Automated Design of Production Scheduling Heuristics: A Review

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Abstract—Hyper-heuristics have recently emerged as a power-² ful approach to automate the design of heuristics for a number 3 of different problems. Production scheduling is a particularly 4 popular application area for which a number of different hyper-5 heuristics have been developed and are shown to be effective, 6 efficient, easy to implement, and reusable in different shop con-7 ditions. In particular, they seem to be a promising way to tackle 8 highly dynamic and stochastic scheduling problems, an aspect 9 that is specifically emphasized in this survey. Despite their success 10 and the substantial number of papers in this area, there is cur-11 rently no systematic discussion of the design choices and critical 12 issues involved in the process of developing such approaches. This 13 paper strives to fill this gap by summarizing the state-of-the-art 14 approaches, suggesting a taxonomy, and providing the inter-15 ested researchers and practitioners with guidelines for the design 16 of hyper-heuristics in production scheduling. This paper also 17 identifies challenges and open questions and highlights various 18 directions for future work.

Index Terms—Evolutionary design, genetic programming (GP),
 hyper-heuristic, scheduling.

I. INTRODUCTION

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²² **S** CHEDULING is concerned with the allocation of limited ²³ **S** CHEDULING is concerned with the basic aim to ensure ²⁴ an effective and efficient use of the available resources. A clas-²⁵ sic problem area is the scheduling of manufacturing systems, ²⁶ in which machines (the resources) have to be allocated to jobs ²⁷ (the tasks) in the best possible way (minimizing or maximizing ²⁸ some objective function). Some of the costs that are typically ²⁹ affected by a production schedule are the holding costs of ³⁰ in-process inventory, contractual penalties for late deliveries, ³¹ setup costs, and the costs of scrap and rework, which illustrate ³² the importance of production scheduling to manufacturers in ³³ their endeavor to become and remain competitive.

A number of exact solution methods that solve deterministic scheduling problems optimally have been proposed in [1]. However, due to the high complexity of most schedulproblems of interest, exact methods are usually unable to solve large instances within a reasonable computational time. Moreover, many problems are stochastic and dynamic,

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i.e., they are subject to change over time due to random, ⁴⁰ stochastic events such as new job arrivals, stochastic processing times, or machine breakdowns. Consequently, many ⁴² researchers and practitioners have turned to heuristics, which ⁴³ deliver acceptable, but not necessarily optimal, solutions in a short computational time. ⁴⁵

In general, heuristics are problem-specific solution meth-46 ods and have to be designed for the problem at hand. 47 Unfortunately, the design of sophisticated heuristics is usu-48 ally a tedious trial-and-error process, with candidate heuristics 49 tested on some instances of the considered problem, modified, 50 and retested until they meet the demands for actual imple-51 mentation, which requires a significant amount of expertise, 52 time, and coding effort. To handle this issue, various meth-53 ods to (partially) automate the design of heuristics have been 54 proposed in the literature, also known as hyper-heuristics. 55

In [2], hyper-heuristics are defined as "an automated 56 methodology for selecting or generating heuristics to solve 57 hard computational search problems." In other words, hyper-58 heuristics explore a search space of heuristics to discover those 59 that work effectively. In this survey, we use fitness to denote 60 the effectiveness of heuristics (discussed in Section III-E), 61 whereas the objective value or function denotes the quality 62 of a schedule. 63

Burke et al. [2] classified hyper-heuristics with respect to 64 the nature of their process, i.e., whether they select or gen-65 erate heuristics. Moreover, they distinguish hyper-heuristics 66 that learn online, i.e., while solving a problem instance, from 67 those that learn offline, i.e., that gather reusable knowledge 68 from a set of training instances. Burke et al. [3] provided 69 a general overview of the state-of-the-art of hyper-heuristic 70 design, covering all categories of hyper-heuristics. The scope 71 of this survey is on (offline) hyper-heuristics for the gen-72 eration of a reusable heuristic, which can be applied to 73 quickly solve new problem instances once it has been gen-74 erated. We deliberately do not cover hyper-heuristics that select a heuristic for every decision point of a particular 76 problem instance (see [4]-[6]), as the generated sequence 77 of heuristics can generally not be reused, nor do we cover 78 hyper-heuristics that learn to select heuristics for a given 79 problem instance (see [7]–[9]), as this is problem clas-80 sification rather than heuristic generation. Hyper-heuristics 81 for the generation of heuristics have also been developed 82 for other problem areas, including bin packing [10]-[12], 83 vehicle routing [13]–[15], timetabling [16], [17], air traffic 84 control [18], [19], and project scheduling [20]. 85

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Problem class		References
Shop configuration	Single machine Parallel machines Job shop	[21], [22], [23], [24], [25], [26], [27], [28], [29], [30], [31], [32], [33] [34] [26], [31], [35], [36], [37], [38], [39], [40], [41], [42], [43], [44], [45], [46], [47], [48], [49], [50], [51], [52], [53], [54], [55], [56]
	Flexible job shop Flow shop	[57], [58], [59], [60], [61], [62], [63], [64], [65][25], [66], [67]
Special processing characteristic	Sequence-dependent setups Job precedence constraints Batch processing Machine eligibility restrictions Dynamic/stochastic environment	[28], [31], [32], [33], [34], [59], [61], [63], [65] [31] [27], [65] [60], [61], [62], [64] [23], [35], [36], [38], [40], [41], [44], [47], [51], [52], [53], [54], [55], [56], [57], [58], [65]
Objective	Completion time-based Due date-based	[23], [24], [25], [27], [28], [29], [31], [34], [36], [37], [39], [42], [44], [45], [46], [48], [49], [50], [52], [53], [55], [57], [61], [64], [66] [22], [23], [25], [26], [27], [29], [30], [31], [33], [34], [35], [36], [38], [40], [41],
	Multi-objective	[43], [47], [48], [49], [50], [54], [55], [58], [59], [60], [62], [63], [66], [64], [65], [67] [32], [51], [56]

 TABLE I

 Scheduling Problems Addressed by Heuristic Generation Hyper-Heuristics

This paper presents a comprehensive review of the literature This paper presents a comprehensive review of the literature Provide the inter-Provide the inter-Provi

99 II. SCHEDULING ENVIRONMENTS AND HEURISTICS

In general, heuristics are designed to be effective for a 100 101 specific problem or class of problems. Production schedul-102 ing problems can be categorized by various properties, two ¹⁰³ important ones being the shop configuration and the objective. The simplest "shop" configuration is a single machine that is 104 105 responsible for processing all jobs. If there is more than one machine available to process a job, this is called a parallel 106 107 machine environment. Multistage problems, which are generally NP-hard [68], are characterized by jobs that consist of a 108 number of processing steps, or operations, that need to be per-109 110 formed on distinct machines in a specified order. Depending 111 on whether all jobs share the same processing order or not, the 112 configuration is called a flow shop or job shop. Flow shops 113 and job shops are further called flexible, if they contain at least ¹¹⁴ one work center that consists of parallel machines [1, Ch. 2]. 115 Objectives can be broadly classified as completion time based, with a focus on the efficiency of the manufacturing system, 116 117 and due date based, with a focus on adherence to promised ¹¹⁸ delivery dates. Table I provides a summary of the scheduling ¹¹⁹ problems that have been addressed in the literature by means 120 of a hyper-heuristic.

¹²¹ In most cases, the heuristics generated by the respec-¹²² tive hyper-heuristic belong to the class of dispatching ¹²³ rules. Dispatching rules are a particularly simple type of scheduling heuristic, which progressively construct solutions 124 by scheduling one operation at a time. Whenever a machine 125 is available and there are jobs waiting to be processed on 126 that machine, dispatching rules compute a priority index for 127 each eligible job as a function of some job attributes (e.g., 128 its processing time or due date), and shop attributes (e.g., 129 the average processing time in the queue of the considered 130 work center), and schedule only the imminent operation of 131 the job with the highest priority. Due to their locally restricted 132 horizon, dispatching rules have very low computational and 133 information requirements, irrespective of the complexity of 134 the overall problem. Moreover, because each scheduling deci- 135 sion is made at the latest possible moment, i.e., immediately 136 before its implementation, dispatching rules naturally possess 137 the ability to quickly react to unexpected changes, which 138 makes them particularly suited for stochastic and dynamic 139 scheduling problems (for a list of papers explicitly addressing 140 stochastic dynamic environments, see Table I). These prop-141 erties, together with their simple and intuitive nature, their 142 ease of implementation and their flexibility to incorporate 143 domain knowledge and expertise [69] explain the wide usage 144 of dispatching rules in practice [70] and the ongoing research 145 on the development of new, more effective dispatching 146 rules (see [71]–[73]). 147

While dispatching rules that have been trained on a 148 set of static, deterministic problem instances could, in 149 principle, be applied to dynamic, stochastic problems, 150 Hildebrandt et al. [44] and Nguyen et al. [49] showed that this 151 does not necessarily lead to good results, and that it is bet- 152 ter to use dynamic, stochastic problems also during training. 153 In terms of hyper-heuristic design, there are some minor differ- 154 ences between using deterministic or stochastic problems for 155 training, which will be discussed in the corresponding sections. 156 In particular, other attributes may be needed (Section III-B), 157 and the fitness function becomes stochastic (Section III-E), 158 which in turn raises issues such as the determination of an 159 appropriate run length of the simulation (Section III-E2). 160 Moreover, the definition of stochastic benchmark problems is 161 also more difficult (Section IV-D). 162

 Supervised learning
 Unsupervised learning

 Parametric representation
 [21], [38], [42], [46]
 [35], [36], [41], [45], [47], [52], [57], [58], [59], [67]

 [24], [30], [39]
 [22], [23], [25], [26], [27], [28], [29], [31], [32], [33], [34], [37], [40], [43], [44], [48], [49], [50], [51], [52], [53], [54], [55], [56], [60], [61], [62], [63], [64], [65], [66]

 TABLE II

 Classification of Hyper-Heuristics for the Generation of Production

 Scheduling Heuristics by Learning Method and Representation

In the literature, dispatching rules are typically designed for 163 164 and tested on a (flexible) job shop problem, which is reflected ¹⁶⁵ in Table I by the relatively large number of studies deal-166 ing with this shop configuration. In general, hyper-heuristics 167 have been used to evolve dispatching rules for a variety of scheduling problems with various objective functions and pro-168 169 cessing characteristics. Also, some recent work has focussed on the development of hyper-heuristics that can evolve a set of 170 171 Pareto-optimal dispatching rules for multiobjective problems. general conclusion of these studies is that hyper-heuristics 172 A ¹⁷³ are able to generate dispatching rules that outperform manually designed benchmark rules. 174

A few researchers have used hyper-heuristics for the gen-175 ¹⁷⁶ eration of other types of scheduling heuristics. Yin et al. [23] evolved the so-called predictive heuristics, which aim to con-177 struct schedules that are robust to unpredictable breakdowns 178 of machines and are shown to outperform a benchmark heuris-179 from the literature. Vázquez-Rodríguez and Ochoa [66] 180 tic evolved variants of the iterative Nawaz, Enscore, and 181 182 Ham (NEH) heuristic [74] for a number of permutation flow 183 shop problems, which are significantly better than the original NEH heuristic and a randomized version. Mascia et al. [67], 184 185 also generated iterative heuristics for a permutation flow shop 186 problem. Park et al. [33] and Nguyen et al. [50] employed 187 a hyper-heuristic for the generation of iterative dispatching 188 rules and variants of a size limited beam search heuristic. 189 These iterative scheduling heuristics evaluate (partial) can-190 didate solutions, and are thus restricted to static and deterministic problems. As in the case of dispatching rules, a key 191 192 component of the above scheduling heuristics is their prior-¹⁹³ ity function (or index), which is generally the part that is 194 evolved by the hyper-heuristic. Hence, the subsequent dis-195 cussion will focus on the evolution of dispatching rules, and 196 priority functions in particular.

197 III. HYPER-HEURISTIC DESIGN CHOICES

Fig. 1 shows a simplified outline of the procedure of a hyper-heuristic for the generation of heuristics. The main components in the design of such a hyper-heuristic concern the encoding or representation of candidate heuristics, which defines the search space, the optimization algorithm to explore this search space, and the fitness function to determine the quality of candidate heuristics. In this survey, we classify the existing hyper-heuristics according to the learning method they

Fig. 1. Basic procedure of a hyper-heuristic for the generation of heuristics.

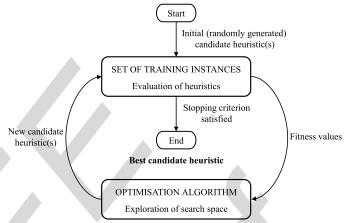
adopt (supervised or unsupervised) and their representation of 2006 candidate heuristics (parametric or grammar based), as shown 2007 in Table II. 2008

Section III-A discusses the two types of learning methods ²⁰⁹ used within hyper-heuristics, followed by a discussion of the ²¹⁰ selection of attributes to be provided to the hyper-heuristic in ²¹¹ Section III-B. The different representations of priority func- ²¹² tions are presented in Section III-C together with suitable ²¹³ optimization algorithms as they are closely tied to the cho- ²¹⁴ sen representation. Section III-D discusses the definition of ²¹⁵ the eligible job set, and Section III-E discusses appropriate ²¹⁶ fitness functions for the evaluation of candidate heuristics. ²¹⁷

1.	Learning Method	2	218
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All hyper-heuristics generate new heuristics by learning ²¹⁹ from a set of training instances. This learning can be supervised or unsupervised. The basic idea of hyper-heuristics using ²²¹ supervised learning in scheduling is to supply the hyperheuristic with a number of very good (preferably optimal) ²²³ schedules that it uses to derive a priority function that reproduces these schedules as closely as possible. These priority ²²⁵ functions can then be used as part of a heuristic, e.g., a ²²⁶ dispatching rule, to solve other problem instances. ²²⁷

A variety of such supervised hyper-heuristics have 228 been proposed in the literature. El-Bouri *et al.* [21] 229 and Eguchi *et al.* [38] developed hyper-heuristics that 230 operate on a neural network representation and use a 231 back-propagation optimization algorithm to learn from 232



²³³ optimal schedules. Ingimundardottir and Runarsson [46] fol-234 lowed the same approach but use logistic regression. Weckman *et al.* [42], also proposed a neural network 235 236 hyper-heuristic, based on a variant of the back-propagation 237 algorithm, but learn from solutions generated by an evolution-238 ary algorithm (EA) instead of optimal solutions. Similarly, Koonce and Tsai [39] employed attribute-oriented induction 239 derive decision rules that reproduce the sequences generto 240 ated by an EA. Li and Olafsson [24] and Olafsson and Li [30] 241 242 developed a hyper-heuristic that generates priority functions in 243 the form of decision trees. However, they only learn from solu-244 tions obtained by some simple dispatching rules, which often 245 generate solutions far from optimal.

While some of the above studies report promising results, 246 major drawback of the supervised learning approach is 247 a 248 that good global schedules can only be obtained for static 249 problems of relatively small size and low complexity, which ²⁵⁰ limits the applicability of these hyper-heuristics. In contrast, ²⁵¹ hyper-heuristics using unsupervised learning generate effective ²⁵² scheduling heuristics by simply applying candidate heuristics a set of problem instances (the training instances), measur-253 to ing their performance, and using this feedback to guide the 254 search toward promising areas of the search space. Hence, 255 unsupervised hyper-heuristics can be applied with relative 256 257 ease to any scheduling problem that can be simulated, which ²⁵⁸ makes them more practical in general. This is also reflected by 259 the fact that most studies in the area of production schedul-²⁶⁰ ing develop hyper-heuristics that are based on unsupervised 261 learning (see Table II), and the subsequent discussion hence 262 concentrates on those.

263 B. Attributes

Irrespective of the representation used, an important design 264 265 decision concerns the selection of adequate job and shop ²⁶⁶ attributes that form the components of the priority functions that can be evolved. To distinguish jobs from each other and 267 be able to prioritize one over another, it is obviously necessary 268 include some job attributes, whereas the inclusion of shop to 269 270 attributes allows for the generation of rules that can adapt to changing shop conditions. Moreover, in the special case of 271 272 evolving iterative scheduling heuristics that make use of the 273 characteristics of candidate solutions in solving a problem, 274 as proposed by Park et al. [33] and Nguyen et al. [49], the ²⁷⁵ hyper-heuristic has to be provided with some attributes related 276 to the current candidate solution, e.g., the realized completion 277 time of a job. In general, the challenge is to select all the 278 relevant attributes while keeping the search space as small as 279 possible. Table III lists a number of promising attributes that 280 have been commonly used in the development of effective dispatching rules in the literature, where shop attributes are 281 282 divided into attributes that concern the work center for which dispatching decision is being made and global attributes. 283 a Attributes should be carefully chosen in consideration of the 284 285 given scheduling problem, e.g., there is no benefit in providing ²⁸⁶ the hyper-heuristic with due date attributes when the objective ²⁸⁷ function is not due date based (see [44]), and certain attributes 288 do not make sense in a dynamic scheduling environment with new jobs arriving all the time (e.g., sum of all processing ²⁸⁹ times).

An important question regarding the selection of attributes 291 attributes in their most basic 292 is whether to include in some aggregate form. To illustrate, a number 293 or researchers provide their hyper-heuristics with the 294 of job due date d_i and the current time t as separate 295 attributes [23], [25], [31], [32], [40], [62], [64]. However, it 296 could be argued that due dates are more meaningful if they 297 are expressed relative to the current time, i.e., $d_j - t$, and 298 that integrating the absolute due date with other job attributes 299 directly will often lead to rules that change their behavior over 300 time, which is generally questionable and particularly unsuit- 301 able for long dynamic scheduling problems. In fact, the term 302 $d_i - t$ appears in many effective manually developed dispatch- 303 ing rules [75]-[77] and several studies on hyper-heuristics 304 have resorted to including due dates (and also release dates 305 and arrival times) in the set of attributes in their relative 306 form [35], [36], [44], [65]. It may make sense to aggregate 307 attributes even further, e.g., to define the non-negative slack 308 $\max(d_j - t - p_j^{r}, 0)$ [29], [31], [41] or the non-negative time 309 to arrival max $(r_i^i - t, 0)$ [29], [31] as one attribute to distin- 310 guish jobs on schedule from late jobs and jobs arriving in the 311 future from waiting jobs, respectively, where r_j^{i} denotes the ³¹² arrival time of job *j* to the work center required for its immi- ³¹³ nent operation. Kuczapski et al. [45] and Baek and Yoon [58] 314 selected a number of composite priority indices of dispatching 315 rules from the literature as attributes for their hyper-heuristics. 316 However, these priority indices may integrate attributes in a 317 suboptimal manner and restrict the hyper-heuristic in its search 318 for a better priority function. Overall, it appears that it is best 319 to provide a hyper-heuristic with attributes in their most basic 320 form and let the hyper-heuristic search for good combinations 321 unless there is a good theoretical foundation for an aggregate 322 attribute, as in the cases above. 323

Another question related to the selection of attributes 324 is whether or not to normalize them to a similar scale. 325 In some cases, this may be necessary to fit a certain 326 representation, e.g., the grammar-based representation by 327 Nguyen et al. [49] or the neural network representa- 328 tion by Branke et al. [52]. Hershauer and Ebert [35], 329 Eguchi et al. [41], and Baek and Yoon [58] also scaled 330 the attributes to a similar range. In a recent study, 331 Branke et al. [52] tested two hyper-heuristics for the gen- 332 eration of dispatching rules (one operating on a parametric, 333 the other on a grammar-based representation of priority func- 334 tions) with and without normalized attributes. They find that 335 normalizing the attributes improves the performance of both 336 hyper-heuristics, especially when the original attributes differ 337 largely in scale. The following sections describe different rep- 338 resentations of priority functions used within hyper-heuristics 339 to combine the individual attributes. 340

C. Representations of Priority Functions

The choice of representation is very important as it determines the range and complexity of the priority functions 343 (or indices) that can be generated by the hyper-heuristic. 344

Туре	Attribute	References
Job	(Imminent) operation processing time	[21], [22], [23], [24], [25], [26], [27], [29], [30]
		[31], [32], [33], [34], [35], [36], [37], [38], [39]
		[40], [41], [42], [43], [44], [45], [46], [47], [48]
		[49], [50], [51], [52], [53], [54], [55], [56], [58]
		[59], [60], [62], [64], [65]
	Processing time of subsequent operation(s)	[25], [36], [37], [44], [49], [51], [52], [53], [54
	Sum of processing times of remaining operations	[25], [26], [31], [35], [37], [38], [39], [40], [41
		[42], [43], [44], [45], [46], [47], [48], [49], [50]
		[51], [52], [53], [54], [55], [56], [57], [58], [60]
		[61], [62], [63], [64], [65], [66], [67]
	Weighted (linearly decreasing) sum of processing times of remaining operations	[66]
	Number of remaining operations	[26], [31], [35], [38], [39], [40], [41], [42], [43
		[44], [45], [47], [48], [49], [50], [51], [52], [53
		[54], [55], [56], [59], [63], [64], [65]
	Release date	[23], [24], [25], [26], [27], [29], [30], [31], [32]
		[33], [34], [44], [48], [51], [52], [53], [56], [60]
		[61], [62], [64], [66]
	Arrival time at considered work centre	[26], [31], [40], [44], [45], [46], [47], [48], [49]
		[50], [51], [52], [53], [54], [55], [56], [57], [61
		[64]
	Due date of job	[21], [22], [23], [24], [25], [26], [27], [29], [30]
		[31], [32], [33], [34], [35], [36], [38], [40], [41]
		[43], [44], [45], [47], [48], [49], [50], [51], [54]
		[55], [56], [58], [59], [60], [61], [62], [63], [64]
		[65], [66], [67]
	Due date of imminent operation	[44], [65]
	Weight	[21], [24], [26], [30], [31], [32], [33], [34], [38]
		[40], [41], [45], [49], [50], [54], [55], [56], [65]
		[66], [67]
	Setup time (given the current setup)	[31], [32], [33], [34], [63], [65]
	Number of machines that can process (imminent) operation	[59], [61], [63]
	Number of successor operations in precedence graph	[31]
	Level of (imminent) operation in precedence graph	[31]
Work	Sum of (imminent) operation processing times of all waiting jobs	[25], [27], [38], [41], [49], [56]
centre	Average (imminent) operation processing times of all waiting jobs	[25], [27], [51], [58], [65]
centre	Minimum processing time of (imminent) operations of waiting jobs	[25], [27], [49], [56]
	Maximum processing time of (imminent) operations of waiting jobs	
		[25], [27], [38], [41], [49], [56]
	Maximum sum of processing times of remaining operations of waiting jobs	[38], [41] [35], [36], [37], [38], [41], [47], [54], [55], [63]
	Number of waiting jobs	[25], [26], [27], [38], [41], [47], [54], [55], [63]
	Minimum due date of waiting jobs	[25], [27]
	Maximum due date of waiting jobs	[25], [27], [38], [41]
	Maximum weight of waiting jobs	[38], [41]
	Average setup time of waiting jobs (given the current setup)	[65]
	Maximum saving in setup time if (imminent) operation is processed on parallel machine	[65]
	Number of waiting jobs of same family	[27], [65]
	Fullness of batch (using currently waiting jobs)	[65]
	Speed of considered machine	[34]
	Sum of speeds of all parallel machines	[34]
Global	Sum of processing times of remaining operations of all jobs requiring considered work centre	[23], [26], [31], [34], [49], [56]
	Number of remaining operations of all jobs requiring considered work centre	[23], [26], [31], [34]
	WINQ = Sum of (imminent) operation processing times (adjusted for number of parallel	[36], [38], [41], [44], [48], [51], [52], [53], [57
	machines) of jobs waiting at next work centre (required for subsequent operation of a job)	[58]
	NINQ = Number of jobs waiting at next work centre	[36], [37], [38], [41], [54]
	Maximum WINQ of waiting jobs	[38], [41]
	Maximum NINQ of waiting jobs	[38], [41]
	Average waiting time of jobs last processed across all work centre in the shop	[54]
	Sum of processing times of remaining operations of all jobs requiring next work centre	[37]
	Sum of processing times of remaining operations of an jobs requiring next work centre Sum of (imminent) operation processing times of jobs waiting at critical work centre	[49], [56]
		[[], []], []]
	(the one with the greatest sum of processing times of all remaining operations)	[40] [56]
	Sum of processing times of remaining operations of all jobs requiring critical work centre	[49], [56]
	Sum of (imminent) operation processing times of jobs waiting at considered work centre	[49], [56]
	that still have to visit critical work centre	
		5 4 0 3 5 5 5 3
	Sum of (imminent) operation processing times of jobs waiting at work centre that still have to visit the work centre with currently largest queue (measured in processing time)	[49], [56]

³⁴⁵ For ease of presentation, priority indices in this paper are ³⁴⁶ defined so that a higher index I_j corresponds to a higher prior-³⁴⁷ ity of a job *j*. To illustrate, the priority index of the well-known minimum slack (MS) rule is given by

$$I_j^{\rm MS} = -\left(d_j - t - p_j^{\rm r}\right) \tag{1} \quad {}_{349}$$

³⁵⁰ where d_j denotes the due date of job *j*, p_j^r the processing time ³⁵¹ of the remaining operations of job *j*, and *t* refers to the cur-³⁵² rent time. Sections III-C1 and III-C2 discuss parametric and ³⁵³ grammar-based representations of priority functions, respec-³⁵⁴ tively. Examples of these two representations are compared ³⁵⁵ empirically in [52].

1) Fixed-Length Parametric Representations: One approach to encoding priority functions is to predefine their basic format and parameterize it. Then, a priority function can be represented by a vector of (real valued) parameter values. A simple and commonly used format is that of the weighted sum [35], [36], [45], [57], that is

362

$$I_j = \sum_{y=1}^{a} w_y x_{y,j} \tag{2}$$

³⁶³ where $x_{y,j}$ denotes one of the *a* attributes of job *j* provided to ³⁶⁴ the hyper-heuristic, and w_y refers to the corresponding weight. ³⁶⁵ To illustrate, if the weighted sum were used as the prede-³⁶⁶ fined format of priority functions and $x_{1,j} = d_j$, $x_{2,j} = t$, ³⁶⁷ and $x_{3,j} = p_j^r$, the priority function of the MS rule would be ³⁶⁸ encoded by the parameter vector $\mathbf{w} = (-1, +1, +1)$. Other ³⁶⁹ simple formats based on if-then-else rules have also been ³⁷⁰ proposed in [47] and [59].

Clearly, a representation based on a simple format such 371 372 as the weighted sum is often too restrictive to allow for the ³⁷³ discovery of the most effective priority functions (see [52]), which motivates the use of more complex representations, 374 375 e.g., based on artificial neural networks [41], [52]. On the 376 other hand, such representations lead to a significantly larger 377 search space, and also to priority functions that are so complex that they defy interpretation. The challenge is to choose 378 format that is as simple as possible without compromising 379 a 380 the ability of the hyper-heuristic to generate effective prior-³⁸¹ ity functions, which is difficult as the complexity required for given problem is normally unknown in advance. One study 382 a 383 that examines the impact of the flexibility of the representation on the results is [52]. 384

One advantage of a parametric representation is that search spaces of real-valued vectors are relatively common, implying the availability of a number of suitable optimization algorithms. In fact, Hooke–Jeeves pattern search [35], [36], simulated annealing [41], and EAs [45], [47], [52], [58], [59] have all been successfully used for the generation of scheduling heuristics based on parametric representations.

2) Variable-Length Grammar-Based Representations: An alternative way of defining the search space of priority funcindividual components can be assembled to yield a valid priority function. Fig. 2 gives an example of a grammar for the generation of priority functions, where the expressions on the left-hand side can be replaced by any of the expressions on the right-hand side (alternative options are separated by the "I" symbol). A specific priority function can then be represented by an expression tree, which is a popular representation the literature (see [27], [29], [31], [37]). Expression trees are composed of leaf nodes, representing terminals such as the attributes in this case, and internal nodes, representing the

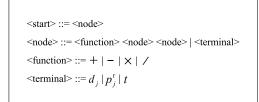


Fig. 2. Simple grammar for the construction of priority functions.

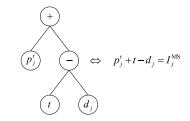


Fig. 3. Expression tree representing the priority function of the MS rule.

functions to combine the terminals with each other. Fig. 3 405 shows an expression tree that could be generated with the 406 grammar from Fig. 2, which encodes the MS priority function, where expression trees are decoded recursively, starting 408 from the root node, and from left to right. 409

The main reason for the popularity of grammar-based representations is that they allow for the generation of priority 411 functions of variable format and length without the requirement to define a basic format in advance. Apart from the 413 attribute (terminal) set, the only input that has to be provided 414 to the hyper-heuristic is a set of suitable functions that it can 415 use to combine the attributes.

Table IV lists the functions that have been used within ⁴¹⁷ hyper-heuristics. It shows that the four basic arithmetic ⁴¹⁸ operators are included in the function set in every of the ⁴¹⁹ reviewed papers, with the division either implemented as protected (returns 1 if divisor is 0) or unprotected (returns a ⁴²¹ very large number if divisor is 0). These operators allow ⁴²² for the reconstruction of many priority functions of the ⁴²³ most effective manually designed rules, which justifies their ⁴²⁴ selection. The function set is further often supplemented ⁴²⁵ with a ternary or quaternary if-then-else (ifte) operator, ⁴²⁶ defined as

ifte
$$(x_1, x_2, x_3) = \begin{cases} x_2 & \text{if } x_1 \ge 0\\ x_3 & \text{otherwise} \end{cases}$$
 (3) 428

and

ifte(x

$$x_1, x_2, x_3, x_4) = \begin{cases} x_3 & \text{if } x_1 \ge x_2 \\ x_4 & \text{otherwise} \end{cases}$$
(4) 430

429

and/or some common mathematical functions such as max or ⁴³¹ min. As in the case of the attribute set, a larger function set ⁴³² increases the size of the search space and the aim should thus ⁴³³ be to select only the most relevant functions. Against this ⁴³⁴ background, the value of including more complex functions ⁴³⁵ such as cos, sin, exp, log, pow, or sqrt that generally do not ⁴³⁶ occur in priority functions of effective rules from the literature, ⁴³⁷ and for which there is no theoretical justification, is question- ⁴³⁸ able. Moreover, some of the above functions can be expressed ⁴³⁹

486

TABLE IV Functions Used Within Hyper-Heuristics Operating on a Grammar-Based Representation of Priority Functions

Туре	Functions	References
Arithmetic	+, -, ×, /	[22], [23], [25], [26], [27], [29], [31], [32], [33], [34], [37], [40], [44], [43], [48], [49], [50], [51], [52], [53] [54], [55], [56], [60], [61], [62], [63], [64], [65], [66]
Logical	ifte (ternary) ifte (quaternary) > ≡ ∧, ∨, ¬	[32], [33], [44], [43], [49], [50], [51], [52], [53], [54], [55], [56], [61], [63], [65] [25], [26], [27], [40] [23], [43], [61], [63] [61] [43], [61] [43], [63]
Mathematical	max min avg sgn pos abs neg sin, cos exp log pow sqrt	

⁴⁴⁰ by means of other functions, e.g., $\max(x_1, x_2) = \text{ifte}(x_1 - x_1, x_2, x_1, x_2)$, $\min(x_1, x_2) = \text{ifte}(x_2 - x_1, x_1, x_2)$, $\operatorname{neg}(x_1) = 0 - x_1$, ⁴⁴² $\operatorname{abs}(x_1) = \max(x_1, \operatorname{neg}(x_1))$, raising the question of whether ⁴⁴³ those functions should be directly provided or whether it ⁴⁴⁴ should be left to the hyper-heuristic to reconstruct them in ⁴⁴⁵ case they are beneficial.

Although one advantage of the grammar-based representa-446 447 tion is that the complexity of the resulting priority indices 448 is potentially unbounded, in practice many researchers have 449 bounded the complexity and search space by limiting the 450 maximum tree depth. Unfortunately, there are no theoreti-451 cal guidelines on the determination of an adequate maximum 452 tree depth for the evolution of priority functions. If it is 453 chosen too small, some high quality heuristics might not 454 be representable and thus the quality of the solutions the 455 algorithm can find is limited. On the other hand, if it is 456 chosen too large, the hyper-heuristic may get lost in the vast search space. The depth used in previous studies varies 457 458 between 6 and 17 [22], [34], [40], [48], [49], [62], [65]. 459 Jakobović and Marasović [31] tested their hyper-heuristic with 460 values ranging from 9 to 17 and find that a maximum tree depth of 14 leads to the best results. However, the best value is 461 ⁴⁶² likely to depend on the given problem as well as the employed 463 optimization algorithm in general.

In addition to job and shop attributes, many researchers have included some (random) constants in the terminal set of their 466 hyper-heuristics [23], [25], [33], [37], [44], [56], [65], [66].

⁴⁶⁷ This enables the hyper-heuristic to weigh attributes differently, ⁴⁶⁸ especially since the latter have different units and can be of ⁴⁶⁹ different magnitude.

⁴⁷⁰ Search spaces of expression trees are typically explored by ⁴⁷¹ means of genetic programming (GP), which is also the pre-⁴⁷² dominant optimization algorithm employed in the literature ⁴⁷³ for the evolution of scheduling heuristics (see [25], [49]). An ⁴⁷⁴ exception by Nie *et al.* [29], [48], [64], who used gene expres-⁴⁷⁵ sion programming (GEP) instead. In [29], they compare their ⁴⁷⁶ GEP hyper-heuristic to a GP hyper-heuristic and report that the former generates slightly better priority functions in most 477 cases, and in much less time. Moreover, the priority func- 478 tions evolved by GEP are shown to be relatively simple and 479 easy to understand, whereas GP has the tendency to evolve 480 unnecessarily large expression trees (see [22], [44], [51], [61]). 481 This phenomenon, also referred to as bloating, is generally 482 undesired as it increases the runtime of GP and leads to priority functions that are more complex but not necessarily more effective. 485

D. Set of Eligible Jobs

In general, a dispatching rule does not only specify a pri- 487 ority function but also the eligible job set, i.e., the jobs from 488 which the next job to be scheduled can be selected. Most dis- 489 patching rules only consider jobs eligible for scheduling that 490 are already waiting at the given work center. This implies that 491 a machine is never left idle if there are jobs waiting to be 492 scheduled, which is also referred to as nondelay scheduling. 493 While nondelay scheduling is generally effective, as it mini- 494 mizes delays due to idle times, it is not necessarily optimal. In 495 fact, it can be beneficial to leave a machine idle in some situa- 496 tions, e.g., in anticipation of a high priority job arriving in the 497 near future that should be processed without delay. In order 498 for dispatching rules to be able to take such a decision, the 499 eligible job set has to be extended to include also jobs arriving 500 in the future. 501

In most of the hyper-heuristics designed to evolve dispatching rules the eligible job set is restricted to waiting jobs [25], [35], [36], [40], [41], [51], [62], [65]. Another common setting is to also include future jobs that are expected to arrive before the shortest operation of waiting jobs can be completed [29], [31], [37], [45], which is in the spirit of the Giffler and Thompson algorithm [78]. Hildebrandt *et al.* [44] tested their hyper-heuristic with and without inclusion of future jobs in the set of eligible jobs and find that the best rule evolved with the extended job set outperforms the best still ⁵¹⁶ account and allow for decisions to leave a machine idle. ⁵¹⁷ Instead of providing a hyper-heuristic with a fixed definition ⁵¹⁸ of the eligible job set, the latter can also be optimized by ⁵¹⁹ the hyper-heuristic itself. Nguyen *et al.* [49] let their hyper-⁵²⁰ heuristic optimize a parameter called the nondelay factor that ⁵²¹ controls the extent to which future jobs are included in the ⁵²² eligible job set. In another paper, Nguyen *et al.* [50] designed ⁵²³ a hyper-heuristic that evolves a separate function (of some ⁵²⁴ shop attributes) for the nondelay factor, which can then adapt ⁵²⁵ to changing shop conditions.

526 E. Evaluation of Candidate Heuristics

In order to know whether a candidate heuristic is effective or not, unsupervised hyper-heuristics need to obtain an estimate of the performance of that heuristic by applying so it to some training instances. The quality of the solutions generated by the heuristic for these instances then determines its fitness, which in turn governs the search behavior of the hyper-heuristic. Hence, in evaluating candidate heuristics, two important decisions to be made concern the selection of training instances and the definition of the fitness function, which are discussed in Sections III-E1 and III-E2, respectively.

1) Training Instances: For reasons of simplicity, a training 538 539 instance is defined in this paper as an instance that pro-540 vides a measure of performance for a given heuristic. This 541 includes static problem instances as well as runs of a stochas-542 tic simulation. Clearly, whether to use static instances or 543 stochastic simulation for the evaluation of candidate heuris-544 tics depends on which problems the heuristic is supposed solve once it has been generated. This is illustrated by 545 to 546 Hildebrandt et al. [44], who tested the dispatching rules from Tay and Ho [62], which have been trained on and shown to 547 548 be effective for static instances, in a long-term simulation with dynamic job arrivals and find that they perform poorly. 549 550 Nguyen et al. [49], also examined the effectiveness of dispatching rules that have been evolved using static instances 551 ⁵⁵² in a long-term simulation. They report that the evolved rules perform well if the shop utilization is equal or less than 80% 553 554 but become worse than some benchmark rules as utilization 555 increases beyond that value. They attribute this to the fact that 556 static instances reflect conditions of low utilization, in which 557 few new jobs arrive over time. Furthermore, the relative per-⁵⁵⁸ formance of evolved scheduling heuristics has been shown to 559 deteriorate with an increasing deviation between the test and ⁵⁶⁰ training instances in terms of job processing orders [37], num-⁵⁶¹ ber of jobs [22], and due date setting [32], [33], [65]. These ⁵⁶² results emphasize the importance of using a set of training ⁵⁶³ instances that reflect the problems the heuristics are likely to 564 encounter in their future use.

Another important factor with regard to the training set is its size, i.e., the number of training instances. If a small training set is chosen, the evolved heuristics are likely to overfit the training instances and not perform well on the 568 unseen test instances, which implies that their reusability 569 is very limited. On the other hand, a larger training set 570 increases the runtime of the hyper-heuristic, without necssarily leading to better heuristics. Geiger and Uzsoy [27] 572 reported that the performance of their hyper-heuristic improves 573 as the number of training instances approaches 10 but does 574 not improve further with a larger training set. In contrast, 575 Jakobović and Marasović [31] found that their hyper-heuristic 576 performs best with the largest of the tested settings (100 training instances), which indicates that the best training set size 578 is highly problem-specific and has to be determined through 579 pilot experiments. 580

Some researchers have argued that a hyper-heuristic can 581 either be used to generate a heuristic that performs reason- 582 ably well for a number of related problems or one that is 583 very effective for the specific problem it has been tailored to, 584 and ineffective otherwise [66], [79]. This argument cannot be 585 supported from a theoretical perspective, as there is no rea- 586 son why two specialized heuristics could not be combined 587 by a hyper-heuristic into one heuristic that analyzes the char- 588 acteristics of a given problem and applies the (specialized) 589 heuristic that is most suitable for it. On the other hand, the 590 generation of more sophisticated heuristics certainly poses a 591 challenge to hyper-heuristics up to an extent where the under- 592 lying relations are merely too complex to be discovered by the 593 hyper-heuristic. In any case, the generation of heuristics that is 594 supposed to perform well on a wider range of problems cer- 595 tainly requires a larger and more heterogeneous training set 596 that covers this problem range, which in turn increases the 597 runtime of the hyper-heuristic. 598

Note that, if the set of problem instances is very large 599 or randomly generated (as is usually the case if a stochas- 600 tic simulation is used for evaluating a dynamic problem) then 601 computational limitations make it necessary to restrict evalu- 602 ation to a subset (sample) of all possible training instances. 603 Effectively, due to the sampling, the fitness function then 604 becomes stochastic. In such cases, to reduce the sampling 605 variance, it has been recommended to evaluate all solutions 606 competing for survival within a generation by the same sub- 607 set of problem instances, while changing the subset from 608 iteration to iteration to make sure individuals that survive 609 several generations are tested on a large variety of problem 610 instances [23], [44], [65]. Furthermore, if the problem used 611 for evaluation consists of a sequence of random jobs dynam- 612 ically generated over time, at least in principle there are two 613 ways to improve the accuracy of evaluation: by increasing 614 the number of problem instances tested, or the number of 615 jobs considered in each problem instance. In such cases, to 616 get a proper estimate of the steady-state behavior of a solu- 617 tion, it is also necessary to discard the first jobs as warm-up 618 period. 619

2) Fitness Function: The application of a scheduling 620 heuristic \mathcal{H} to a number of training instances T = 621 {1, 2, ..., |T|} results in performance measures $z_i(\mathcal{H})$, the 622 objective value reached by the heuristic on instance *i*. These 623 measures have to be integrated by means of a fitness func- 624 tion f(.) to determine the overall fitness of the heuristic. 625

626 The following fitness functions have been proposed in the 627 literature.

⁶²⁸ 1) Sum [or average] of objective values [22], [23], [31], ⁶²⁹ [32], [41], [48], [52]

$$f(\mathcal{H}) = \left[\frac{1}{|T|}\right] \sum_{i=1}^{|T|} z_i(\mathcal{H})$$

⁶³¹ 2) Average relative objective value [44]

$$f(\mathcal{H}) = \frac{1}{|T|} \sum_{i=1}^{|T|} \frac{z_i(\mathcal{H})}{z_i(\operatorname{ref})}.$$

⁶³³ 3) Sum [or average] of relative deviations [33], [49], [66]

$$f(\mathcal{H}) = \left[\frac{1}{|T|}\right] \sum_{i=1}^{|T|} \frac{z_i(\mathcal{H}) - z_i(\text{ref})}{z_i(\text{ref})}$$

where z_i (ref) denotes a reference objective value for instance *i*, obtained by some other solution method. Which fitness aggregation is most desirable depends very much on the intentions of the designer of the algorithm.

The sum (or average) of objective values concentrates on 640 erforming well on problem instances with a large potential 641 642 for improvement. To illustrate this, consider a problem where the objective is to minimize the mean tardiness of all jobs so 643 that the fitness is computed as the sum of mean tardiness val-644 ues obtained from all training instances. If there is one training 645 646 instance in the set with a very tight due date setting, the tar-647 diness of this instance will be much higher than that of other 648 training instances and therefore strongly correlate with the fit-649 ness value. In consequence, the hyper-heuristic will focus its 650 search on heuristics that can solve well this particular instance while largely ignoring their performance on other instances. 651

Alternatively, one can use the average relative objective 652 value or the sum (or average) of relative deviations as the 653 654 fitness function, which are equivalent for hyper-heuristics 655 that operate on the ranks of fitness values rather than the values itself. These fitness functions reduce the weight of 656 657 difficult training instances by relating the objective value of 658 each instance to a reference value before combining them. Their disadvantage is that they require good reference values, 659 660 which are generally only available for well-studied benchmark instances from the literature for which (near)-optimal solu-661 662 tions are known. In all other cases, reference values have to 663 be obtained, typically by applying some benchmark heuris-⁶⁶⁴ tic(s) to the problem [44], [49], [66], which may or may not 665 yield good results.

Another issue that arises only if candidate heuristics are evaluated with long simulation runs is that some heuristics, which may be present particularly at the start of the run of her hyper-heuristic, can lead to an unstable system, i.e., the number of jobs in the shop grows steadily. The fitness of these inferior heuristics may then be never obtained and excessive time wasted in the attempt. To prevent this from happening, Hildebrandt *et al.* [44] proposed to monitor the number of jobs in the shop during a simulation run and abort the run if a preset threshold value for the number of jobs is exceeded. The fitness of these heuristics is then largely reduced by a penalty, which 676 ensures that they are quickly discarded. In a follow-up paper, 677 Branke *et al.* [52] showed that this measure reliably detects the 678 inferior heuristics without prematurely stopping the evaluation 679 of good heuristics. 680

IV. ISSUES AND CHALLENGES

The research on hyper-heuristics for the automated design 682 of scheduling heuristics is still in an early stage and there 683 remain a number of open questions and challenges. The following sections discuss some of the main issues that future 685 work should focus on. 686

A. Evolving Sets of Heuristics

A common theme of the existing studies on hyper-heuristics 6688 is that they predominantly address simple scheduling environments and/or employ a hyper-heuristic to evolve a single 690 scheduling heuristic. However, in more complex scheduling 691 environments, there may be a number of interrelated deci-692 sion problems that have to be resolved, e.g., the formation 693 and scheduling of batches in the presence of batch process-694 ing machines, the assignment of operations to machines and 695 scheduling of these machines in parallel machine environ-696 ments, or the coordination of resources in environments with 697 multiple resource constraints. This raises the question of how 698 to best deal with such more complex scenarios.

The most straightforward approach seems to be to design a 700 set of heuristics, one for each decision, and to simply encode 701 the set of heuristics as one individual. This approach is fol-702 lowed by Nie et al. [64], who developed a hyper-heuristic for 703 flexible job shop problems that operates on a search space 704 in which each individual consists of two functions, one for 705 routing, i.e., assigning operations to machines, the other for 706 sequencing, i.e., scheduling the machines. Their results show 707 that this hyper-heuristic can evolve sets of heuristics that 708 outperform the single sequencing heuristics evolved by a con-709 ventional hyper-heuristic (and combined with some benchmark 710 routing heuristic). On the other hand, the drawback of encod- 711 ing multiple heuristics as one individual is that the search 712 space of the hyper-heuristic grows exponentially in the num- 713 ber of heuristics, which limits the approach to the generation 714 of small sets of heuristics. 715

One possible solution to overcome the issue of search space 716 size may be the use of coevolution, which implies a division 717 of the search space into several sub-spaces handled separately 718 by different subpopulations. Nguyen *et al.* [56] designed a 719 coevolutionary hyper-heuristic, in which a subpopulation of 720 priority functions used for dispatching is coevolved with a 721 separate subpopulation of functions for due date assignment. 722 The hyper-heuristics, in which the two functions are repre-724 sented by a single individual. Another option is to generate 725 heuristics sequentially, in a similar way as proposed in [57]. 726 However, such hyper-heuristics implicitly assume that the rel-727 ative effectiveness of candidate heuristics at one stage is more 728 quent stages, and thus cannot be expected to perform very well 730

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731 if there are strong interdependencies between the subprob-732 lems. On the other hand, it may be sufficient to resolve several 733 subproblems using the same heuristic in some situations. To 734 illustrate, Park et al. [33] addressed an order acceptance and scheduling problem with a hyper-heuristic that evolves two 735 separate functions for the acceptance and dispatching of jobs. 736 They compare its performance to a hyper-heuristic evolving a 737 single function to handle both decisions and find that the latter 738 more effective, indicating that there is no need for a separate is 739 740 heuristic to deal with the acceptance of jobs, possibly due to the strong correlation underlying the two decisions. Hence, it 741 742 is important to carefully assess the interrelations between the subproblems to be solved prior to designing a hyper-heuristic. 743 In multimachine problems, it has been found beneficial to 744 745 use different dispatching rules at different machines if the lat-746 ter vary with respect to their positions in the shop [80]-[83] 747 and/or workload [84]-[86]. Consequently, some researchers 748 have developed hyper-heuristics that generate sets of machine-749 specific rules by selecting a (potentially) different rule for each work center from a number of given rules [65], [87], [88] 750 or by evolving several composite rules, where each rule is 751 tailored to a certain work center [25], [55], [57], [58]. The 752 search space then grows exponentially in the number of 753 work centers or machines, which can be very large in shops 754 755 of realistic size, requiring some specific measures in addition to the above techniques to deal with this problem. 756 757 Miyashita [40] proposed a hyper-heuristic that is based on predetermined classification of machines into bottlenecks а 758 759 and nonbottlenecks and evolves one rule for each class of 760 machines. Similarly, Jakobović and Budin [26] designed a 761 hyper-heuristic that optimizes the classification of machines while searching for good dispatching rules for each class. 762 763 More specifically, each individual consists of three functions, where one of them is a discriminating function of attributes 764 765 relating to the workload of a machine that determines which of 766 the two dispatching rules, encoded by the other two (priority) 767 functions, to apply. The best rule sets evolved by these hyper-⁷⁶⁸ heuristics are generally shown to outperform single benchmark rules. However, the results of the study by Nguyen et al. [50], 769 who examined the performance of the three-function hyper-770 heuristic by Jakobović and Budin [26] more closely, show that 771 772 the effectiveness among the evolved rule sets varies a lot more 773 than that among the single rules evolved by a conventional 774 hyper-heuristic. This indicates that the former struggles with 775 the more complex search space. In summary, there appears 776 to be some potential for future work on intelligent designs of 777 hyper-heuristics for more complex scheduling environments.

778 B. Attribute Selection and Construction

As discussed in Section III-B, a main challenge in designing an effective hyper-heuristic is to provide it with all the relevant problem attributes while excluding any redundant or irrelevant attributes. Otherwise, the search space could be either res too restrictive or unnecessarily large, which both hinder the hyper-heuristic in its ability to generate effective heuristics.

A few studies have performed some analysis of the best reference evolved heuristics in order to identify important attributes. Branke et al. [52] leaved out each of the attributes present in 787 the priority functions of their evolved dispatching rules one-788 by-one and examine the performance of these rules without the 789 respective attribute. Their analysis shows that some attributes, 790 specifically, those that also appear in the most effective rules 791 from the literature, are substantially more important for the 792 performance of the evolved rules than others. Eguchi et al. [41] 793 examined the first-order correlation between attribute values 794 and priority values (applying the best evolved dispatching rule) 795 to determine the relevance of attributes, and eliminate ineffec-796 tive attributes in this way. Nguyen et al. [49] conducted a 797 high-level analysis of the occurrence frequency of attributes 798 in the priority functions of the best dispatching rules evolved. 799 They find that the relevance of attributes is problem-specific 800 to some extent though some seem to be generally more impor- 801 tant than others. This highlights a major drawback of any 802 post-generation analysis, which is that the gained insights may 803 only be applicable to the given problem (class), for which an 804 effective heuristic has already been generated, and therefore, 805 be of limited value. Instead, future hyper-heuristics should 806 be designed to perform the tasks of selecting and construct- 807 ing suitable attributes automatically and simultaneously to the 808 optimization. 809

C. Interpretability and Trust

The interpretability of evolved heuristics is a crucial aspect 811 to gain the trust of users, i.e., operators or managers, par- 812 ticularly since hyper-heuristics are black box optimizers. 813 Unfortunately, there is some evidence that more complex 814 scheduling environments (for which the use of hyper-heuristics 815 is most promising) often require heuristics of a certain com- 816 plexity so that simply choosing an easy-to-interpret representa- 817 tion will result in heuristics of comparatively low quality [52]. 818 On the other hand, it may be possible to allow for open- 819 ended evolution and still search for the simplest representation 820 of a well-performing heuristic, or to generate good tradeoffs 821 between quality and interpretability, by means of multiobjec- 822 tive methods, with heuristic complexity being one objective to 823 be minimized. Online rule simplification techniques [89]-[91] 824 can also be applied to improve the readability of evolved 825 heuristics. 826

In many cases, it may be possible to simplify heuristics after they have been generated without significantly compromising their performance. This is particularly true for heuristics evolved on the basis of a grammar-based representation, which typically contain redundant components, e.g., if-thenelse operators where the condition is always true or false. The following simplication, these heuristics may then be analyzed manually and linked to some human-designed heuristics to facilitate interpretation [25], [27], [49], [56], [66]. However, fully understanding evolved heuristics is still a challenging task, especially when dealing with complex environments, which stresses the need for some methodological support.

In the literature, a few tools and methods to understand ⁸³⁹ the behavior of scheduling heuristics, specifically dispatching ⁸⁴⁰ rules, have been developed and used. Branke *et al.* [52] ana- ⁸⁴¹ lyzed dispatching rules by visualizing their priority indices ⁸⁴²

843 as functions of the attributes they incorporate. Clearly, such 844 a visualization is only possible for a very limited number of 845 attributes. Branke and Pickardt [92] proposed a method that 846 identifies weaknesses in the decision logic of a given dis-847 patching rule. All in all, future research on hyper-heuristics 848 should place more emphasis on the issue of interpretability of 849 heuristics and controlling or reducing their complexity.

850 D. Comparison of Algorithms

In order to give recommendations on when it is ben-851 852 eficial to use a hyper-heuristic and how to design it, 853 extensive and meaningful performance comparisons of 854 evolved heuristics with more sophisticated (global) solu-855 tion algorithms as well as between different hyper-856 heuristics are needed. So far, such comparisons have 857 been rather limited (see [28], [32], [33], [45], [49], [56] and 858 [29], [43], [45], [49], [52], respectively). Intuitively, hyper-⁸⁵⁹ heuristic approaches have strengths compared to global opti-860 mization approaches in particular in dynamic and stochastic 861 environments where a quick reaction is important. But as observed in [93], they also become more competitive as 862 863 the problem size (and thus the search space for the global optimizer) increases. 864

One reason for the limited number of comparisons may be 865 866 that hyper-heuristics possess several properties that make a fair ⁸⁶⁷ comparison particularly difficult. For example, not only are the ⁸⁶⁸ hyper-heuristics stochastic algorithms with many parameters 869 to tune, but also is the evaluation function often a stochas-870 tic simulation, resulting in stochastic fitness values. Also, the running time for the simulations can be quite substantial, and, 871 872 to make things worse, the running time to evaluate a particu-873 lar dispatching rule strongly depends on the rule itself, as the 874 time to calculate the priority value and the number of jobs 875 in the system depends on the rule itself. This implies that a 876 comparison of hyper-heuristics based on the same number of 877 function evaluations has limited validity.

Irrespective of the challenges faced, an important pre-878 879 requisite for systematic algorithm comparisons is the avail-880 ability of suitable benchmark problems and algorithms. For 881 reusable heuristics, it is further important to clearly dis-882 tinguish between training and test problem instances-the ⁸⁸³ hyper-heuristic may use the training instances during opti-884 mization, while the generated heuristics have to be tested 885 on a separate, previously unseen set of test instances. While 886 libraries exist for static, deterministic scheduling problem ⁸⁸⁷ instances, e.g., the OR-Library [94] (which has also been used to test hyper-heuristics [33], [50]), the most promising ⁸⁸⁹ applications for hyper-heuristics seem to be in the area of 890 dynamic, stochastic problems, which are much more diffi-⁸⁹¹ cult to define. A possible benchmark are the dynamic job and ⁸⁹² flow shop problems designed by Rajendran and Holthaus [95] ⁸⁹³ and Holthaus and Raiendran [96] for the purpose of compar-⁸⁹⁴ ing dispatching rules from the literature. We have used these ⁸⁹⁵ problems in several hyper-heuristic studies [44], [51]–[53], ⁸⁹⁶ and have published some results online [97]. Still, especially ⁸⁹⁷ for more complex dynamic scheduling problems, the publica-⁸⁹⁸ tion of entire simulators (see jasima [97]) would greatly help replicability and facilitate fair comparisons. Also, the generated heuristics should be published in addition to the obtained results, ideally in a format that can be directly plugged into a simulator. 902

E. Computational Time

A major drawback of hyper-heuristics based on unsupervised learning is their high computational requirements. Even though the obtained heuristics typically can be executed very fast, a run of the hyper-heuristic itself can last many hours, especially if the evaluation of the many candidate heuristics, evolved during the search, involves extensive simulation runs. Measures to reduce the computational time of unsupervised hyper-heuristics should consequently focus on the fitness evaluations, which usually take up the most time by far.

Some approaches have been proposed to reduce the time 913 spent on evaluations. References [44], [52], and [65] equipped 914 their hyper-heuristic with a mechanism that monitors the num- 915 ber of jobs in the shop to detect heuristics that cause an 916 unstable system (see Section III-E2) and terminates the eval- 917 uation of heuristics once a preset threshold value is exceeded. 918 This idea could be extended to save time on the evaluation 919 of other inferior candidate heuristics by monitoring similar 920 values, e.g., the objective value after a predefined number of 921 completed jobs. Branke et al. [52] suggested a duplicate detec- 922 tion technique to avoid evaluating the same candidate heuristic 923 twice. Thereby, two priority functions are considered equiv- 924 alent if they provide the same ranking on a set of randomly 925 generated "dummy" operations. Hildebrandt and Branke [53] 926 investigated the use of surrogate models to approximate the 927 fitness of candidate heuristics, i.e., dispatching rules, in a GP 928 hyper-heuristic. By employing a distance measure based on 929 the behaviors of rules, the proposed surrogate-supported GP 930 hyper-heuristic can reduce the computational cost and improve 931 the convergence speed, which indicates that the development 932 of surrogate models is a promising direction for future work. 933

As previously discussed, another aspect that influences the ⁹³⁴ computational time is the complexity of candidate heuristics. ⁹³⁵ It is well-known that optimization algorithms that operate on ⁹³⁶ a variable-length grammar-based representation, such as GP, ⁹³⁷ are liable to bloating [98], i.e., they gradually evolve larger ⁹³⁸ and more complex individuals that are not necessarily better, ⁹³⁹ but require more time to be evaluated. Thus, controlling or ⁹⁴⁰ reducing the complexity of the heuristics that are evolved is ⁹⁴¹ also important for efficiency reasons, and the development of ⁹⁴² effective bloating control [99], [100] and online program sim-⁹⁴³ plication techniques [89]–[91] should also be of concern to ⁹⁴⁴ hyper-heuristic research in the future.

F. Overfitting and Robustness

Given the high computational requirements of unsupervised hyper-heuristics, it is highly desirable that the resulting 948 heuristics are reusable. However, like other forms of machine 949 learning, hyper-heuristics carry the risk of generating heuristics that overfit the problem instances used in the training 951 stage and perform poorly on all other instances, limiting their 952 reusability. In fact, overfitting has been observed in connection 953

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⁹⁵⁴ with hyper-heuristics by several researchers, in particular when ⁹⁵⁵ more complex representations [40], [49] and/or small sets of ⁹⁵⁶ training instances [33] are used. Hence, this issue should be ⁹⁵⁷ taken into account when designing a hyper-heuristic.

A concept closely related to overfitting is that of the 958 959 robustness of heuristics, i.e., their ability to cope with unfore-960 seen changes in the scheduling environment. Clearly, if the performance of the heuristics evolved was to strongly dete-961 ⁹⁶² riorate in the event of a minor change, e.g., in the job ⁹⁶³ arrival pattern, this would question the practicality of the ⁹⁶⁴ approach. So far, a few studies have examined the robust-965 ness of dispatching rules evolved by various hyper-heuristics. ⁹⁶⁶ Their results can be summarized in that evolved rules show be reasonably robust, including to changes in the num-967 to ⁹⁶⁸ ber of machines [44], processing time distribution [44], [51], ⁹⁶⁹ job arrival pattern [44], [65], shop utilization [51], [52], [65], 970 and due date setting [51], [65]. These studies can be fur-971 ther extended, e.g., by examining the limits of the robustness 972 of evolved heuristics, i.e., when they become worse than 973 benchmark heuristics. On the other hand, if the changes 974 to the scheduling environment are more pronounced, there 975 is always the option to simply rerun the hyper-heuristic to 976 generate a new heuristic for the altered problem. How to 977 determine the point at which rerunning the hyper-heuristic 978 becomes beneficial is another challenging question to be 979 investigated.

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V. CONCLUSION

In recent years, hyper-heuristics have demonstrated their ability to automatically generate very competitive heurisgenerate heuristics automatically, it becomes feasible to tailor heuristics to the specific production environment, and to change them quickly whenever the environment changes. In this sense, hyper-heuristics have the potential to revolutionize production scheduling as they allow problem-specific heuristics to be applied successfully in settings where the traditional way of a human expert designing heuristics would be too expensive, or simply too time consuming.

This paper constitutes the first comprehensive review 993 of hyper-heuristics for the automated design of produc-994 tion scheduling heuristics, providing a simple taxonomy and 995 focussing on key design choices such as the learning method, 996 attributes, representation, and fitness evaluation. Moreover, a 997 number of the issues and challenges that should be addressed 998 in the future have been discussed, including the generation 999 of rule sets, algorithm comparison, interpretability of the 1000 resulting heuristics, computational time, and overfitting and 1001 robustness.

The review is aimed for researchers as well as practitiontional ers. Researchers who aim to further advance the technique total of hyper-heuristic scheduling are provided with a compretheory hensive review of the state-of-the-art and a discussion of the open issues suitable for future work. Also, we have total established a website that may serve as a starting point for future algorithm comparisons on dynamic, stochastic benchtor mark problems. Practitioners in scheduling, on the other hand, can use this paper to compose a suitable hyper-heuristic and 1010 make the appropriate design choices for their particular appli- 1011 cation. This paper contains guidelines, for example, on how 1012 to select attributes, what fitness function to choose, and what 1013 representation might be the most appropriate.

Currently, the vast majority of papers fall into the category 1015 of unsupervised learning with open-ended grammar-based 1016 evolution. Clearly, some of the less explored areas may deserve 1017 more attention, and the work reviewed here may benefit from 1018 cross-fertilization also with other hyper-heuristic concepts, 1019 such as hyper-heuristics for heuristic selection [3]. Whereas 1020 the approaches to select dispatching rules mostly use machine 1021 learning algorithms such as artificial neural networks, deci-1022 sion trees or support vector machines, heuristic generation 1023 approaches mostly apply heuristic search methods such as EAs 1024 or tabu search, and so far there is very little overlap between 1025 the two areas. A first example that combines heuristic gener- 1026 ation and heuristic selection (but both based on EAs) can be 1027 found in [65]. Another promising direction may be to auto- 1028 matically gear metaheuristics to a particular problem domain, 1029 such as in [101]. 1030

Finally, our review has focussed on the generation of 1031 reusable heuristics for production scheduling. The research in 1032 this area may benefit from ideas developed for hyper-heuristics 1033 in related problem domains like timetabling [16], [17] or 1034 project scheduling [20]. Given their potential and the various 1035 open problems, research on hyper-heuristics for the design of 1036 production scheduling heuristics is likely to continue yet for 1037 some time.

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