibilities in which M can be constructed from two matroids with two elements in common.

Suppose that M is the union of M_1 and M_2 , where $E(M_1)$ has 5 elements, $E(M_2)$ has 2 elements, with two elements in common as shown in Figure 5.1. Let $E(M) = \{a, b, c, x, y\}$, $E(M_1) = \{a, b, c, x, y\}$ and $E(M_2) = \{x, y\}$. Since any 2-element set in $U_{2,5}$ is independent, and because a and b are only in $E(M_1)$, then $\{a, b\}$ is independent in M_1 . Also, since both M_1 and M_2 are loopless, then any 1-element set in M_2 is independent. So $\{x\}$ is independent in M_2 . But then by matroid union definition, we should have $\{a, b\} \cup \{x\} = \{a, b, x\}$ is independent in M. Contradicting the fact that $M \cong U_{2,5}$ is a rank-2 matroid. Therefore it is not possible for M to have a $U_{2,5}$ -minor when $E(M_1)$ has 5 elements, and $E(M_2)$ has 2 elements, with two elements in common.

Now suppose that M is the union of M_1 and M_2 , where $E(M_1)$ has 4 elements, and $E(M_2)$ has 3 elements, with two elements in common as shown in Figure 5.2. Let $E(M) = \{a, b, c, x, y\}$, $E(M_1) = \{b, c, x, y\}$, and $E(M_2) = \{x, y, a\}$. $\{b, c\}$ is independent in M. So $\{b, c\}$ must be independent in M_1 . Since M_2 is loopless, then $\{a\}$ must be independent in M_2 . Therefore $\{b, c\} \cup \{a\} = \{a, b, c\}$ must be independent in M, contradicting the fact that M is a rank-2 matroid. So M cannot have a $U_{2,5}$ -minor when |E(M)| = 5.

Case 1.2. Suppose M has a $U_{3,5}$ -minor, then since |E(M)| = 5, then $M \cong U_{3,5}$. We will examine the two different possibilities in which M can be constructed from two matroids with two elements in common.

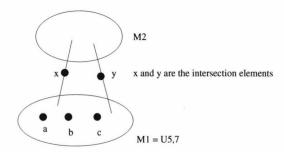


Figure 5.1: $E(M_1)$ has 5 elements and $E(M_2)$ has 2 elements.

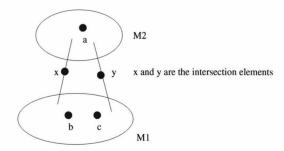


Figure 5.2: $E(M_1)$ has 4 elements and $E(M_2)$ has 3 elements.

Suppose that M is the union of M_1 and M_2 , where $E(M_1)$ has 5 elements, and $E(M_2)$ has 2 elements, with two elements in common as shown in Figure 5.1. Let $E(M) = \{a, b, c, x, y\}$, $E(M_1) = \{a, b, c, x, y\}$ and $E(M_2) = \{x, y\}$. Since any 3-element set in $U_{3,5}$ is independent, and because a, b and c are only in $E(M_1)$, then it must that $\{a, b, c\}$ is independent in M_1 . Element x is not a loop in M_2 , so $\{x\}$ is independent in M_2 . It follows that $\{a, b, c\} \cup \{x\} = \{a, b, c, x\}$ is independent in M, contradicting the fact that $M \cong U_{3,5}$ is a rank-3 matroid.

Now suppose that M is the union of M_1 and M_2 where M_1 has 4 elements, M_2 has 3 elements, with two elements in common as shown in Figure 5.2. Let $E(M) = \{a, b, c, x, y\}$, $E(M_1) = \{b, c, x, y\}$ and $E(M_2) = \{x, y, a\}$. If one of the two matroids is rank-0, then M is indecomposable, giving a contradiction. So one of the two matroids must be rank-1, and the other must be rank-2. Since any two-element set in M is independent, then $\{b, c\}$ must be independent in M. So that $\{b, c\}$ must be independent in M_1 .

Therefore M_1 is the rank-2 matroid, and M_2 is the rank-1 matroid. It follows that any 2-element set in M_2 is a circuit. Since $\{a, x, b\}$ is independent in M, it is the union of two independent sets. But $\{a, x\}$ is a circuit in M_2 . Therefore it must be that $\{b, x\}$ is independent in M_1 . Similarly, the fact that $\{a, x, c\}$ is independent in M, implies that $\{c, x\}$ must be independent in M_1 . Also since the elements b and c are only in $E(M_1)$, then $\{b, c\}$ must be independent in M_1 as well. In fact by following this reasoning we can show that any 2-element set in M_1 is independent. But M_1 is rank-2, meaning any set with size 3 is a circuit. Since M_1 has four elements, then it follows that $M_1 \cong U_{2,4}$. This contradicts the fact that M_1 is binary. So we conclude that when |E(M)| = 5, M has no $U_{2,5}$ or $U_{3,5}$ -minors. Note that this case does not require the fact that the union is 3-connected.

Case 2. Suppose that |E(M)| = 6. We break this into two further cases, one for each minor.

Case 2.1. Suppose M has a $U_{2,5}$ -minor. Since |E(M)| = 6, there are two different ways we can obtain a $U_{2,5}$ -minor from M.

The first one is when $M \setminus p \cong U_{2,5}$ for some $p \in E(M)$. By Lemma 5.3.5, $M \setminus p \cong U_{2,5}$ implies that $M \setminus z$ is non-binary for all $z \in E(M) - \{p\}$. But z can be an intersection element of $E(M_1)$ and $E(M_2)$. So that by deleting z from M, we get by Lemma 5.3.3 a binary matroid, giving a contradiction. Therefore $U_{2,5}$ cannot be obtained by deleting an element from M.

The second possibility is when $M/p \cong U_{2,5}$. By Lemma 5.3.4, $M/p \cong U_{2,5}$ implies that $M \setminus z$ is non-binary for all $z \in E(M) - \{p\}$. If z is an intersection element of $E(M_1)$ and $E(M_2)$, then by Lemma 5.3.3 $M \setminus z$ is a binary matroid. So we get a contradiction. Therefore it is not possible to obtain $U_{2,5}$ from M by contracting or deleting an element.

Case 2.2. Suppose M has a $U_{3,5}$ -minor, then since |E(M)| = 6, there are two different ways we can obtain a $U_{3,5}$ -minor from M.

The first one is when $M \setminus p \cong U_{3,5}$ for some $p \in E(M)$. If p is not in $E(M_1) \cap E(M_2)$, then suppose without loss of generality that $p \in E(M_1)$. By Proposition 2.4.1 $M \setminus p = (M_1 \setminus p) \vee M_2$. In this case $M \setminus p$ is a 5-element decomposable matroid that is the union of two binary matroids with two elements in common. It follows by Case 1.2 that we cannot have $M \setminus p \cong U_{3,5}$. If on the other hand p was in the intersection, then deleting

p from M leaves $E(M_1)$ and $E(M_2)$ with only one element in common. So $M \setminus p$ by Proposition 3.2.1 is a series connection of two matroids. But $U_{3,5}$ is not a series connection of two matroids, so this contradicts the fact that $M \setminus p \cong U_{3,5}$.

The second possibility is when $M/p \cong U_{3,5}$. If p is not in $E(M_1) \cap E(M_2)$, then we lose no generality if we assume that $p \in E(M_1)$. By Proposition 2.4.5 $M/p = M_1/p \vee M_2$. But M/p is a 5-element decomposable matroid that is the union of two binary matroids that have two elements in common. So we arrive at the same contradiction shown in Case 1.2. If p is in the intersection, then let us consider the dual of M/p, to see what possibilities there are for M. Since $M/p \cong U_{3,5}$, then $M^* \setminus p \cong U_{2,5}$, and we have the following possibilities:

- (i) Either p is a coloop in M^* , implying that p is a loop in M, contradicting the fact that M is connected.
- (ii) p is a loop in M^* , implying that p is a coloop in M, contradicting the fact that M is connected.
- (iii) M^* has p in a parallel class with another element, implying that M has p in a series pair with another element. But $M \setminus p$ must be connected by Proposition 1.8.16, giving a contradiction.
- (iv) Or $M^* \cong U_{2,6}$, implying that $M \cong U_{4,6}$. Choose an element z not in $E(M_1) \cap E(M_2)$. Assume without loss of generality that $z \in E(M_1)$. Then by Proposition 2.4.5 $M/z = (M_1/z) \vee M_2$. But $M/z \cong U_{3,5}$ is a 5-element decomposable matroid that is the union of two binary matroids with two elements in common, contradicting what was shown in Case 1.2.
- Case 3. Suppose that |E(M)| = 7. We break this into two further cases, one for each minor.
- Case 3.1. Suppose M has a $U_{2,5}$ -minor, then since |E(M)| = 7, there are three different ways we can obtain a $U_{2,5}$ -minor from M. First one is deleting two elements from M, second one is deleting one element and contracting another, and the third one is contracting two elements from M.

Consider the two ways involving a deletion. Suppose that either $M \setminus p, q \cong U_{2,5}, \ M \setminus p/q \cong U_{2,5}$, or $M \setminus q/p \cong U_{2,5}$. If the deletion element, say p, is in the intersection, then $M \setminus p$ is the union of two binary matroids that have one element in common. It follows by Proposition 3.2.1 that $M \setminus p$ is

a series connection of two binary matroids. So that by Proposition 1.8.14 $M \setminus p$ is a binary matroid too. Hence $M \setminus p$ does not have a $U_{2,5}$ -minor, contradicting the assumption. If none of the deletion elements are in the intersection, then we can delete one of them say p, and get by Proposition 2.4.1 $M \setminus p$ as a 6-element decomposable matroid. But $M \setminus p$ is the union of two binary matroids that have two elements in common, and has a $U_{2,5}$ -minor, which contradicts what was shown in Case 2.1.

Now consider $M/p/q \cong U_{2,5}$. If only one of the elements p and q is in the intersection, say q, then by Proposition 2.4.5 M/p is a 6-element decomposable matroid. It is the union of two binary matroids that have two elements in common, and it has a $U_{2,5}$ -minor. But that contradicts Case 2.1. Similarly, if none of the elements p and q are in the intersection, then M/p and M/q also give the same contradiction. Suppose on the other hand that both p and q are in the intersection. Let N = M/p, so that $N/q \cong U_{2,5}$. It follows that $N^* \setminus q \cong U_{3,5}$. Let us now consider the possibilities for what N could be:

- (i) If q is a coloop in N^* , then it implies that N^* is disconnected, and hence by Proposition 1.5.4 N is disconnected. But that contradicts the fact that M is 3-connected by Proposition 1.8.16.
- (ii) q is a loop in N^* , implying that q is a coloop in N meaning N is disconnected, contradicting Proposition 1.8.16.
- (iii) q is in parallel to another element in N^* , implying that N has a series pair. But if N = M/p has a series pair, then M must have a series pair too, which means M can be disconnected by removing one element, contradicting the fact that it is 3-connected.

Having made the observations (i-iii), we see that element q must be on the same plane as the other elements of $U_{3,5}$. So that N^* must be rank-3. Moreover, q can be collinear with at most two elements in N^* , meaning q cannot create a 2-separation in N^* . So that N^* is 3-connected. But N^* has 6 elements, meaning its corank exceeds two. So it follows from Corollary 5.3.6 that there are only 5 possibilities for N^* , namely $M(\mathcal{W}_3)$, \mathcal{W}^3 , Q_6 , P_6 and $U_{3,6}$ as shown in Figure 5.3. From those 5 matroids, however, only the ones shown in Figure 5.3(a, b and c) have $N^* \setminus q \cong U_{3,5}$. If we examine their duals as shown in Figure 5.3(d, e and f) respectively, we will find that

no matter where the intersection element q is located in N, $N \setminus q$ always has a $U_{2,4}$ -minor. Hence $N \setminus q$ is not binary. If $N \setminus q$ is non-binary, then $M/p \setminus q = M \setminus q/p$ is non-binary. But by Lemma 5.3.3 $M \setminus q$ is supposed to be binary, contradiction. Therefore when |E(M)| = 7, M has no $U_{2,5}$ -minor.

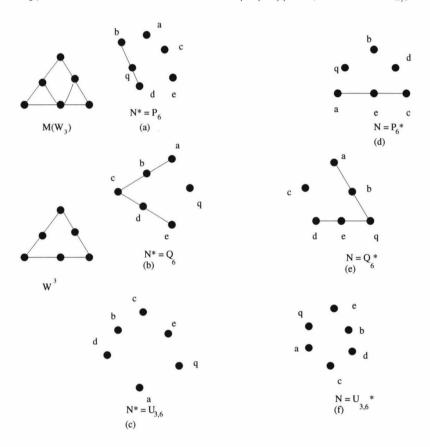


Figure 5.3: Matroid with $U_{3,5}$ -minor.

Case 3.2. Suppose M has a $U_{3,5}$ -minor, since |E(M)| = 7, then we consider the different ways M can have a $U_{3,5}$ -minor.

The two ways involving deletions are $M \setminus p, q \cong U_{3,5}$, $M \setminus p/q \cong U_{3,5}$ or $M \setminus q/p \cong U_{3,5}$. If none of the elements p and q are in the intersection, then $M \setminus p$ having a $U_{3,5}$ -minor contradicts what was shown in Case 2.1. If any of the deletion elements is in the intersection, say p, then $M \setminus p$ is a series connection of two binary matroids. Hence by Proposition 1.8.14 is binary, contradicting the fact that $M \setminus p$ has a $U_{3,5}$ -minor.

The third way is when $M/p/q \cong U_{3,5}$. Suppose none of the two elements are in the intersection. Then we can contract one element, say p, and get by Proposition 2.4.5 M/p as a 6-element decomposable matroid, constructed from the union of two binary matroids with two elements in common, and it must be 3-connected. That contradicts what was shown in Case 2.2. Suppose one of the elements, say p, is in the intersection and q is not. Then M/q is a 6-element decomposable matroid, constructed from the union of two binary matroids with two elements in common, which also must be 3-connected, again contradicting what was shown in Case 2.2. If on the other hand both p and q are in the intersection, then since $M/p/q \cong U_{3,5}$, it follows that $M^* \setminus p \setminus q \cong U_{2,5}$. Now we look at the different possibilities for M^* .

 M^* is a 5-point line in addition to the two other elements p and q. Because p and q are intersection elements, then by Theorem 4.2.1 they are not fixed in M.

- (i) If either p or q is a loop or coloop in M^* , then it is a coloop or loop in M respectively, leading to contradictions as was shown in Case 3.1.
- (ii) If either p or q is in parallel with any other element z in M^* , then that makes $\{p, z\}$ or $\{q, z\}$ a series pair in M. But then it follows that M can be disconnected by deleting one of the series pair elements, contradicting the fact that M is 3-connected. We would get the same contradiction if p and q are parallel to each other.
- (iii) If p or q is in a series pair with any other element z in M^* , then that makes $\{p, z\}$ or $\{q, z\}$ a parallel pair in M respectively, and hence fixing p or q in M. Thereby contradicting the fact that p and q are both non-fixed in M. The same contradiction follows if p and q make a series pair together.

Therefore the only place left for p and q in M^* is on the 5-point line, such that $\{p,q\}$ is independent on the line. Thus making $M^* \cong U_{2,7}$. Hence $M \cong U_{5,7}$. Choose $z \in E(M)$, such that z is not an intersection element. Assume without loss of generality that $z \in E(M_1)$. Since (M_1/z) is binary by Proposition 1.7.2, and M_2 is binary by assumption. Then it follows that $M/z = (M_1/z) \vee M_2 \cong U_{4,6}$ is a 6-element decomposable matroid made from the union of two binary matroids. Moreover, the two matroids still have two elements in common, M/z has a $U_{3,5}$ -minor, and must be 3-connected as well. This contradicts what was shown in Case 2.2. Therefore we conclude

that when |E(M)| = 7, M has no $U_{3,5}$ -minor. Thus we have established a base case for our induction.

Inductive step. Assume that $|E(M)| = |E(M_1) \cup E(M_2)| \ge 8$. Also assume that the result holds whenever $|E(M)| \le 7$, for all pairs of binary matroids M_1 and M_2 such that $E(M_1) \cap E(M_2) = \{x,y\}$. Suppose the result fails for M. Then there exists $X, Y \subseteq E(M)$, such that $(M_1 \vee M_2) \setminus X/Y \cong U_{2,5}$ or $U_{3,5}$. Since $|E(M)| \ge 8$, then there exists z in $(X \cup Y) - \{x,y\}$. Without loss of generality, assume z is in $E(M_1)$, and also assume that z is in X. Then by Proposition 2.4.1 $M \setminus z = (M_1 \setminus z) \vee M_2$. It means that $M \setminus z$ has a $U_{2,5}$ or $U_{3,5}$ -minor, contradicting the fact that the result holds for any number of elements less than the size of E(M). Similarly for the contraction, if z is in $E(M_1)$, and we let $z \in Y$, then since z is not in $E(M_1) \cap E(M_2)$, it follows by Proposition 2.4.5 $M/z = (M_1/z) \vee M_2$. But if M/z has a $U_{2,5}$ or $U_{3,5}$ -minor, then that contradicts the fact that the result holds for any number of elements less than the size of E(M). Therefore we conclude that M has no $U_{2,5}$ or $U_{3,5}$ -minor.

To complete the proof of Theorem 5.3.1, we use a result that has been proven by M.J.Piff and D.J.A.Welsh, see [7] for more details. We also state it here for the reader's convenience. Note that the theorem makes no assumption about the ground sets being the same.

Theorem 5.3.8. Let M_1 and M_2 be matroids on E_1 and E_2 respectively. There exists n such that if M_1 and M_2 are representable over a field F, with $|F| \geq n$, then $M_1 \vee M_2$ is also representable over F.

Proof of Theorem 5.3.1. Since M_1 and M_2 are regular, then they are both binary and ternary, and therefore by Lemma 5.3.7 we know that $M_1 \vee M_2$ has no $U_{2,5}$ or $U_{3,5}$ -minors. By Theorem 1.7.6, if we show that M has no F_7 or F_7^* -minors, then that would complete the proof of M being ternary.

Suppose now that M_1 and M_2 are both regular, such that $M_1 \vee M_2$ has an F_7 or F_7^* -minors. Since M_1 and M_2 are regular, then by Theorem 1.7.7 they are representable over every field, and hence they are representable over the reals. It follows then by Theorem 5.3.8 that $M_1 \vee M_2$ is also representable over the reals. But F_7 and F_7^* are minors of $M_1 \vee M_2$ by assumption. So by Proposition 1.7.2 we must have F_7 and F_7^* as representable over the reals

as well, contradicting Theorem 1.7.4. This concludes the proof for Theorem 5.3.1 that $M = M_1 \vee M_2$ is a ternary matroid.

As a result of combining Theorem 5.3.1 with Theorem 3.4.1 we get the following corollary.

Corollary 5.3.9. Let M_1 and M_2 be two loopless regular matroids with $E(M_1) \cap E(M_2) = \{x, y\}$. If M_1 and M_2 are both connected, neither of them has a series pair, $|E(M_1)|, |E(M_2)| \geq 3$ and x and y are split in M_1 and split in M_2 , then $M = M_1 \vee M_2$ is a 3-connected ternary matroid.

5.4 Decomposing ternary matroids

The aim of this section is to prove the following theorem, which is a partial converse to Theorem 5.3.1.

Theorem 5.4.1. Let $M = M_1 \vee M_2$ be a 3-connected ternary matroid, such that $|E(M_1) \cap E(M_2)| = 2$. If M_1 and M_2 are 3-connected, then M_1 and M_2 are both regular matroids.

A matroid M uses a set T, if $T \subseteq E(M)$. A class \mathcal{N} of matroids is t-rounded, if every member of \mathcal{N} is (t+1)-connected, and the following condition holds: If M is a (t+1)-connected matroid having an \mathcal{N} -minor, from the class and X is a subset of E(M) with at most t elements, then M has an \mathcal{N} -minor using X. We are particularly interested in 2-roundedness. The following results will be used in the proof.

Proposition 5.4.2. [5, Proposition 11.2.13] Every 3-connected matroid M having at least four elements has a minor isomorphic to $U_{2,4}$ or $M(K_4)$.

Proposition 5.4.3. [5, Corollary 11.2.14] Every 3-connected binary matroid having at least four elements has a minor isomorphic to $M(K_4)$.

Proposition 5.4.4. [5, Proposition 11.3.8] Let M be a 3-connected matroid having a $U_{2,4}$ -minor and suppose that e and f are distinct elements of M. Then M has $U_{2,4}$ -minor using $\{e, f\}$.

Theorem 5.4.5. [5, Proposition 11.3.9] Let t be 1 or 2 and \mathcal{N} be a collection of (t+1)-connected matroids. Then \mathcal{N} is t-rounded if and only if the following condition holds: If M is a (t+1)-connected matroid having an \mathcal{N} -minor such that |E(M) - E(N)| = 1, and X is a subset of E(M) with at most t elements, then M has an \mathcal{N} -minor using X.

The following follows from Theorem 5.4.5.

Proposition 5.4.6. [5] The collection of matroids $\{U_{2,4}, M(K_4)\}$ is 2-rounded.

Next we prove two lemmas that will also be used in proving Theorem 5.4.1.

Lemma 5.4.7. Let N be a connected matroid that has at least 3 elements, such that N has no series pairs. If $E(N) = A \cup \{x,y\}$, then N/A consists entirely of loops.

Proof. We first observe that $E(N/A) = \{x, y\}$. Let us consider what possibilities x and y have in N/A.

- Case 1. $\{x,y\}$ are independent in N/A. Meaning that x and y are coloops of N/A. But a coloop cannot be generated by a series of contractions. It follows that N must have had a coloop to start with, contradicting the fact that it is connected.
- Case 2. $\{x, y\}$ is a parallel pair in N/A. Here there are two possibilities for x and y to be a parallel pair.
- Case 2.1. N has a minor containing three elements x, y and a, where $\{x, y\}$ is a parallel pair, and a as a coloop in this minor, as shown in Figure 5.4(a). It follows that a must have been a coloop of N, giving a contradiction.
- Case 2.2. N has a minor isomorphic to $U_{2,3}$ containing the elements x, y and a, as shown in Figure 5.4(b). It follows that $\{x, y\}$ is a series pair of this minor, and therefore must have been a series pair of N, giving a contradiction.
- Case 3. One of x and y is a loop in N/A, and the other is not. Assume without loss of generality that x is a loop in N/A. It follows that N must have a minor where x is in a parallel class, and y is not. So that y is a coloop of this minor, and hence a coloop of N, giving a contradiction.



Figure 5.4: x and y as parallel pairs.

Therefore it follows that the only possibility left for x and y, is for x to be a loop in N/A and y to be a loop in N/A. Hence N/A consists entirely of loops.

Lemma 5.4.8. Let $M = M_1 \vee M_2$, such that $E(M_1) \cap E(M_2) = \{x, y\}$. If both M_1 and M_2 are connected, $|E(M_1)| \geq 3$, $|E(M_2)| \geq 3$, and neither of them has $\{x, y\}$ as a series pair, then $M/(E(M_1) - \{x, y\}) = M_2$. Similarly, $M/(E(M_2) - \{x, y\}) = M_1$.

Proof. Let $E(M_2) - \{x, y\} = A$, and $E(M_1) - \{x, y\} = B$. Without loss of generality, consider M/A. By Proposition 2.4.5 $M/A = M_1 \vee M_2/A$. But $E(M_2/A) = \{x, y\}$. So it follows by Lemma 5.4.7 that M_2/A consists entirely of loops. Hence $M/A = M_1$. Similarly, $M/B = M_2$.

Now we prove Theorem 5.4.1.

Proof of Theorem 5.4.1. Let $E(M_1) \cap E(M_2) = \{x,y\}$. Since M is the union of two matroids with exactly two elements in common, and M is 3-connected, then by Theorem 3.4.1 it follows that both M_1 and M_2 have to be connected. Moreover, each must have size at least 3, and neither of them has a series pair. It follows then by Lemma 5.4.8 that both M_1 and M_2 are minors of M. Hence by Proposition 1.7.2 they are both ternary. So what we need to show now to complete the proof, is that M_1 and M_2 are both binary. Assume that both M_1 and M_2 are 3-connected. If both M_1 and M_2 are binary, then we are done. So consider the following possibilities.

Case 1.1. Only one of M_1 and M_2 is not binary. Assume without loss of generality that M_1 is not binary and M_2 is binary. By Proposition 5.4.4, M_1 has a $U_{2,4}$ -minor using x and y. So let $N_1 = M_1 \setminus A/B \cong U_{2,4}$ where A and B are subsets of $E(M_1)$ not containing x or y. So that by Propositions

2.4.1 and 2.4.5 $M \setminus A/B = (M_1 \setminus A/B) \vee M_2 = N_1 \vee M_2$. Since M_2 is binary and 3-connected, then by Propositions 5.4.3 and 5.4.6, it must have a minor isomorphic to $M(K_4)$ using x and y. Call this minor N_2 . From $M(K_4)$ we can obtain a $U_{1,3}$ -minor. It can be obtained by contracting two elements and deleting one element from $M(K_4)$, other than x and y. So let $c, d \in E(N_2)$ be the contraction elements, and $e \in E(N_2)$ be the deletion element. Then $M \setminus A/B/c/d \setminus e \cong (U_{2,4} \vee U_{1,3})$, where both $U_{2,4}$ and $U_{1,3}$ use $\{x,y\}$. But in this case $U_{2,4} \vee U_{1,3} \cong U_{3,5}$ as shown in Figure 5.5, which is not ternary, contradicting the fact that M is ternary.

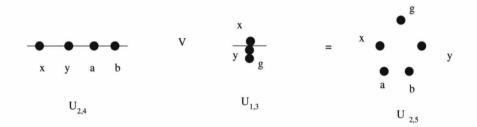


Figure 5.5: $U_{2,5}$ -minor.

Case 1.2. Both M_1 and M_2 are not binary. By Proposition 5.4.4 both M_1 and M_2 have a $U_{2,4}$ -minor using $\{x,y\}$. So that $M_1 \setminus A/B \cong U_{2,4}$, where $A, B \subseteq E(M_1)$ not containing x and y, and $M_2 \setminus C/D \cong U_{2,4}$, where $C, D \subseteq E(M_2)$ not containing x and y. By Propositions 2.4.1 and 2.4.5 $M \setminus A/B \setminus C/D = (U_{2,4} \vee U_{2,4}) \cong U_{4,6}$. But $U_{4,6}$ is not ternary, contradicting the fact that M is ternary. Hence we have shown that if M_1 and M_2 are 3-connected, then they are regular.

As noted earlier, Theorem 5.4.1 is a partial converse of Theorem 5.3.1. The next proposition is not used in this thesis, but together with Conjecture 6.1.2, which will be stated in the next chapter, makes a complete proof for a full converse of Theorem 5.3.1.

Proposition 5.4.9. Let M be a non-binary connected but not 3-connected matroid, where $\{x,y\} \subseteq E(M)$. If x and y are split in M, then M has a $U_{2,4}$ -minor using x and y.

Proof. The proof is by induction, first we prove the theorem is true for the base case when |E(M)| = 4.

Base case step. If |E(M)| = 4, then since $\{x, y\} \subseteq E(M)$, and M has a $U_{2,4}$ -minor, then $M \cong U_{2,4}$ using x and y.

Inductive step. Assume now that $|E(M)| = n \geq 5$, and assume that the theorem holds for all $|E(M)| \leq n$. If M has a 2-separation (A, B) that splits x and y, then let us assume without loss of generality that $x \in A$ and $y \in B$. Let M_A and M_B be two matroids such that $M = M_A \oplus_2 M_B$. By Proposition 1.8.21 there must exist an intersection element between M_A and M_B . So let $E(M_A) \cap E(M_B) = \{p\}$ be the base-point. By property of 2-sum $M_A \setminus p = M|A$ and $M_B \setminus p = M|B$. If both M_A and M_B are binary, then by Proposition 1.8.11, their 2-sum, which is M, is also binary, giving a contradiction. So at least one of M_A and M_B is non-binary. Assume without loss of generality that M_B is non-binary. If M_B has a $U_{2,4}$ -minor using x and y, then we are done, so assume that is does not have a $U_{2,4}$ -minor using x and y.

Consider M_B . Replace element p by x in M_B , and call the resulting matroid N. N is a minor of M isomorphic to M_B . So N is non-binary. If

N is 3-connected, then by Proposition 5.4.4 it has a $U_{2,4}$ -minor using x and y and we are done. So let us suppose that it is not 3-connected, so that it has a 2-separation (P,Q). If every 2-separation in N splits x and y, then the inductive argument implies that N has a $U_{2,4}$ -minor using x and y and we are done. So let us suppose that x and y are not split in N. Without loss of generality assume that both $x, y \in P$.

Now replace x by p in N. The resulting matroid is M_B again like it was at the start. The difference now, however, is that we know that M_B has a 2-separation such that p and y are not split in M_B . Suppose M_B has a 2-separation (P',Q), such that both $p,y \in P'$. Then let M(P') and M(Q) be two matroids such that $M_B = M(P') \oplus_2 M(Q)$, and let $E(M(P')) \cap E(M(Q)) = \{z\}$ be the base-point, so that $P' \cup \{z\} = E(M(P'))$, $Q \cup \{z\} = E(M(Q))$. We have the following:

$$M = M_A \oplus_2 (M(P') \oplus_2 M(Q)). \tag{5.4.1}$$

Let (Q, R) be a partition of M, such that $R = P' - \{p\} \cup A$. Due to the fact that the 2-sum of M_A and M_B was applied between M_A and the set P' which contains p, we get

$$M_A \oplus_2 (M(P') \oplus_2 M(Q)) = (M_A \oplus_2 M(P')) \oplus_2 M(Q).$$
 (5.4.2)

Where $E(M_A \oplus_2 M(P')) = R \cup \{z\}$. Therefore (R,Q) is a 2-separation of M. Moreover, x and y are both in the same part of the 2-separation R. But that contradicts the assumption that M has no 2-separation that does not split x and y. Therefore N must split x and y. Since $|E(N)| \leq n$, then N has a $U_{2,4}$ -minor using x and y. Since this was the case for |E(M)| = n, then the theorem holds for all values of n by induction.

Chapter 6

What is next?

6.1 Conjectures

In this chapter, we state some conjectures that we believe are true, but did not spend a lot of time thinking about due to severe time constraints. To resolve these issues would strengthen some of the results presented in the preceding chapters.

Chapter 5: Decomposable ternary matroids

Section 5.2. Theorem 5.2.1 shows that if a matroid M is decomposable, then it must have at least two non-fixed elements p and q. The next conjecture is the converse of Theorem 5.2.1.

Conjecture 6.1.1. Let M be a 3-connected ternary matroid that has two non-fixed elements p and q. Let (A, B) be a separation of $M \setminus p, q$. Then M is decomposable such that $M = (M/A) \vee (M/B)$.

Section 5.4. Theorem 5.4.1 has shown that a 3-connected ternary matroid decomposes into two regular matroid, provided that the two matroids are 3-connected. To be able to have a stronger version of the theorem, the following conjecture is needed.

Conjecture 6.1.2. Let M be a connected but not 3-connected binary matroid. If x and y are split in M, where $\{x,y\} \subseteq E(M)$, then M has a minor isomorphic to the matroid shown in Figure 6.1.



Figure 6.1: Minor of a binary matroid that has x and y spilt.

If the above conjecture is true, then by combining it with Proposition 5.4.9, we can prove the following conjecture, which is a full converse of Theorem 5.3.1.

Conjecture 6.1.3. Let $M = M_1 \vee M_2$ be a 3-connected ternary matroid, such that $|E(M_1) \cap E(M_2)| = 2$. Then both M_1 and M_2 are regular matroids.

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